Abstract—An energy-efficient orthogonal multiple access relay channel (MARC) system is developed where accumulator and differential detection (DD) are used at the source and relay nodes, respectively. However, the weak decoding capability of DD degrades the frame error rate (FER) performance of the orthogonal MARC system if the conventional decoded-and-forward relaying strategy is applied. In this paper, a novel joint adaptive network-channel coding (JANCC) technique is proposed to support DD by making efficient use of the erroneous estimates output from DD. In the JANCC technique, the destination constructs a vector identifying the indexes of the source nodes whose information parts contain errors, and sends it to the relay to request a retransmission. The relay performs network coding by taking the exclusive-OR (XOR)-operation only over the stored estimates specified by the identifier vector, which aims to avoid unnecessary erroneous estimates being coded. In addition, a bit-flipping probability \( p_{\text{nc}} \) is obtained between the two sequences, one is the network-coded sequence sent from the relay, and the other is their corresponding XOR-ed information sequence. The decoding algorithm of JANCC exploits the probability \( p_{\text{nc}} \) at the destination to update the Log-Likelihood Ratio during the iterative decoding process. Hence, the information sequences received at the destination are able to be recovered even though the redundancy forwarded from the relay is generated from the erroneous estimates. Compared with the system where the iterative decoding is performed at the relay, the utilization of DD significantly reduces the computational complexity, which leads to meaningful power saving with only a small loss in the FER performance.

Index Terms—Multiple Access Relay Channel, different detection, automatic repeat request, joint channel-network coding.

I. INTRODUCTION

COOPERATIVE wireless networks have attracted a lot of attention of the wireless communication research community recently, since they provide a lot of design flexibility in the form of, e.g., diversity-multiplexing tradeoff (DMT), coverage extension, and multiple users’ Quality-of-Service (QoS) management. One of the cooperative network structures is Multiple Access Relay [1], which consists of \( N \) source nodes, one relay and one common destination; the role of the relay is to assist the source nodes to improve the probability of successful transmission to the destination. Many excellent joint network-and-channel coding (JNCC) techniques for orthogonal multiple access relay channel (MARC) with \( N = 2 \) can be found in the literature [2]–[7]. However, only a few publications consider orthogonal MARC systems with an arbitrary number of source nodes in their designs. In many cases the designs for orthogonal MARC systems with \( N \) greater than two require higher complexity in structure.

In addition, the powerful codes, such as turbo codes [8] or low density parity check (LDPC) codes [9], are applied at the relay in the most of the techniques for network designs [2]–[5], [10]. Although the application of powerful codes guarantees the high probability of successful decoding at the relay, the application causes heavy computation load; furthermore, as the number of source nodes increases, due to the heavy computation for the signal processing of each source node, power consumption is also expected to increase, which reduces the battery longevity of the relay.

In the conventional decoded-and-forward (DF)-based orthogonal MARC network designs [11], [12], the relay only forwards the successfully decoded information data transmitted from the source nodes, and discards the sequence estimates containing errors to avoid the error propagation in the JNCC decoding process. However, discarding the erroneous sequence estimates obtained at the relay degrades the transmission efficiency, because the estimates still contain a lot of useful information, which are helpful for reconstructing the transmitted signals from the source nodes at the destination. Hence, several emerging design schemes have been proposed [13]–[15] for preserving the information of the erroneous estimates at the relay, and forward the estimates in an analog form or their quantized versions. However, even with proper quantization techniques, relaying the signals requires more bandwidth for the representation of the soft bits.

Motivated by the disadvantageous properties of the previous works, an energy-efficient and low-complexity orthogonal MARC system combined with automatic repeat request (ARQ) [16] was developed in our previous work [17], where an accumulator (ACC) [18], equivalent to a differential encoder, is used at each source node after a recursive systematic convolutional (RSC) encoder. The use of the ACC is to ensure that the extrinsic information transfer (EXIT) convergence tunnel open until a point very close to the (1.0, 1.0) mutual information point [19]. In the broadcast (BC) phase, a very simple differential detection (DD) is performed at the relay.
and the information part is only extracted [20] from the output of the differential detector; the iterative decoding between the decoders of the RSC code and ACC is not performed.

However in [17], only two source nodes were assumed in the system model. In addition, although we have observed the bit-flipping probability $p_{nc}$ between the network-coded sequence sent from the relay and their corresponding XOR-ed information sequence, the aid of higher layer protocol setting is still needed for the destination to acquire the knowledge of the probability $p_{nc}$, which is less practical. Furthermore, all the information estimates output from the DD were network coded at the relay, where unnecessary erroneous estimates may be coded, which increases the bit-flipping probability $p_{nc}$, and thus, the correction capability of the decoder at the destination is degraded.

This paper aims to further improve our previous work in [17]. The system model in [17] is first generalized with an arbitrary number of source nodes. We then propose a novel joint adaptive network-channel coding (JANCC) technique, where after completing the decoding of all the information sequences transmitted from the source nodes, the destination constructs a vector identifying the indexes of the source nodes whose information parts contain errors. The identifier vector is then sent to the relay to request a retransmission of the information sequences decoded in error at the destination, and the relay performs network coding by taking the XOR-operation only over its received estimates of the information sequences specified by the identifier vector. The decoding of JANCC is performed at the destination, where the bit-flipping probability $p_{nc}$ is estimated by using the proposed algorithm and effectively utilized in the decoding process. Finally, the reduction in terms of the computational complexity is analyzed.

The main contributions of the paper are summarized as follows.

- By making the efficient use of the erroneous estimates of the information sequences received at the relay and by forwarding the coded version of them to the destination as additional redundancy, it is possible to correctly recover the information estimates transmitted via the source-destination (SD) links, which improves the frame error rate (FER) performance over the conventional DF relay strategies for MAC relaying systems.
- Lower bit-flipping probability $p_{nc}$ can be obtained with adaptively performing network coding on the estimates extracted from the DD, which enhances the decoding capability of the JANCC and thus FER performance can be improved.
- The computational complexity of the relay is significantly reduced by replacing the iterative decoding of an iteratively-decodable powerful codes by simply extracting the information (systematic) part output from DD. This advantageous characteristic is more meaningful, when the number of source nodes increases.
- The energy consumption for the channel estimation can be eliminated at the relay by utilizing DD because it does not need channel estimation. This advantageous characteristic is, again, more meaningful, when the number of source nodes increases.

The rest of this paper is organized as follows. Section II introduces the developed orthogonal MARC system model. A detailed description of the proposed JANCC and its corresponding decoding techniques are presented in Section III. Section IV presents simulation results to demonstrate the performance of the developed system combined with the proposed JANCC techniques, in terms of average FER and the computational complexity. The average throughput evaluation results are also provided in Section IV. Finally, Section V concludes this paper with some concluding remarks.

II. SYSTEM MODEL

Fig. 1 depicts a block diagram of the developed orthogonal MARC system assumed in this paper, where there are $N$ source nodes, one common relay node $R$, and one common destination node $D$. Each node is equipped with a single antenna, and $R$ is assumed to be connected with $D$ via a half-duplex link. As shown in Fig. 1, each source node $S_i$, $i = 1, 2, ..., N$, generates its binary and cyclic redundancy check (CRC) [21] encoded information sequence $u_i = \{u_i(k)\}_{k=1}^K$. All the information sequences are assumed to be statistically independent.

Each sequence $u_i$ is interleaved and encoded by a serially concatenated convolutional code (SCCC) where the SCCC is composed of a rate-$K/M$ RSC encoder $C_S$ and rate-1 ACC. The corresponding coded bit sequence $x_i = \{x_i(m)\}_{m=1}^M$ is modulated by binary phase-shift keying (BPSK) and broadcasted to $R$ and $D$ at the each source node’s dedicated independent time slot.

Without loss of generality, all transmitted signals have unit power. All the links are assumed to suffer from block Rayleigh fading, and thus the signal vectors received at $R$ and $D$ are

$$y_{iR} = h_{iR} \cdot x_i + n_{iR}$$
$$y_{iD} = h_{iD} \cdot x_i + n_{iD}$$

where $i = 1, 2, ..., N$, where $h_{iR}$ and $h_{iD}$ indicate the channel coefficients of $SR$ and $SD$ links with the source node $S_i$. $n_{iR}$ and $n_{iD}$ indicate the vectors of independent zero-mean complex additive white Gaussian noise (AWGN) at $R$ and $D$ with the source node $S_i$, respectively, with variance $\sigma_{iR}^2 = \sigma_{iD}^2 = \sigma^2$ per dimension. Block Rayleigh fading is assumed, with which $h_{iR}$ and $h_{iD}$ are assumed to be constant over one coded sequence but vary independently transmission-by-transmission and link-by-link.

If $D$ identifies that all users’ information sequences are successfully decoded by the CRC error detection, acknowledgement (ACK) is sent back to the source nodes, and the source nodes broadcast their following information sequences at their dedicated time slots. Otherwise $D$ sends to $R$ a redundancy retransmission request via the feedback channel\(^1\), and once receiving the retransmission request, $R$ transmits

\(^1\)The feedback channel is assumed to be error-free. This assumption is reasonable because the retransmission request can be represented by $N$ bits, as detailed in Section III, which is protected by using a powerful channel coding.
Joint adaptive network-channel coding (JANCC)

The received signal vector \( y_{iD} = \{y_{iD}(m)\}_{m=1}^{M} \) of the \( i \)th slot at \( D \) is demodulated to obtain the corresponding soft channel values, as

\[
L(y_{iD} | x_{i}) = 2 \cdot \mathbb{R} \left\{ \frac{y_{iD} h_{iD}^{*}}{\sigma^2} \right\}
\]

where the notations \( L(\cdot) \) and \( \mathbb{R}(\cdot) \) denote the Log-Likelihood Ratio (LLR)\(^3\) and a function that takes the real part of its argument, respectively, and \(^*\) indicates complex conjugation.

The soft channel outputs \( L(y_{iD} | x_{i}) \) for the \( i \)th slot, are iteratively decoded by the decoder of the SCCC with the Log-MAP algorithm [23]. The estimated sequence \( \hat{u}_i \) of \( u_i \) is obtained by making binary hard decision on the \textit{a posteriori} LLR \( \{L(u_i(k) | y_{iD})\}_{k=1}^{K} \) output from the decoder of \( C_S \), when no relevant increase in LLR is obtained by iterations.

The sequences \( \hat{u}_i \) are CRC-decoded for error detection after completing the iterative decoding of the SCCC, and the

\(^3L(x) = \ln \frac{P_{x}(x=1)}{P_{x}(x=0)}\), where \( P_{x}(\cdot) \) denotes the probability of its argument.
the updated LLR sequence for each user are stored in an $N$-bit identifier vector $\mathbf{F} = \{f(i)\}_{i=1}^{N}$, where

$$f(i) = \begin{cases} 0, & \tilde{u}_i = u_i \\ 1, & \tilde{u}_i \neq u_i \end{cases}$$

and a set $\mathcal{F}$ is defined as

$$\mathcal{F} = \{i : f(i) = 1\}. \quad (6)$$

If $\mathcal{F} \neq \emptyset$, $D$ sends to $R$ an redundancy-retransmission request with the vector $\mathbf{F}$ via the feedback channel and attempts to recover the information sequences decoded in error with additional redundancy forwarded from $R$.

As shown in Fig. 1, once receiving the retransmission request, the relay $R$ performs network coding by taking a sequence-by-sequence XOR-operation, notated as $\oplus$, over the received estimates of the information sequences specified by $\mathbf{F}$ even though the specified estimates may contain errors, as

$$\mathbf{u}_R = \bigoplus_{j \in \mathcal{F}} \tilde{u}_j. \quad (7)$$

The relay node $R$ then re-encodes the interleaved version of the network-coded bit sequence $\mathbf{u}_R = \{u_R(k)\}_{k=1}^{K}$ by using a rate-$K/M$ RSC encoder $C_R$, and finally, the corresponding RSC-coded bit sequence $\mathbf{x}_R = \{x_R(m)\}_{m=1}^{M}$ is forwarded to $D$ using BPSK modulation as additional redundancy.

**B. Decoding scheme and algorithm of JANCC**

Fig. 2 shows a factor graph representation of the proposed decoding scheme of JANCC, where $C_S^{-1}$ and $C_R^{-1}$, respectively, indicate the RSC decoders corresponding to $C_S$ and $C_R$ used in the source nodes and $R$. The notations $L_a$ and $L_e$ shown in Fig. 2 denote *a priori* and *extrinsic* LLR of their argument variables, respectively. Once $D$ receives additional redundancy forwarded from $R$, the decoding of JANCC is initiated to recover the information sequences decoded in error at $D$ via the $SD$ links.

As shown in Fig. 2, the LLR updating function $f_c(\cdot)$ [20] is utilized in the proposed JANCC decoding scheme to obtain the updated LLR sequence $\mathcal{L}'$, as

$$\mathcal{L}'(k) = f_c(\mathcal{L}(k), p_{nc}) = \ln \left( \frac{1 - p_{nc}}{1 - p_{nc} \cdot e^{\mathcal{L}(k)}} + p_{nc} \cdot e^{\mathcal{L}(k)} \right)$$

where $\mathcal{L}$ represents the input LLR sequence to be updated by the function $f_c(\cdot)$, and $p_{nc}$ is defined as

$$p_{nc} = \frac{\sum_{k=1}^{K} |u_R(k) - u_\oplus(k)|}{K} \quad (9)$$

where the sequence $\mathbf{u}_\oplus$ is generated by the source nodes specified by $\mathbf{F}$ as

$$\mathbf{u}_\oplus = \bigoplus_{j \in \mathcal{F}} \tilde{u}_j. \quad (10)$$

The knowledge of the bit-flipping probability $p_{nc}$ is exploited during the decoding process of JANCC in the LLR updating function $f_c(\cdot)$ to avoid the error propagation if one/some of the sequences $\tilde{u}_i$ obtained at $R$ is/are erroneous. It should be emphasized that the probability $p_{nc}$ can be estimated only at $D$ without having to involve any higher layer protocols, as [24, eq. (9)]

$$\hat{p}_{nc} = G(L_e(\Pi_R[\mathbf{u}_\oplus]), L_e(\Pi_R[\mathbf{u}_R]), \mathcal{G}) = G(L_e(\mathbf{u}_\oplus), L_e(\mathbf{u}_R), \mathcal{G})$$

$$= \frac{1}{|\mathcal{G}|} \sum_{g \in \mathcal{G}} \Pr(u_\oplus(g) = 1) \Pr(u_R(g) = 0) + \Pr(u_\oplus(g) = 0) \Pr(u_R(g) = 1)$$

$$= \frac{1}{|\mathcal{G}|} \sum_{g \in \mathcal{G}} \frac{e^{L_e(u_\oplus(g))}}{1 + e^{L_e(u_\oplus(g))}} + \frac{e^{L_e(u_R(g))}}{1 + e^{L_e(u_R(g))}}$$

$$= \frac{1}{|\mathcal{G}|} \sum_{g \in \mathcal{G}} \frac{e^{L_e(u_\oplus(g))} + e^{L_e(u_R(g))}}{1 + e^{L_e(u_\oplus(g))}} \quad (11)$$

where $\hat{p}_{nc}$ is the estimated value of the probability $p_{nc}$, and $\Pi_R[\cdot]$ denotes interleaving by $\Pi_R$. The set $\mathcal{G}$ is defined as

$$\mathcal{G} = \{g : |L_e(u_\oplus(g))| \leq T\} \quad (12)$$

with $T$ representing a threshold value. However, since the $SD$ links suffer from block fading, we modified [24, eq. (9)] as follows: the value of $T$ decreases during every iterative operation of the loop until a sufficiently large number of $L_e(u_\oplus(g))$ in the sequence $u_\oplus$ is guaranteed (three-quarter of the entire sequence-length at least in this paper). Algorithm 1 summarizes the decoding process of JANCC, where $\sum_{\oplus}$ and $\Pi_R[\cdot]$ represent the boxplus operation [25] and de-interleaving by $\Pi_R$.

**IV. NUMERICAL RESULTS**

In this section we provide the results of the simulations conducted to evaluate the performance of the developed energy-efficient orthogonal MARC system combined with the proposed JANCC techniques, which is referred to as DDEX-JANCC. In the simulations, the length of the information sequence $K$ is 1024 bits, and the number of source nodes and retransmission are three and one, respectively. The RSC encoder $C_S$, $C_R$ and ACC with the generator polynomials $(G_r, G)$ are $([7, 5], [7, 5])$ and $([3, 3], [3, 3])$, respectively, where the generator polynomial is expressed in an octal form and $G_r$ specifies the feedback polynomial. The code rate $K/M$ of $C_S$ and $C_R$ is 1/2, and the Log-Map algorithm is used in all the decoding process. The iterations $I$, interval $\epsilon$, ratio $\rho$, threshold $T^{(0)}$ and times $n$ used in Algorithm 1 are set to 10, 0.05, 3/4, 3 and 10, respectively.

Two other alternative orthogonal MARC systems are also provided in the simulations for comparisons. One system uses the DDEX detection strategy while the other performs the
Fig. 2. Proposed decoding scheme of JANCC represented by factor graph, where ACC$^{-1}$ and SCCC$^{-1}$ denote the ACC decoder and the decoder of the SCCC applied by $S_1$, respectively.

Algorithm 1: Decoding process of JANCC

Pre-defined: iterations $I$, interval $\epsilon$, ratio $\rho$, threshold $T(0)$, times $n$

Input: $L_e(u_j)$, $\forall j \in F$

Initialization: $\hat{p}_{ne} = 0.25$

for $i$ to $I$ do

$\forall j \in F$ do

$\tilde{u}_j$ = $\left\{\begin{array}{ll} u_j & \text{if } |\mathcal{Q}| < \rho \cdot \#L_e(u_{\circ}) \\
\text{Exit for;}
end\end{array}\right.$

$T(l) = T(l-1) - \epsilon$;

if $T(l) \leq 0$ then

Store all the indexes of $L_e(u_{\circ})$ to set $\mathcal{Q}$;

Exit for;

end

$L_\alpha(\Pi_R[u_R]) = f_c(L_e(\Pi_R[u_R]), \hat{p}_{ne})$;

Calculate $L_e(\Pi_R[u_R])$ by $C_R^{-1}$;

$\hat{p}_{ne} = G(L_e(\Pi_R[u_R]), L_e(\Pi_R[u_R]), \mathcal{Q})$;

$\hat{p}_{ne} = G(L_e(\Pi_R[u_R]), L_e(\Pi_R[u_R]), \mathcal{Q})$;

for all $j \in F$ do

Calculate $L_a(u_j)$ by boxplus operation;

Feed $L_a(u_j)$ to SCCC$^{-1}$;

Calculate $L_e(u_j)$ by SCCC$^{-1}$ with $I$ iterations;

end

end

Calculate $\tilde{u}_j$, $\forall j \in F$;

iterative decoding of the SCCC at $R$ to obtain the sequences $\tilde{u}_i$. The JNCC technique [11] based on the conventional DF strategy, referred to as DFJNCC$^4$, is applied by both systems after obtaining $\tilde{u}_i$. According to their detection strategy applied at $R$, we refer these two systems to as DDEX-DFJNCC and SCCC-DFJNCC, respectively. The outage probability for the orthogonal MARC systems with the conventional DF strategy is numerically calculated by Monte Carlo methods, where we generated an enough number of random fading coefficients $(h_{iR}, h_{iD}$ and $h_{RD})$ and measured the probability that the conditions provided in [11, eqs. (4)-(6)] are not satisfied.

The structure for all the orthogonal MARC systems are assumed to be symmetric in geometric gain where all the $SR$ links have the same average SNR ($\Gamma_iR = \text{constant } \forall i$) and all the $SD$ links as well ($\Gamma_iD = \text{constant } \forall i$). Finally, we define FER as follows in this paper: number of the information sequences which cannot be correctly decoded at $D$ even with the help of $R$, divided by the total number of the sequences transmitted by all the source nodes.

Fig. 3 demonstrates the FER performance when $\Gamma_iR = \Gamma_iD + 0.5 \, \text{dB}$. First of all, we investigate the effect of the detection technique applied at $R$, assuming the conventional DF strategy. For the DDEX-DFJNCC system, as can be observed in Fig. 3, there is roughly a 3 dB loss in average $SD$ SNR at FER $= 10^{-2}$ compared with the SCCC-DFJNCC system. The loss is due to the fact that DDEX only extracts the estimates of the information sequences output from DD at $R$. However, if DDEX is combined with the proposed JANCC and its corresponding decoding techniques, i.e., DDEX-JANCC, the loss can be reduced to only 1.5 dB by utilizing the erroneous estimates of the transmitted source information received at $R$.

Furthermore, it can also be observed in Fig. 3 that by employing the proposed JANCC decoding algorithm

$^4$For fair comparisons, the encoder of [11, Fig. 2] is modified by replacing the Turbo Encoder with $C_R$. In addition, the interleaver $\Pi_R$ is added and placed between the XOR-operation and $C_R$. }
DDEX strategy. The DDEX-JANCC system can achieve 1 dB SR because the better quality of the sequences can be clearly observed in Fig. 4 that the gap in FER between the accuracy of the sequences $e^+$ and $e^−$ is very close to the $R_D$ achieved due to low computational complexity required.

A. Impact of the SR link quality

In order to consider the impact of the quality of the SR links to the FER performance of the DDEX-JANCC system, another scenario was tested where $\Gamma_{IR} = \Gamma_{ID} + 6.5$ dB. It can be clearly observed in Fig. 4 that the gap in FER between the DDEX-DFJNCC and SCCC-DFJNCC systems is reduced, because the better quality of the SR links achieves the higher accuracy of the sequences $\hat{u}_i$, obtained at $R$, even by using the DDEX strategy. The DDEX-JANCC system can achieve 1 dB gain over DDEX-DFJNCC at FER = $10^{-2}$, and only 0.5 dB away from SCCC-DFJNCC. However, the FER performance of the DDEX-JANCC system is roughly 3.5 dB away from the DF outage probability [11]. This observation indicatively means that the optimal code design for the $SR$ links is left as a future study, when the quality of the $SR$ links is good.

B. Computational Complexity Evaluation

We also make a comparison between the DDEX-JANCC and SCCC-DFJNCC systems in terms of computational complexity. The main difference between the two systems is due mainly to the detection strategy applied at $R$, so we calculate the number of addition (ADD), multiplication (MUL) and comparison (COMP) operations needed to obtain one estimated information sequence $\hat{u}_i$, where all the operations are on real numbers.

It is assumed $\hat{u}_i$ is obtained by performing the iterative decoding of the SCCC on the received vector $y_{IR}$, where the Log-MAP algorithm is used on the inner and outer codes with memory length $m_1$ and $m_2$, respectively. The number of operations needed per bit in the Log-MAP decoding algorithm is provided in [23, Sec. 4.1]. Since the lengths of sequences decoded by the inner and outer decoders are $M$ and $K$ bits, respectively, the required number of operations per sequence of each constituent code is approximately $M$ and $K$ times higher, respectively. The total computational cost for the iterative decoding of the SCCC per iteration is obtained by summing the number of operations in the decoding of both codes, and thus, $(15 \cdot 2^{m_1} + 9) \cdot M + (15 \cdot 2^{m_2} + 9) \cdot K$ ADDs, $8 \cdot M + 8 \cdot K$ MULs and $(5 \cdot 2^{m_1} - 2) \cdot M + (5 \cdot 2^{m_2} - 2) \cdot K$ COMPs in total are required.

However, when using the DDEX strategy, only one complex multiplication per bit, equivalent to two ADDs and four MULs, is required by the DD strategy, and since $y_{IR}$ is an $M$-bit sequence, $2M$ ADDs and $4M$ MULs in total are required in the detection of $y_{IR}$. The sequence $\hat{u}_i$ is simply extracted from the output of DD and thus no operations are required.

It is clear that the required number of operations in each category in SCCC decoding is much higher than that in DDEX, where the cost for the correction function used in the Log-MAP algorithm is not considered. In addition, assuming ADD, MUL and COMP have equal complexity, with $m_1 = 1$, $m_2 = 2$, $K/M = 1/2$ and 10 iterations set in this paper, the complexity of SCCC decoding is at least 170 times higher than DDEX, which leads to significant power saving for computation at $R$.

C. Average throughput

For the proposed JANCC strategy, unlike DFJNCC, the specified estimates of the information sequences $\hat{u}_j$, $\forall j \in F$, are, regardless of whether they are correct or not, always jointly network-channel coded at $R$ and forwarded to $D$ when $R$ received the retransmission request from $D$. Hence, in order to identify the probability for recovering the information sequence correctly at $D$ with applying DFJNCC and JANCC at $R$, we also evaluate the average throughput of all the orthogonal MARC systems used in the simulations.
TABLE I
ACHIEVED THROUGHPUT WITH/ WITHOUT RETRANSMISSION.

<table>
<thead>
<tr>
<th>$D$ sends request</th>
<th>DFJNCC</th>
<th>JANCC</th>
</tr>
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<tbody>
<tr>
<td>$R$ forwarding</td>
<td></td>
<td></td>
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<tr>
<td>No</td>
<td></td>
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<tr>
<td>Yes</td>
<td></td>
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</tr>
<tr>
<td>$\text{correctly decoded msgs at } D$</td>
<td>$N$</td>
<td>$N-1$</td>
</tr>
<tr>
<td></td>
<td></td>
<td>$N-1$</td>
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<td></td>
<td></td>
<td>$1$</td>
</tr>
<tr>
<td></td>
<td></td>
<td>$0$</td>
</tr>
<tr>
<td>Achieved throughput</td>
<td>$\frac{N}{N+1}$</td>
<td>$\frac{N}{N+1}$</td>
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<tr>
<td></td>
<td></td>
<td>$N$</td>
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<td>$\frac{N}{N+1}$</td>
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<td>$1$</td>
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<td></td>
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<td>$0$</td>
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</table>

Fig. 5. Average throughput when $\Gamma_{iR} = \Gamma_{iD} + 0.5$ dB, $\Gamma_{RD} = 20$ dB, and the number of source nodes $N = 3$.

Fig. 6. Average throughput when $\Gamma_{iR} = \Gamma_{iD} + 6.5$ dB, $\Gamma_{RD} = 20$ dB, and the number of source nodes $N = 3$.

systems for $\Gamma_{iR} = \Gamma_{iD} + 0.5$ dB. The average throughput of the DDEX-JANCC system is even lower than the point-to-point transmission in the low $SD$ SNR regime, which indicates the help of the proposed JANCC technique for the recovery of the information sequence at $D$ is limited when the $SR$ link quality is poor. The probability of $\tilde{u}_j$, $\forall j \in \mathcal{F}$ after the detection of DDEX containing a large amount of errors is relatively high when the quality of the $SR$ links is poor, and thus, the bit-flipping probability between $u_D$ and $u_R$ increases if any one of the sequences $u_j$ contains a large amount of errors, which decreases the correction capability of JANCC decoding.

As the quality of the $SR$ links improves, as shown in Fig. 5 and Fig. 6, the performance of the DDEX-JANCC system is apparently improved and outperforms the performance of DDEX-DFJNCC in terms of the average throughput efficiency. Furthermore, the average throughput of the DDEX-JANCC system asymptotically approaches to that with SCCC-DFJNCC, as shown in Fig. 6. Therefore, it can be concluded that the average throughput efficiency of the DDEX-JANCC system depends on the quality of the $SR$ links.

V. CONCLUSION

This paper has investigated the developed energy-efficient orthogonal MARC system. A low-complexity and non-coherent detection strategy DDEX was applied at the relay $R$ to reduce the energy consumption. The proposed JANCC and its corresponding decoding techniques were employed to utilize the erroneous estimates output from DDEX. A vector is constructed in the JANCC technique to identify the indexes of the source nodes whose information sequences are decoded in error at the destination $D$. The relay $R$ performs network coding only on the estimates specified by the vector when receiving the retransmission request, which aims to avoid unnecessary erroneous estimates being network coded. The bit-flipping probability $p_{nc}$ is observed between the two sequences, one is the network-coded sequence sent from $R$, and the other is their corresponding XOR-ed information sequence obtained at $D$. The proposed JANCC decoding algorithm estimated and exploited the probability $p_{nc}$ to update the Log-Likelihood Ratio during the iterative decoding process. In the case of the system setup assumed in this paper, computational complexity of the DDEX strategy is at least 170 times lower than the iterative decoding of the SCCC performed at $R$, in exchange for only a $0.5 - 1.5$ dB loss in the FER performance. However, a significant reduction in power consumption can be achieved due to low computational complexity required and the elimination of channel estimation, which is suitable for summarizes all the combinations for the cases that require retransmission request for applying DFJNCC and JANCC at $R$. The achieved throughput is defined as the number of correctly decoded information sequences at $D$ totalling over all the source nodes, and the average throughput is calculated by averaging the achieved throughput under all the combinations according to the corresponding occurrence probabilities of those cases.

Fig. 5 shows the average throughput efficiencies of all the
the relaying systems where the availability of the resources is limited.

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Pen-Shun Lu (S’09) received the B.Sc. in electrical engineering from National Sun Yat-Sen University (NSYSU), Taiwan, in 2003 and the M.Sc in wireless communications from the University of Southampton, UK, in 2007. Currently he is pursuing the PhD degrees in University of Oulu, Finland and Japan Advanced Institute of Science and Technology (JAIST), Japan. His current research interests include joint source and channel coding, distributed coding, iterative detection and cooperative communications.

Xiaobo Zhou (S’11–M’13) received the B.Sc. in Electronic Information Science and Technology from University of Science and Technology of China (USTC), Hefei, China, in 2007, the M.E. in Computer Application Technology from Graduate University of Chinese Academy of Science (GU-CAS), Beijing, China, in 2010, and the Ph.D. degree from School of Information Science, Japan Advanced Institute of Science and Technology (JAIST), Ishikawa, Japan, in 2013. Since October 2013, he is with the School of Information Science, JAIST as a researcher. His research interests include coding techniques, joint source-channel coding, cooperative wireless communications and network information theory.

Khoirul Anwar (S’03–M’08) graduated (cum laude) from the department of Electrical Engineering (Telecommunications), Institut Teknologi Bandung (ITB), Bandung, Indonesia in 2000. He received Master and Doctor Degrees from Graduate School of Information Science, Nara Institute of Science and Technology (NAIST) in 2005 and 2008, respectively. Since then, he has been appointed as an assistant professor in NAIST. He received best student paper award from the IEEE Radio and Wireless Symposium 2006 (RWS ’06), California-USA, Best Paper of Indonesian Student Association (ISA ’07), Kyoto, Japan in 2007, and Award for Innovation, Congress of Indonesian Diaspora (CID), Los Angeles, USA, July 2012. Since September 2008, he is with the School of Information Science, Japan Advanced Institute of Science and Technology (JAIST) as an assistant professor. His research interests are network information theory, error control coding, iterative decoding and signal processing for wireless communications. Dr. Anwar is a member of IEEE, information theory society, communication society and IEICE Japan.
Tad Matsumoto (S’84—SM’95—F’10) received his B.S., M.S., and Ph.D. degrees from Keio University, Yokohama, Japan, in 1978, 1980, and 1991, respectively, all in electrical engineering. He joined Nippon Telegraph and Telephone Corporation (NTT) in April 1980. Since he engaged in NTT, he was involved in a lot of research and development projects, all for mobile wireless communications systems. In July 1992, he transferred to NTT DoCoMo, where he researched Code-Division Multiple-Access techniques for Mobile Communication Systems. In April 1994, he transferred to NTT America, where he served as a Senior Technical Advisor of a joint project between NTT and NEXTEL Communications. In March 1996, he returned to NTT DoCoMo, where he served as a Head of the Radio Signal Processing Laboratory. In March 2002, he moved to University of Oulu, Finland, where he served as a Professor at Centre for Wireless Communications. In 2006, he served as a Visiting Professor at Ilmenau University of Technology, Ilmenau, Germany, funded by the German MERCATOR Visiting Professorship Program. In April 2007, he returned to Japan and since then he has been serving as a Professor at Japan Advanced Institute of Science and Technology (JAIST), while also keeping a part-time position at University of Oulu. Prof. Matsumoto has been appointed as a Finland Distinguished Professor for a period from January 2008 through December 2012, funded by the Finnish National Technology Agency (Tekes) and Finnish Academy, under which he preserves the rights to participate in and apply to European and Finnish national projects. Prof. Matsumoto is a recipient of IEEE VTS Outstanding Service Award (2001), Nokia Foundation Visiting Fellow Scholarship Award (2002), IEEE Japan Council Award for Distinguished Service to the Society (2006), IEEE Vehicular Technology Society James R. Evans Avant Garde Award (2006), and Thuringen State Research Award for Advanced Applied Science (2006), 2007 Best Paper Award of Institute of Electrical, Communication, and Information Engineers (IEICE) of Japan (2008), Telecom System Technology Award by the Telecommunications Advancement Foundation (2009), IEEE Communication Letters Exemplifying Reviewer Award (2011), UK Royal Academy of Engineering Distinguished Visiting Fellow Award (2012) and Nikkei Electronic Wireless Japan Awards (2013). He is serving as an IEEE Vehicular Technology Distinguished Lecturer since July 2011. He is a Fellow of IEEE and a member of IEICE.