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Author(s)	Qian, Shen; Zhou, Xiaobo; He, Xin; He, Jiguang; Juntti, Markku; Matsumoto, Tad
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Description	

Performance Analysis for Lossy-Forward Relaying over Nakagami- m Fading Channels

Shen Qian, *Student Member, IEEE*, Xiaobo Zhou, *Member, IEEE*, Xin He, *Member, IEEE*, Jiguang He, *Student Member, IEEE*, Markku Juntti, *Senior Member, IEEE*, and Tad Matsumoto, *Fellow, IEEE*

Abstract—We investigate the performance of three-node lossy-forward (LF) relaying over independent block Nakagami- m fading channels. Based on the theorem of source coding with side information, the exact outage probability expression for arbitrary values of the shape factor m is derived under the assumptions of both the Gaussian codebook capacity and the constellation constrained capacity. The difference in outage probability between the two codebook models of capacity is found to be very minor. Furthermore, an accurate high signal-to-noise ratio approximation for the outage probability is obtained. It clearly identifies the equivalent diversity order and coding gain of the LF relaying. It is shown that the LF relaying is superior to conventional decode-and-forward relaying in terms of the outage probability and the ϵ -outage achievable rate. Moreover, with the LF relaying, the optimal location for the relay (R), which minimizes the outage probability, is found to be the point having same distance to the source (S) and the destination (D) when the S-R and R-D links experience same level of fading. The accuracy of the analytical results is verified by a series of Monte Carlo simulations.

Index Terms—outage probability, relay channels, lossy forward, Nakagami- m fading, equivalent diversity order, coding gain, ϵ -outage achievable rate, optimal relay location

I. INTRODUCTION

COOPERATIVE communication is a scheme which offers spatial diversity to mitigate the effect of fading in the propagation medium [1], [2]. Among the various relaying protocols, decode-and-forward (DF) relaying has drawn significant and practical attentions, and been widely studied. Several analyses about the diversity and multiplexing gains as well as the diversity-multiplexing tradeoff of the DF relaying have been conducted for half duplex and full duplex relaying [3]–[5]. Various practical implementations of the

DF relaying have been performed for single-antenna [6], [7] and multiple-antenna scenarios [8], [9]. Theoretical outage probability, which provides useful performance metrics for designing practical coding techniques [10], is analyzed for the DF relaying in [11]–[13].

Recently, a new cooperative communication scheme based on the DF protocol allowing source-relay (S-R) link errors is proposed in [14]. The scheme is referred to as lossy forward (LF) relaying in this paper. Unlike the conventional DF strategy, the decoded information sequence at relay (R) is interleaved, re-encoded, and transmitted to destination (D), even though error is detected during the decoding process at R. The two information sequences, one from the source (S) and the other from R, received at D are correlated. The correlation knowledge can be estimated and utilized in a joint decoder at D. For example, iterative processing between two decoders (one is for S and the other is for R) for updating log-likelihood ratio (LLR) via a LLR modification function, proposed in [15], reduces the decoding error probability. The basic idea behind the LF relaying is the relay system, as a whole, can be seen as a distributed turbo code, and hence it can achieve turbo-cliff-like bit-error-rate (BER) performance in additive white Gaussian noise (AWGN) channels [16].

The outage probability of the LF relaying has been evaluated in [17], where the admissible rate region is determined by the Slepian-Wolf correlated source coding theorem. It was found soon after [17] was published, that analyzing the exact rate region of the LF scheme falls into the source coding with side information in network information theory [18], [19]. Based on this finding, the exact and approximated outage probabilities were analyzed in [20]. It was shown that the LF relaying can achieve lower outage probability in Rayleigh fading channels compared to the DF relaying. The technique is further extended to wireless sensor networks (WSNs) and a simple, general, efficient power allocation scheme for arbitrary number of sensors is derived [21]. In [22], the outage probability of the LF relaying is analyzed, where the source-destination (S-D) link is assumed to be suffering from block Rayleigh fading, whereas the relay-destination (R-D) link from block Rician fading. Brulatout *et al.* propose a multiple access channel (MAC) protocol in [23] specifically designed to optimally operate with the LF paradigm.

However, there are several drawbacks in the previous work related to the LF system: 1) Performance analyses for the LF relaying are all based on numerical calculation. Diversity and coding gain cannot be identified separately from the outage probability expression; 2) In [17], [20], [22], Rayleigh and

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Shen Qian and Tad Matsumoto are with the School of Information Science, Japan Advanced Institute of Science and Technology (JAIST), 1-1 Asahidai, Nomi, Ishikawa, Japan 923-1292 (e-mail: {shen.qian, matsumoto}@jaist.ac.jp), and the Centre for Wireless Communication, University of Oulu, FI-90014 Finland.

Xiaobo Zhou is with the School of Computer Science and Technology, Tianjin University, Tianjin, China. (e-mail: xiaobo.zhou@tju.edu.cn).

Xin He is with the School of Mathematics and Computer Science, Anhui Normal University, Anhui, China. (e-mail: xin.he@ahnu.edu.cn).

Jiguang He and Markku Juntti are with the Centre for Wireless Communication, University of Oulu, FI-90014 Finland (e-mail: {jhe, markku.juntti}@oulu.fi).

Rician distributions are assumed as fading channel models, since they are widely used to describe the statistical properties of the fading channels in wireless mobile communication systems. However, Nakagami- m distribution is considered to be more practical and better matched to field measurement data than Rayleigh and Rician distributions [24]; 3) Even though [20] has derived the exact outage probability of the LF relaying based on the rate region determined by the theorem of source coding with side information, Gaussian codebook capacity (GCC) is assumed when describe the relationship between the Hamming distortion of the S-R link and the receive block-wise instantaneous signal-to-noise ratio (SNR). However, in practice, channel capacity is upper bounded by the modulation constellation size.

The main objective of this paper is to eliminate the aforementioned drawbacks and to theoretically investigate outage probability of the LF relaying, where all the links are assumed to experience block Nakagami- m fading. The major contributions of this paper are as follows.

- The exact outage probability for the LF relaying over Nakagami- m fading channels is derived. The outage expressions are further approached by accurate enough approximations in high SNR regime. Explicit expressions of the equivalent diversity order and coding gain for arbitrary values of the shape factor m are then obtained from the approximated outage expressions. It is shown that the equivalent diversity order (outage curve decay) of the LF relaying is limited to either the S-R link or R-D link, having smaller shape factor m .
- The outage probabilities of the LF system using the GCC and constellation constrained capacity (CCC) are derived. The numerical results indicate that the difference between the GCC and CCC based outage probabilities of the LF relaying is negligible.
- We introduce another performance metric, ϵ -outage achievable rate, which indicates the largest transmission rate such that the outage probability is not larger than a threshold value ϵ [25], [26]. We identify from the analytical results that in high SNR region, the ϵ -outage achievable rate of the LF relaying increases rapider than that of the conventional DF relaying.
- The optimal relay location is investigated for minimizing the outage probability of the LF relaying. The relay R is considered to be located on a line parallel to the S-D link, as shown in Fig. 1. It is shown that the optimal relay location is always at the point where the S-R and R-D links have equal distance, as long as the values of shape factor m of the S-R and R-D links are equivalent.

The rest of this paper is organized as follows: the system model used to analyze the outage probability of the LF relaying is presented in Section II. The admissible rate region and outage probability definition are described in Section III. The exact outage probabilities of the LF relaying based on GCC and CCC are also derived in Section III. Equivalent diversity order and coding gain with the LF relaying are provided in Section IV. In Section V, ϵ -outage achievable rate is analyzed. The optimal relay location is investigated

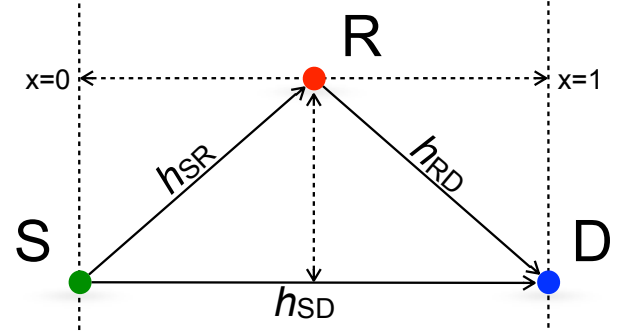


Fig. 1. The schematic of one relay aided communication system.

in Section VI. Finally, the conclusion is drawn in Section VII.

II. SYSTEM MODEL

We consider a simple three-node relaying system as shown in Fig. 1. The source S communicates with the destination D with the help of a single relay R. The location of R varies in a line parallel to the S-D link between $x = 0$ (nearest to S) and $x = 1$ (nearest to D). We assume time-division transmission, where the overall transmission is divided into two time slots. At the first time slot, the original uniformly distributed binary information sequence b_S is broadcast from S. The relay R aims to recover the information sequence and transmits it to D at the second time slot, i.e., in orthogonal transmission. S does not transmit during the second time slot.

A. LF Relaying

In conventional DF relaying, R keeps silent if error/errors is/are detected after decoding in the information sequence sent through the S-R link. With the LF relaying, after receiving the signal from S, R attempts to recover b_S . Although the decoding result of b_S at R, denoted by b_R , may be found to contain errors, R interleaves the information sequence b_R , re-encodes the interleaved sequence, and forwards it to D.

When only the information sequence is concerned before encoding at S and after decoding at R, the S-R link can be virtually modeled by a binary symmetric channel (BSC) model with a crossover probability p_f . More specifically, p_f represents the bit flipping probability between the information sequence obtained after decoding at R and the original information sequence sent from S. Hence, $b_R = b_S \oplus e$, where $\oplus$ denotes the modulo-2 addition and e is a realization of a binary random variable E with $\Pr(E = 1) = p_f$ which stays the same within each block while changes block-by-block with its value determined by the instantaneous SNR.

At D, after receiving the signals from S and R, joint decoding is performed to retrieve the original information b_S . Iterative decoding is utilized between two decoders for decoding the messages sent from S and R. During the decoding process, the S-R link error probabilities can be estimated at D by making a comparison between the LLRs output from the two decoders. The estimated p_f value is used as the correlation knowledge between b_S and b_R [14]. The LLRs of

the systematic bits are exchanged between the two decoders via the interleaver/de-interleaver. Therefore, the system, as a whole, can be viewed as a distributed turbo code.

B. Channel Model

The signals received at D and R during the first time slot, $y_{D,1}$ and $y_{R,1}$, respectively, and the signal received at D in the second time slot, $y_{D,2}$, are expressed as¹

$$y_{D,1} = \sqrt{G_{SD}}h_{SD}x_1 + n_{D,1}, \quad (1)$$

$$y_{R,1} = \sqrt{G_{SR}}h_{SR}x_1 + n_{R,1}, \quad (2)$$

$$y_{D,2} = \sqrt{G_{RD}}h_{RD}x_2 + n_{D,2}, \quad (3)$$

respectively, where G_{ij} ($i \in (S, R), j \in (R, D), i \neq j$) are the geometric gains related to the transmit distance of each link. x_1 and x_2 denote the modulated symbols corresponding to the coded and interleaved information sequence, transmitted from S and R, respectively. h_{ij} denotes the complex channel gain and $n_{j,1}$ and $n_{j,2}$ are zero-mean AWGN with variance of $N_0/2$ per dimension. It is assumed that $E[|h_{ij}|^2] = 1$ and h_{ij} stays constant over one block duration due to the block fading assumption.

The transmit power of each symbol is denoted as E_s . Therefore, the average and instantaneous SNRs are expressed as $\bar{\gamma}_{ij} = G_{ij} \frac{E_s}{N_0}$ and $\gamma_{ij} = |h_{ij}|^2 \bar{\gamma}_{ij}$ ($i \in (S, R), j \in (R, D), i \neq j$), respectively. For the sake of simplicity, the variations due to shadowing and the fading frequency selectivity are not taken into account.

We assume that all the links (i.e., the S-R, S-D, and R-D links) suffer from independent block Nakagami- m fading, with the probability density function (pdf) of γ_{ij} given by

$$p(\gamma_{ij}) = \frac{m_{ij}^{m_{ij}} (\gamma_{ij})^{m_{ij}-1}}{(\bar{\gamma}_{ij})^{m_{ij}} \Gamma(m_{ij})} \exp\left(-\frac{m_{ij}\gamma_{ij}}{\bar{\gamma}_{ij}}\right), m_{ij} > 0.5, \quad (4)$$

where $\Gamma(\cdot)$ is the Gamma function. The shape factor m_{ij} represents the severity of the fading variation of the corresponding link.

In cooperative networks such as vehicle-to-vehicle (V2V) and/or vehicle-to-infrastructure (V2I) communications, the network topology and the fading distributions of the links change over time. A reasonable scenario assumption is that S-D link suffers from severe fading, the destination needs the help of a relay via S-R and R-D links, which suffer from moderate fading. Hence, we set the shape factor of the S-D link $m_{SD} = 1$, corresponding to Rayleigh fading, while for other links (the S-R and R-D links), we set their corresponding m values arbitrarily as parameters.

III. OUTAGE PROBABILITY ANALYSIS

In this section, the definition of admissible rate region and the derivation for outage probability of the LF relaying are provided. The outage probabilities are derived based on both GCC and CCC assumptions.

¹The symbol indexes are omitted in (1), (2), and (3) for conciseness.

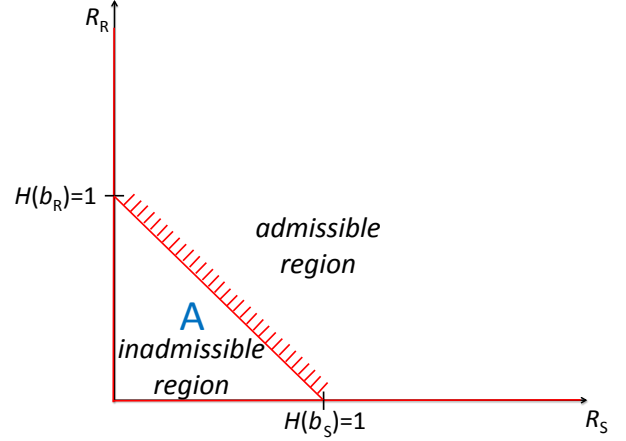


Fig. 2. Rate region for S and R when $p_f = 0$; the red solid line with bars separates the admissible and inadmissible regions.

A. Source Coding with Side Information based Admissible Rate Region for S and R

Since D aims to recover b_S only, the coded signal transmitted from R can be considered as side information of b_S . According to the theorem of source coding with side information [19, Section 10.4], D can recover b_S losslessly, if the source coding rate pair of S and R, R_S and R_R , satisfies

$$\begin{cases} R_S & \geq H(b_S|\hat{b}_R), \\ R_R & \geq I(b_R;\hat{b}_R), \end{cases} \quad (5)$$

where $H(\cdot|\cdot)$ and $I(\cdot;\cdot)$ denote the conditional entropy and the mutual information between their arguments, respectively, and $\hat{b}_R$ is the estimate of b_R obtained at D.

With the block fading assumption, we use a BSC model to represent the R-D link (helper channel), block by block, as $\hat{b}_R = b_R \oplus e'$ with $\Pr(E' = 1) = p'_f$, where e' is a realization of a binary random variable E' . Since the source is assumed to be binary, uniform, and independently and identically distributed (i.i.d), (5) can be expressed as

$$\begin{cases} R_S & \geq H(p_f * p'_f), \\ R_R & \geq H(b_R) - H(\hat{b}_R|b_R) = 1 - H(p'_f), \end{cases} \quad (6)$$

where $p_f * p'_f = (1 - p_f)p'_f + (1 - p'_f)p_f$ and $H(\cdot)$ denotes the binary entropy function.

$p_f = 0$ indicates perfect decoding at R, and hence $H(b_S|b_R) = H(b_R|b_S) = 0$. In this case, the inadmissible rate region becomes the triangle area A as shown in Fig. 2. When $0 < p_f \leq 0.5$, the inadmissible region, which can be divided into two areas, B and C, is shown in Fig. 3. The rate region defined in (6) indicates that, even with $0 \leq R_R \leq H(b_R)$, b_R can be partially recovered at D, and b_R containing errors can serve as the side information for losslessly recovering b_S . In the case $R_R > H(b_R)$, the conditions in (6) become to $R_S \geq H(p_f)$ and $R_R \geq 1$.

B. Outage Event of LF relaying

If the rate pair (R_S, R_R) falls into the inadmissible regions in Fig. 2 or Fig. 3, the outage event occurs and D cannot

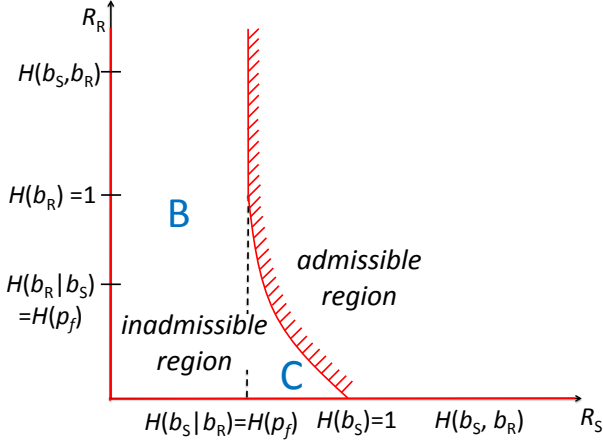


Fig. 3. Rate region for S and R when $p_f \neq 0$; the red solid line with bars separates the admissible and inadmissible regions.

guarantee the reconstruction of b_S with an arbitrarily small error probability. Since $p_f = 0$ and $0 < p_f \leq 0.5$ are distinctive, the outage probability of the LF relaying can be expressed as

$$P_{\text{out}}^{\text{LF}} = P_A + P_B + P_C, \quad (7)$$

where P_A , P_B , and P_C denote the probabilities that (R_S, R_R) falls into the inadmissible areas A, B, and C, respectively. Taking into account the impact of p_f and p'_f , P_A , P_B , and P_C can further be expressed as

$$P_A = \Pr[p_f = 0, 0 \leq R_S < 1, 0 \leq R_R < H(p_f * p'_f)], \quad (8)$$

$$P_B = \Pr[0 < p_f \leq 0.5, 0 \leq R_S < H(p_f), R_R \geq 0], \quad (9)$$

$$P_C = \Pr[0 < p_f \leq 0.5, H(p_f) \leq R_S < 1, 0 \leq R_R < H(p_f * p'_f)]. \quad (10)$$

1) *Outage Derivation with Gaussian Codebook Capacity (GCC)*: For calculating the outage probability, first we establish the relationship between γ_{SD} and R_S , and that between γ_{RD} and R_R . According to the Shannon's lossless source channel separation theorem [27], if

$$R_i \cdot R_{c,iD} \leq C_{iD}, \quad (i \in S, R) \quad (11)$$

is satisfied, the error probability can be arbitrarily small at the destination. C_{iD} and $R_{c,iD}$, respectively, denote the channel capacity of the i -D link and the normalized spectrum efficiency of the corresponding transmission chain. The spectrum efficiency includes the channel coding rate and the modulation constellation size. With the assumption that Gaussian codebook is used, the channel capacity C_{iD} of each link can be expressed as

$$C_{iD} = \frac{E^n}{2} \log_2 \left(1 + \frac{2\gamma_{iD}}{E^n} \right), \quad (i \in S, R), \quad (12)$$

where E^n denotes the signaling dimensionality. Hence, the relationship between the instantaneous channel SNR γ_{iD} and

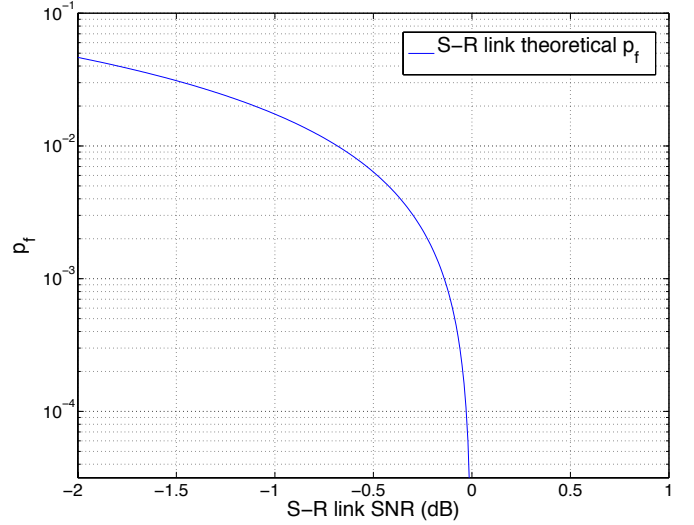


Fig. 4. S-R link Hamming distortion p_f versus SNR, where $E_n = 2$ and $R_{c,SR} = 1$.

its corresponding source coding rate R_i is given by

$$R_i \leq \Theta(\gamma_{iD}) = \frac{E^n}{2R_{c,iD}} \log_2 \left(1 + \frac{2\gamma_{iD}}{E^n} \right), \quad (i \in S, R) \quad (13)$$

with its inverse inequality

$$\gamma_{iD} \geq \Theta^{-1}(R_i) = \frac{E^n}{2} \left(2^{\frac{2R_i R_{c,iD}}{E^n}} - 1 \right). \quad (14)$$

Then, we establish the relationship between γ_{SR} and p_f . Since p_f only depends on the quality of the S-R link, according to Shannon's lossy source channel separation theorem [28], we have

$$R(\mathcal{D}) \cdot R_{c,SR} \leq C_{SR}, \quad (15)$$

where $R(\mathcal{D})$ denotes the source rate-distortion function with the distortion measure $\mathcal{D}$. For the Hamming distortion measure, $R(\mathcal{D}) = 1 - H(\mathcal{D})$. Following [29], we set p_f to the minimum value of $\mathcal{D}$ yielding the specified rate $R(\mathcal{D})$ in (15) for a specific instantaneous S-R link SNR γ_{SR} .

Since we assume that Gaussian codebook is used for the S-R link transmission, the relationship between the required instantaneous channel SNR γ_{SR} and its corresponding source rate $R(\mathcal{D})$ with distortion $\mathcal{D}$ is given by

$$\gamma_{SR} \geq \frac{E^n}{2} \left(2^{\frac{2R(\mathcal{D}) R_{c,SR}}{E^n}} - 1 \right). \quad (16)$$

Then, we can obtain the relationship between p_f and γ_{SR} as

$$p_f = H^{-1} \left(1 - \frac{\frac{E^n}{2} \log_2 \left(1 + \frac{2\gamma_{SR}}{E^n} \right)}{R_{c,SR}} \right), \quad (17)$$

with $H^{-1}(\cdot)$ denoting the inverse function of $H(\cdot)$. The relationship between the S-R link SNR and p_f is shown in Fig. 4.

With the assumption that each link suffers from statistically independent block Nakagami- m fading, each term of the

outage probability expression with the LF relaying in (7) can be expressed as

$$\begin{aligned}
 P_A &= \Pr[\gamma_{SR} > \Theta^{-1}(1), \Theta^{-1}(0) < \gamma_{SD} \leq \Theta^{-1}(1), \\
 &\quad \Theta^{-1}(0) \leq \gamma_{RD} < \Theta^{-1}(1 - \Theta(\gamma_{SD}))] \\
 &= \int_{\gamma_{SR}=\Theta^{-1}(1)}^{\Theta^{-1}(\infty)} \int_{\gamma_{SD}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \int_{\gamma_{RD}=\Theta^{-1}(0)}^{\Theta^{-1}(1-\Theta(\gamma_{SD}))} \\
 &\quad p(\gamma_{SR})p(\gamma_{SD})p(\gamma_{RD})d\gamma_{SR}d\gamma_{SD}d\gamma_{RD} \\
 &= \frac{1}{\bar{\gamma}_{SD}} \left(1 - \left[\frac{\gamma(m_{SR}, m_{SR} \frac{1}{\bar{\gamma}_{SR}})}{\Gamma(m_{SR})} \right] \right) \\
 &\quad \cdot \int_{\gamma_{SD}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \exp\left(-\frac{\gamma_{SD}}{\bar{\gamma}_{SD}}\right) \\
 &\quad \cdot \left[\frac{\gamma(m_{RD}, m_{RD} \frac{\Theta^{-1}(1-\Theta(\gamma_{SD}))}{\bar{\gamma}_{RD}})}{\Gamma(m_{RD})} \right] d\gamma_{SD}, \quad (18)
 \end{aligned}$$

$$\begin{aligned}
 P_B &= \Pr[\Theta^{-1}(0) \leq \gamma_{SR} \leq \Theta^{-1}(1), \Theta^{-1}(0) < \gamma_{SD} \\
 &\quad \leq \Theta^{-1}(1 - \Theta(\gamma_{SR})), \gamma_{RD} \geq \Theta^{-1}(0)] \\
 &= \int_{\gamma_{SR}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \int_{\gamma_{SD}=\Theta^{-1}(0)}^{\Theta^{-1}(1-\Theta(\gamma_{SR}))} \int_{\gamma_{RD}=\Theta^{-1}(0)}^{\Theta^{-1}(\infty)} \\
 &\quad p(\gamma_{SR})p(\gamma_{SD})p(\gamma_{RD})d\gamma_{SR}d\gamma_{SD}d\gamma_{RD} \\
 &= \int_{\gamma_{SR}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \frac{m_{SR}^{m_{SR}} (\gamma_{SR})^{m_{SR}-1}}{(\bar{\gamma}_{SR})^{m_{SR}} \Gamma(m_{SR})} \exp\left(-\frac{m_{SR}\gamma_{SR}}{\bar{\gamma}_{SR}}\right) \\
 &\quad \cdot \left[1 - \exp\left(-\frac{\Theta^{-1}(1-\Theta(\gamma_{SR}))}{\bar{\gamma}_{SD}}\right) \right] d\gamma_{SR} \quad (19)
 \end{aligned}$$

and

$$\begin{aligned}
 P_C &= \Pr\{\Theta^{-1}(0) \leq \gamma_{SR} \leq \Theta^{-1}(1), \Theta^{-1}[1 - \Theta(\gamma_{SR})] \\
 &\quad < \gamma_{SD} \leq \Theta^{-1}(1), \Theta^{-1}(0) \leq \gamma_{RD} \geq \Theta^{-1}[\xi(\gamma_{SD}, \gamma_{SR})]\} \\
 &= \int_{\gamma_{SR}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \int_{\gamma_{SD}=\Theta^{-1}(1-\Theta(\gamma_{SR}))}^{\Theta^{-1}(1)} \int_{\gamma_{RD}=\Theta^{-1}(0)}^{\Theta^{-1}[\xi(\gamma_{SD}, \gamma_{SR})]} \\
 &\quad p(\gamma_{SR})p(\gamma_{SD})p(\gamma_{RD})d\gamma_{SR}d\gamma_{SD}d\gamma_{RD} \\
 &= \frac{1}{\bar{\gamma}_{SD}} \int_{\gamma_{SR}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \int_{\gamma_{SD}=\Theta^{-1}(1-\Theta(\gamma_{SR}))}^{\Theta^{-1}(1)} \\
 &\quad \exp\left(-\frac{\gamma_{SD}}{\bar{\gamma}_{SD}}\right) \frac{m_{SR}^{m_{SR}} (\gamma_{SR})^{m_{SR}-1}}{(\bar{\gamma}_{SR})^{m_{SR}} \Gamma(m_{SR})} \exp\left(-\frac{m_{SR}\gamma_{SR}}{\bar{\gamma}_{SR}}\right) \\
 &\quad \cdot \left[\frac{\gamma(m_{RD}, m_{RD} \frac{\xi(\gamma_{SD}, \gamma_{SR})}{\bar{\gamma}_{RD}})}{\Gamma(m_{RD})} \right] d\gamma_{SD}d\gamma_{SR}, \quad (20)
 \end{aligned}$$

where $\xi(\gamma_{SD}, \gamma_{SR}) = H\{H^{-1}[1 - \Theta(\gamma_{SD})] * H^{-1}[1 - \Theta(\gamma_{SR})]\}$ and $\gamma(\cdot, \cdot)$ is the lower incomplete gamma function.

Fig. 5 presents the outage probabilities of the LF relaying, denoted as P_{out}^{LF} . As a reference, the outage probability of a conventional DF system, denoted as P_{out}^{DF} , is also plotted. It is found that LF achieves lower outage probability than DF. This is because LF always forwards information sequences from the relay to destination, and it converts the system into a distributed turbo code. Note that the theoretical outage

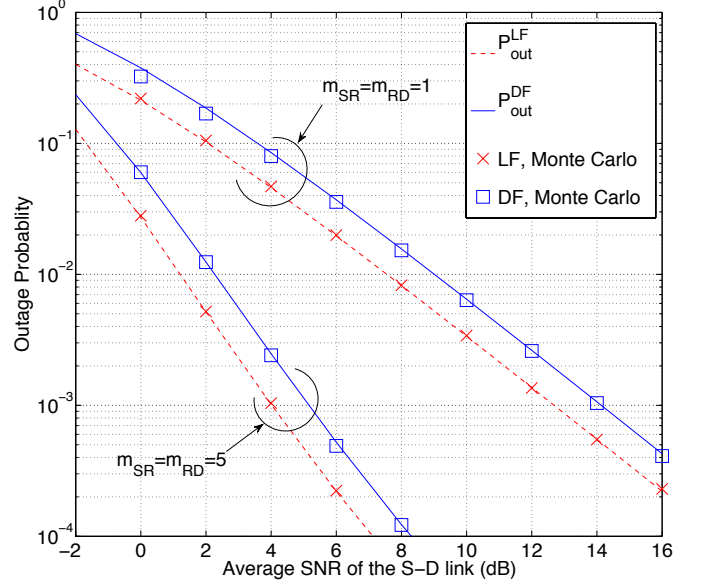


Fig. 5. Outage probability comparison between LF and DF, where $m_{SD} = 1$ and $G_{SD} = G_{RD} = G_{SR}$.

probabilities exactly match their corresponding simulation results obtained by Monte Carlo method.

2) *Outage Derivation with Constellation Constrained Capacity (CCC)*: According to Shannon's source channel separation theorem, if $R_i \cdot R_{c,iD} \leq C_{cc}(\gamma_{iD})$, ($i \in S, R$) is satisfied, the error probability can be made arbitrarily small at the destination, where $C_{cc}(\gamma_{iD})$ is the channel CCC with the instantaneous SNR γ_{iD} of the i -D link. Since the CCC has no closed-form expression, numerical evaluation yields a relationship between γ_{iD} and R_i as $\gamma_{iD} \geq C_{cc}^{-1}(R_i \cdot R_{c,iD})$, where C_{cc}^{-1} represents the inverse function of channel CCC.

Similarly, the relationship between γ_{SR} and p_f is given by $R(\mathcal{D}) \cdot R_{c,SR} \leq C_{cc}(\gamma_{SR})$ and $\gamma_{SR} \geq C_{cc}^{-1}(R(\mathcal{D}) \cdot R_{c,SR})$, where the minimum distortion $\mathcal{D}$ is equivalent to p_f for a given instantaneous S-R link SNR γ_{SR} . Then, by using the method presented in the previous subsection, the outage probability of the LF relaying with CCC can be calculated through a lookup table at each SNR value.

Fig. 6 compares the outage probabilities of the LF relaying derived from GCC and CCC, P_{out}^{GCC} and P_{out}^{CCC} , respectively.² The quaternary CCC is calculated via Monte Carlo simulations to create a lookup table. The Monte Carlo simulation results of GCC based LF relaying are also presented. We can observe from Fig. 6 that the theoretical outage probabilities with the LF relaying, calculated by using numerical integral in (18), (19), and (20), well match their corresponding Monte Carlo simulation results. It is also found that, the difference between P_{out}^{GCC} and P_{out}^{CCC} is negligible. This is because in low SNR regime the difference between GCC and CCC is very small. On the contrary, in high SNR region, even though the difference between GCC and CCC becomes large, p_f approaches zero very quickly, resulting in lossless transmission

²In the numerical calculation, $E^n = 2$ and, without loss of generality, $R_{c,ij} = 1$ ($i \in (S, R), j \in (R, D), i \neq j$).

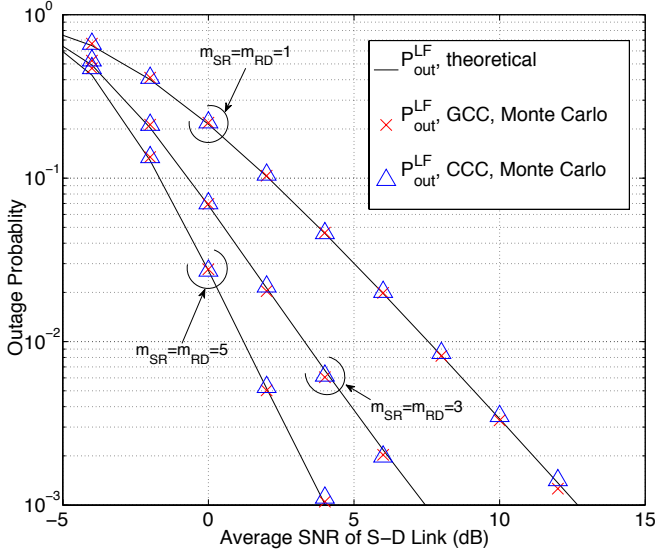


Fig. 6. Comparison of the GCC and CCC based outage probability of the LF relaying, where $m_{SR} = m_{RD}$, $m_{SD} = 1$, and $G_{SD} = G_{RD} = G_{SR}$.

over the S-R link. Therefore, the difference between GCC and CCC based outage probabilities is negligibly small.

IV. EQUIVALENT DIVERSITY ORDER AND CODING GAIN

The equivalent diversity order and coding gain can be obtained by approximating (18), (19), and (20) at high SNR region. Invoking the series representation of incomplete gamma function $\gamma(a, x) = \sum_{n=0}^{\infty} \frac{(-1)^n x^{a+n}}{n!(a+n)}$ [30, equation 8.354.1], together with the approximation $p(\gamma_{ij}) \approx \frac{m_{ij}(\gamma_{ij})^{m_{ij}-1}}{(\bar{\gamma}_{ij})^{m_{ij}} \Gamma(m_{ij})}$ [31], after several steps of mathematical manipulations, each term of the outage probability expression with the LF relaying over Nakagami- m fading channel can be approximated as

$$P_A \approx A \cdot A' \cdot \bar{\gamma}_{SD}^{(-m_{SD})} \bar{\gamma}_{RD}^{(-m_{RD})}, \quad (21)$$

$$P_B \approx B \cdot \bar{\gamma}_{SD}^{(-m_{SD})} \bar{\gamma}_{SR}^{(-m_{SR})}, \quad (22)$$

$$P_C \approx C \cdot \bar{\gamma}_{SD}^{(-m_{SD})} \bar{\gamma}_{RD}^{(-m_{RD})} \bar{\gamma}_{SR}^{(-m_{SR})}, \quad (23)$$

where

$$A = \frac{m_{SD}^{m_{SD}} m_{RD}^{m_{RD}-1}}{\Gamma(m_{SD}) \Gamma(m_{RD})} \int_{\gamma_{SD}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \frac{\left(\frac{2}{1+\gamma_{SD}} - 1\right)^{m_{RD}}}{\gamma_{SD}^{1-m_{SD}}} d\gamma_{SD}, \quad (24)$$

$$B = \frac{m_{SR}^{m_{SR}} m_{SD}^{m_{SD}-1}}{\Gamma(m_{SD}) \Gamma(m_{SR})} \int_{\gamma_{SR}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \frac{\left(\frac{2}{1+\gamma_{SR}} - 1\right)^{m_{SD}}}{\gamma_{SR}^{m_{SR}-1}} d\gamma_{SR}, \quad (25)$$

$$C = \int_{\gamma_{SR}=\Theta^{-1}(0)}^{\Theta^{-1}(1)} \int_{\gamma_{SD}=\Theta^{-1}(1-\Theta(\gamma_{SR}))}^{\Theta^{-1}(1)} \gamma_{SD}^{m_{SD}-1} \gamma_{SR}^{m_{SR}-1} \frac{m_{SR}^{m_{SR}} m_{SD}^{m_{SD}} m_{RD}^{m_{RD}-1}}{\Gamma(m_{SD}) \Gamma(m_{RD}) \Gamma(m_{SR})} \xi^{m_{SD}}(\gamma_{SD}, \gamma_{SR}) d\gamma_{SR} d\gamma_{SD}. \quad (26)$$

$A' = 1 - \frac{\gamma(m_{SR}, m_{SR} \frac{1}{\bar{\gamma}_{SR}})}{\Gamma(m_{SR})}$ indicates the probability of $p_f = 0$ (i.e., decoding error can be made arbitrarily small at R).

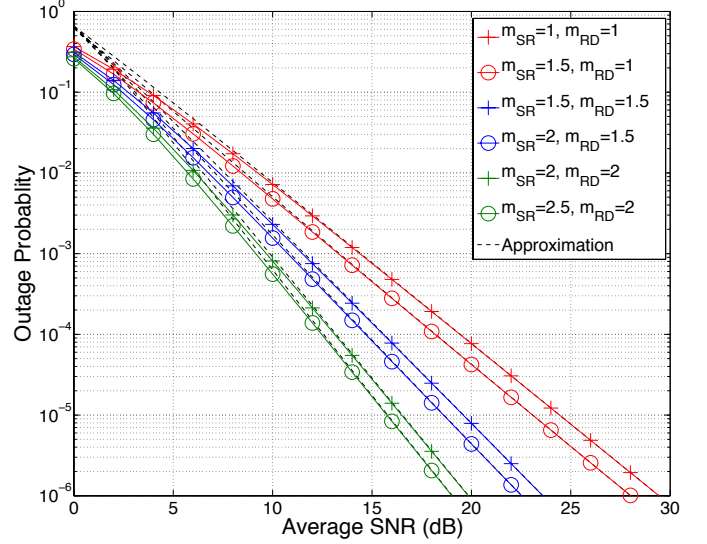


Fig. 7. Comparison of outage curves obtained by using the numerical calculation via (18), (19), and (20) and approximation via (21), (22), and (23), where $m_{SD} = 1$ and $G_{SD} = G_{RD} = G_{SR}$.

Fig. 7 shows that at high SNR region the approximated outage curves obtained from (21), (22), and (23) well match the numerically calculated curves obtained from (18), (19), and (20) using a numerical integral method. From this observation, it can be concluded that the approximation is sufficiently accurate to calculate the outage probability.

Since $\bar{\gamma}_{SD}$, $\bar{\gamma}_{RD}$, and $\bar{\gamma}_{SR}$ in (21), (22), and (23), are linear to a certain representative value $\bar{\gamma}$, and the exponent terms stay unchanged even if we replace $\bar{\gamma}_{SD}$, $\bar{\gamma}_{RD}$, and $\bar{\gamma}_{SR}$ by $\bar{\gamma}$, we have

$$P_A \approx A \cdot A' \cdot \bar{\gamma}^{(-m_{SD}-m_{RD})}, \quad (27)$$

$$P_B \approx B \cdot \bar{\gamma}^{(-m_{SD}-m_{SR})}, \quad (28)$$

$$P_C \approx C \cdot \bar{\gamma}^{(-m_{SD}-m_{RD}-m_{SR})}. \quad (29)$$

It is observed from (27), (28), and (29) that, $\bar{\gamma}^{(-m_{SD}-m_{RD}-m_{SR})}$ is upper bounded by $\bar{\gamma}^{(-m_{SD}-m_{RD})}$ or $\bar{\gamma}^{(-m_{SD}-m_{SR})}$ when $\bar{\gamma}$ approaches infinity. Moreover, it is not difficult to find that A' equals one asymptotically at high SNR regime. Therefore, the overall outage probability of the LF relaying with Nakagami- m fading can be formulated as

$$P_{out}^{LF} = (G_c \cdot \bar{\gamma})^{(-G_d)}, \quad (30)$$

where

$$G_d = m_{SD} + \min(m_{SR}, m_{RD}) \quad (31)$$

and

$$G_c = \begin{cases} \frac{1}{(G_d \sqrt{B})}, & m_{SR} < m_{RD} \\ \frac{1}{(G_d \sqrt{A+B})}, & m_{SR} = m_{RD} \\ \frac{1}{(G_d \sqrt{A})}, & m_{SR} > m_{RD} \end{cases} \quad (32)$$

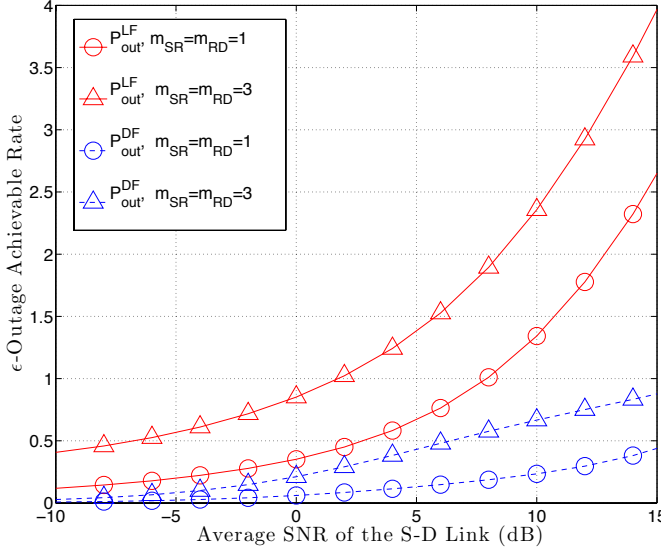


Fig. 8. The ϵ -outage achievable rates of LF and DF for $\epsilon = 0.01$, where $m_{SD} = 1$ and $G_{SD} = G_{RD} = G_{SR}$.

are the equivalent diversity order³ and coding gain [31] achieved by the LF relaying, respectively. We can see from (31) that, the equivalent diversity order of the LF relaying system is restricted by the less reliable channel of either the S-R or R-D link. The outage curves shown in Fig. 7 also illustrate that the decay in outage curve cannot be obtained by only increasing the m value of either the S-R or R-D link. The outage curves become sharper only when m values of both the S-R and R-D links increase simultaneously, which is consistent to the conclusion drawn in (31) regarding equivalent diversity order.

V. ϵ -OUTAGE ACHIEVABLE RATE

The ϵ -outage achievable rate R^ϵ is defined as the largest achievable transmission rate R_c such that the outage probability P_{out} is not larger than a threshold ϵ ,

$$R^\epsilon = \sup\{R_c : P_{out} \leq \epsilon\}, \quad 0 \leq \epsilon \leq 1. \quad (33)$$

For LF and DF, the ϵ -outage achievable rates R_{LF}^ϵ and R_{DF}^ϵ are, respectively, given by

$$R_{LF}^\epsilon = \sup\{R_c : P_{out}^{LF} \leq \epsilon\}, \quad 0 \leq \epsilon \leq 1, \quad (34)$$

$$R_{DF}^\epsilon = \sup\{R_c : P_{out}^{DF} \leq \epsilon\}, \quad 0 \leq \epsilon \leq 1, \quad (35)$$

The ϵ -outage achievable rates of the LF and DF relaying are illustrated in Fig. 8, where ϵ is set at 0.01. It is found that for both LF and DF, R_c^ϵ increases as the values of m_{SR} and m_{RD} increase. However, LF always outperforms DF in terms of R_c^ϵ . It can also be observed from Fig. 8 that at high

³ G_d indicates the decay of the outage probability curve versus average SNR $G_{ij} \frac{E_s}{N_0}$ because G_d appears in the exponent part in (30). Therefore, the decay of the outage curve is concerned and it has the equivalent meaning as the diversity order. However, according to the definition, G_d can take arbitrary real number which is related to the shape factor m_{ij} . In this sense, this paper calls the decay of the curve equivalent diversity order, which does not follow the standard definition of the diversity order.

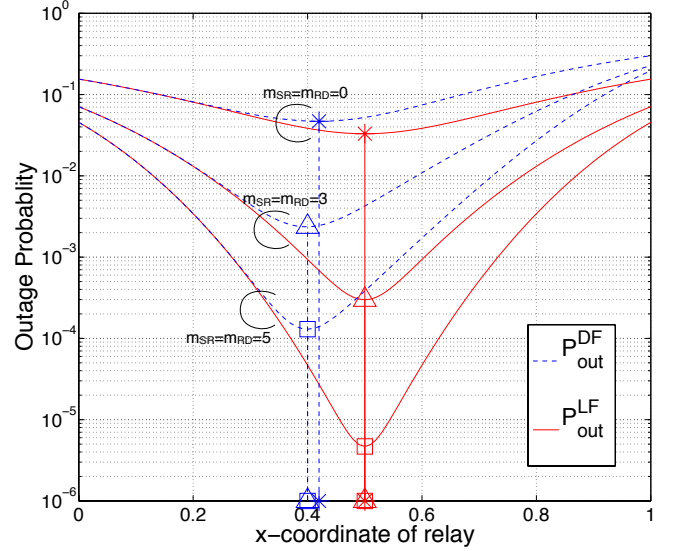


Fig. 9. The optimal relay positions of the LF and DF relaying where $m_{SD} = 1$ and $\alpha = 3.52$. The vertical lines indicate the points where the outage probabilities are minimized.

SNR region, the LF relaying is more advantageous than DF in terms of ϵ -outage achievable rate.

VI. OPTIMAL RELAY LOCATION

Let d_{SD} , d_{RD} , and d_{SR} denote the distances between S and D, R and D, and S and R respectively. With G_{SD} being normalized to unity, G_{SR} and G_{RD} can be defined as $G_{SR} = \left(\frac{d_{SD}}{d_{SR}}\right)^\alpha$ and $G_{RD} = \left(\frac{d_{SD}}{d_{RD}}\right)^\alpha$, respectively, where α is the path loss exponent. Then, the average SNRs of the S-R and R-D links can be given as

$$\bar{\gamma}_{SR} = \bar{\gamma}_{SD} + 10\log_{10}(G_{SR})(\text{dB}), \quad (36)$$

$$\bar{\gamma}_{RD} = \bar{\gamma}_{SD} + 10\log_{10}(G_{RD})(\text{dB}). \quad (37)$$

By substituting (36) and (37) into (18), (19), and (20), we can obtain the outage probability expression with respect to the position of R.

Fig. 9 shows the impact of the relay location on the outage probability, with $\bar{\gamma}_{SD} = 3$ dB. R is located on the line parallel to the S-D link between $x = 0$ and $x = 1$, as shown in Fig. 1. The outage probabilities of the DF relaying, are also depicted as the references. With DF, the lowest outage probability can be achieved when R is located close to S, while, interestingly, the lowest outage probability can be achieved when R is located at the midpoint with LF as long as $m_{SR} = m_{RD}$. It is also observed that the outage curves are symmetric with respect to the midpoint of the S-D link with the LF relaying. This is because with the LF relaying, the errors due to the S-R link can be corrected at the destination, and therefore, the midpoint ($d_{SR} = d_{RD}$) is the optimal point where the contributions of the S-R and R-D links are balanced. Note that, for ensuring certain outage probability, the LF relaying has a larger range to find a relay than the DF relaying.

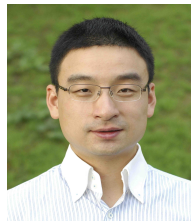
VII. CONCLUSION

The outage probability of the LF relaying has been analyzed over the independent Nakagami- m fading channels. The exact expression of the outage probability has been derived. It has been found that through a reasonable approximation of the outage expression, the equivalent diversity order and coding gain of the LF relaying can be derived. It has been revealed that the equivalent diversity order (outage curve decay) is dominated by the less reliable channel of either the S-R or R-D link. Compared with the DF relaying, LF can achieve lower outage probability. Moreover, it has been found that the outage probability of the LF relaying derived from GCC and that from CCC are almost the same. The ϵ -outage achievable rate of the LF relaying is always larger, compared with that of the DF relaying. Finally, it has been shown that since the LF relaying utilizes the contributions from both the S-R and R-D links, the optimal R position is exactly in the midline between S and D, so long as $m_{SR} = m_{RD}$. The observation about the relay location suggests that the search area for a relay (helper) can be increased by the LF relaying compared to the DF relaying.

Since the channel model practically depends on the position of nodes, joint analysis of the actual propagation property, the impact of node position, i.e., average SNR, and dynamic channel variation characteristics in terms of the higher order statistics is left for future study.

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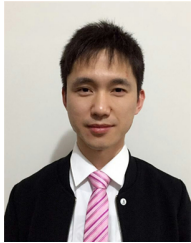


Shen Qian graduated from the School of Telecommunication Engineering (STE), Xidian University, Xi'an, China, in 2003. He received the master degree from School of Information Science, Japan Advanced Institute of Science and Technology (JAIST), in 2014. He is currently working toward the Ph.D degree in JAIST and Dr.Sc. degree in University of Oulu, under a dual doctoral degree program. His research interests are network information theory, cooperative wireless communications, and optimal power allocation problems.



Xiaobo Zhou (S'11-M'13) received the B.Sc. in electronic information science and technology from the University of Science and Technology of China (USTC), Hefei, China, the M.E. in computer application technology from Graduate University of Chinese Academy of Science (GUCAS), Beijing, China, and the Ph.D. degree from School of Information Science, Japan Advanced Institute of Science and Technology (JAIST), Ishikawa, Japan, in 2007, 2010, and 2013, respectively. From April 2014 to March 2015, he was a Researcher with the Department of

Communications Engineering, University of Oulu, Oulu, Finland. Currently, he is an Associate Professor with the School of Computer Science and Technology, Tianjin University, Tianjin, China. His research interests include cooperative wireless communications, network information theory, cloud computing, data centers, software-defined networking and future internet architecture.



Xin He received the D.Sc.(Tech.) degree in communications engineering from the University of Oulu, Oulu, Finland, and Ph.D. degree (with distinction) from the School of Information Science, Japan Advanced Institute of Science and Technology (JAIST), Ishikawa, Japan, in 2016. He is currently with the School of Mathematics and Computer Science, Anhui Normal University, Wuhu, China. His research interests include network information theory, energy harvesting networks and optimal coding/decoding design.



Jiguang He (S'16) received the B.Eng. degree in communications engineering from Harbin Institute of Technology, Harbin, China, in 2010, and M.Sc. degree also in communications engineering from Xiamen University, Xiamen, China, in 2013. From September 2013 to March 2015, he was with Key Laboratory of Millimeter Waves at City University of Hong Kong, conducting research on channel tracking over millimeter wave MIMO systems. Since June 2015, he has been with Centre for Wireless

Communications Engineering, University of Oulu, Oulu, Finland, studying toward his Ph.D degree. His research interests span cooperative communications, network information theory, joint source and channel coding, and distributed source coding.



Markku Juntti (S'93-M'98-SM'04) received the M.Sc. and Dr.Sc. degrees in electrical engineering from the University of Oulu, Oulu, Finland, in 1993 and 1997, respectively.

Dr. Juntti was with University of Oulu in 1992-1998. In academic year 1994-95, he was a Visiting Scholar at Rice University, Houston, Texas. In 1999-2000, he was a Senior Specialist with Nokia Networks. Dr. Juntti has been a professor of communications engineering since 2000 at University of Oulu, Centre for Wireless Communications (CWC), where he leads the Communications Signal Processing (CSP) Research Group. In 2014-2017, he also serves as Dean of University of Oulu Graduate School and currently also as Head of CWC Radio Technologies (RT) Research Unit. His research interests include signal processing for wireless networks as well as communication and information theory. He is an author or co-author in some 350 papers published in international journals and conference records as well as in books WCDMA for UMTS and Signal Processing Handbook. Dr. Juntti is also an Adjunct Professor at Department of Electrical and Computer Engineering, Rice University, Houston, Texas, USA.

Dr. Juntti is an Editor of IEEE Transactions on Communications and was an Associate Editor for IEEE Transactions on Vehicular Technology in 2002-2008. He was Secretary of IEEE Communication Society Finland Chapter in 1996-97 and the Chairman for years 2000-01. He has been Secretary of the Technical Program Committee (TPC) of the 2001 IEEE International Conference on Communications (ICC 2001), and the Co-Chair of the Technical Program Committee of 2004 Nordic Radio Symposium and 2006 IEEE International Symposium on Personal, Indoor and Mobile Radio Communications (PIMRC 2006), and the General Chair of 2011 IEEE Communication Theory Workshop (CTW 2011). He has served as Co-Chair of the Signal Processing for Communications Symposium of Globecom 2014 Signal Processing for Communications Symposium and IEEE GlobalSIP 2016 Symposium on Transceivers and Signal Processing for 5G Wireless and mm-Wave Systems.



Tad Matsumoto (S'84-SM'95-F'10) received the B.S., M.S., and Ph.D. degrees from Keio University, Yokohama, Japan, in 1978, 1980, and 1991, respectively, all in electrical engineering.

He joined Nippon Telegraph and Telephone Corporation (NTT) in April 1980. In July 1992, he transferred to NTT DoCoMo, Tokyo, Japan, where he researched Code-Division Multiple-Access techniques for Mobile Communication Systems. In April 1994, he transferred to NTT America, where he served as a Senior Technical Advisor of a joint project between NTT and NEXTEL Communications. In March 1996, he returned to NTT DoCoMo, where he served as the Head of the Radio Signal Processing Laboratory until August of 2001. He worked on adaptive signal processing, multiple-input multiple-output turbo signal detection, interference cancellation, and space-time coding techniques for broadband mobile communications. In March 2002, he moved to University of Oulu, Finland, where he served as a Professor with Centre for Wireless Communications. In 2006, he served as a Visiting Professor at Ilmenau University of Technology, Ilmenau, Germany, funded by the German MERCATOR Visiting Professorship Program. Since April 2007, he has been serving as a Professor with Japan Advanced Institute of Science and Technology (JAIST), Japan, while also keeping a cross appointment professorship at University of Oulu. He was appointed a Finnish Distinguished Professor for a period from January 2008 to December 2012, funded by the Finnish National Technology Agency (Tekes), Helsinki, Finland, and Finnish Academy, under which he preserves the rights to participate in and apply to European and Finnish national projects. He was the recipient of the IEEE VTS Outstanding Service Award (2001), Nokia Foundation Visiting Fellow Scholarship Award (2002), IEEE Japan Council Award for Distinguished Service to the Society (2006), IEEE Vehicular Technology Society James R. Evans Avant Garde Award (2006), and Thuringen State Research Award for Advanced Applied Science (2006), 2007 Best Paper Award of Institute of Electrical, Communication, and Information Engineers of Japan (2008), Telecom System Technology Award by the Telecommunications Advancement Foundation (2009), IEEE Communications Letters Exemplifying Reviewer Award (2011), and Nikkei Wireless Japan Award (2012). He is a member of IEICE. He served as an IEEE Vehicular Technology Distinguished Lecturer during the term July 2011-June 2015.