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Elliptic Curves over $F_p$
Suitable for Cryptosystems

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Abstract. Koblitz ([5]) and Miller ([6]) proposed a method by which the group of points on an elliptic curve over a finite field can be used for the public key cryptosystems instead of a finite field. To realize signature or identification schemes by a smart card, we need less data size stored in a smart card and less computation amount by it. In this paper, we show how to construct such elliptic curves while keeping security high.

1 Introduction

Public key cryptosystems based on the discrete logarithm problem on an elliptic curve (EDLP) can offer small key length cryptosystems. If an elliptic curve is chosen to avoid the Menezes-Okamoto-Vanstone reduction ([9]), then the only known attacks on EDLP are the Pollard $\rho$-method ([11]) and the Pohlig-Hellman method ([10]). So up to the present, such elliptic curve cryptosystems on $E/F_q$ are secure if $\#E(F_q)$ is divisible by a prime only more than 30 digits ([3]).

If we use an elliptic curve $E/F_q$ for digital signature or identification by a smart card ([12]), data size and computation amount of signature generation should be as small as possible. We may publish only the $x$-coordinate $x(P)$ of a public key $P$ and one bit necessary to recover the $y$-coordinate $y(P)$ of $P$ since the public key of an elliptic curve point is 2 times as large as the definition field $F_q$. Then we can reduce the data size to one half. But it will cause the computation amount to recover $y(P)$.

In this paper, we investigate an elliptic curve suitable for cryptosystems, in the sense that it requires less data size and less computation, while maintaining the security. We also show the advantage of our elliptic curve in the case of the Schnorr’s digital signature scheme on an elliptic curve.

This paper is organized as follows. Section 2 summarizes the addition formula of an elliptic curve ([13]). Section 3 describes the Schnorr signature on an elliptic curve, and show the data size and the computation amount for two cases, the basic version and the reducing-data version. Section 4 discusses the elliptic curve which gives cryptosystems that reduce both of data sizes and the computation amount.

2 Addition formula of Elliptic curve

Cryptosystems on an elliptic curve $E/F_q$, for example the Diffie-Hellman key distribution and ElGamal cryptosystems, require the computation of $kP$ ($P \in$
$E(F_q)).$ We will discuss the computation amount of $kP.$ For simplicity, we neglect addition, subtraction and multiplication by a small constant in $F_q$ because they are much faster than multiplication and division in $F_q.$

Let $K$ be a finite field $F_q$ of characteristic $\neq 2, 3.$ An elliptic curve over $K$ is given as follows,

$$E : y^2 = x^3 + ax + b \ (a, b \in K, 4a^3 + 27b^2 \neq 0).$$

Then the set of $K$-rational points on $E$ (with a special element $O$ at infinity), denoted $E(K)$, is a finite abelian group, where $E(K) = \{(x, y) \in K^2 | y^2 = x^3 + ax + b \} \cup \{O\}.$ For the curve $E,$ the addition formulas in the affine coordinate are the following. Let $P = (x_1, y_1), \ Q = (x_2, y_2)$ and $P + Q = (x_3, y_3)$ be points on $E(K)$.

- **Curve addition formula in the affine coordinates** $(P \neq \pm Q)$
  
  \begin{align*}
  x_3 &= \lambda^2 - x_1 - x_2, \\
  y_3 &= \lambda(x_1 - x_3) - y_1, \\
  \lambda &= \frac{y_2 - y_1}{x_2 - x_1},
  \end{align*}

- **Curve doubling formula in the affine coordinates** $(P = Q)$
  
  \begin{align*}
  x_3 &= \lambda^2 - 2x_1, \\
  y_3 &= \lambda(x_1 - x_3) - y_1, \\
  \lambda &= \frac{3x_1^2 + a}{2y_1}.
  \end{align*}

The formula (1) requires two multiplications and one division in $K,$ while the formula (2) requires three multiplications and one division in $K.$ The computation amount of division in $K$ is more than that of multiplication in $K.$ So we often use the projective coordinates to avoid divisions in $K.$ The addition formulas in the projective coordinates are the following. Let $P = (X_1, Y_1, Z_1), \ Q = (X_2, Y_2, Z_2)$ and $P + Q = (X_3, Y_3, Z_3)$.

- **Curve addition formula in the projective coordinates** $(P \neq \pm Q)$
  
  \begin{align*}
  X_3 &= vA, \\
  Y_3 &= u(v^2X_1Z_2 - A) - v^3Y_1Z_2, \\
  Z_3 &= u^3Z_1Z_2,
  \end{align*}

  where $u = Y_2Z_1 - Y_1Z_2, \ v = X_2Z_1 - X_1Z_2, \ t = X_2Z_1 + X_1Z_2, \ A = u^2Z_1Z_2 - v^2t;$

- **Curve doubling formula in the projective coordinates** $(P = Q)$
  
  \begin{align*}
  X_3 &= 2hs, \\
  Y_3 &= w(4B - h) - 8Y_1^2s^2, \\
  Z_3 &= 8s^3,
  \end{align*}

  where $w = aZ_1^2 + 3X_1^2, \ s = Y_1Z_1, \ B = X_1Y_1s, \ h = w^2 - 8B.$ The formula (3) requires 15 multiplications, while the formula (4) requires 12 multiplications.
For the use of cryptosystems, we may set \( z(P) = Z_1 \) to one in the formula (3). Then the formula (3) requires 12 multiplications.

Subtractions are as expensive as additions over elliptic curves. So the computation amount of \( kP \) by the addition-subtraction method ([2, 8]) is less than that by the binary method, while both methods need memory storage only for \( P \). We assume to compute \( kP \) by the addition-subtraction method. The computation by the addition-subtraction requires \( n \) times of curve doubling and \( \frac{3n}{2} \) times of curve adding on the average, where \( n = \vert K \vert \). Computation of \( kP \) in the projective coordinate requires one division and two multiplications in the final stage. Since \( n \) is larger than about 100, the computations in the projective coordinates are faster than that in the affine coordinates if the ratio of the computation amount of division in \( K \) to that of multiplication in \( K \) is larger than 9. In order to compare the computation amount of Schnorr signature scheme on a finite field and on an elliptic curve, we assume to compute \( kP \) in the projective coordinate by the addition-subtraction method and compute the power residue by the binary method.

3 Elliptic curve cryptosystems

If \( E(K) \) and a basepoint \( P \in E(K) \) are carefully chosen, then the only known attacks on the cryptosystems are the square root attacks. EDLP on such \( E \) to the base \( P \) is secure up to the present ([3]), if the order of \( P \), \( \text{ord}(P) \), is divisible by more than a 30-digit prime. Here we summarize the Schnorr signature on such an elliptic curve and establish a basis for evaluation of the elliptic curve proposed in the next chapter. Let \( M \in \mathbb{Z} \) be a message. User A sends the message \( M \) to user B with her or his signature of \( M \).

- Initialization
  - system parameter
    - \( E : y^2 = x^3 + ax + b \) (\( a, b \in F_p \); \( p \) is a prime of \( n(\geq 97) \) bits).
    - \( P \in E(F_p) \): a basepoint (chosen as the above).
    - \( l = \text{ord}(P) \) (\( l \) is \( m(\geq 97) \) bits).
  - a one-way hash function \( h : \mathbb{Z}_t \times \mathbb{Z} \to \{0, \ldots, 2^t - 1\} \), where \( t \) is the security parameter.

- Key generation
  - User A randomly chooses an integer \( s \), a secret key, and makes public the point \( P_A = -sP \) as a public key.

- Signature generation
  1. Pick a random number \( k \in \{1, \ldots, l\} \) and compute
     \[ R = kP = (r_x, r_y). \] (5)
     Here \( r_x = x(R) \) and \( r_y = y(R) \).
  2. Compute \( e := h(r_x, M) \in \{0, \ldots, 2^t - 1\} \).
  3. Compute \( y \equiv k + se \pmod{l} \) and output the signature \((e, y)\).
– Signature verification

1. Compute \( \overline{R} = yP + eP_A = (\overline{r_x}, \overline{r_y}) \) and check that \( e = h(\overline{r_x}, M) \).

As we described in Section 2, the computation of \( kP \) requires \( m \) curve doublings and \( \frac{m}{2} \) curve additions on the average, where \( k \) is a \( m \)-bit number. Extending the addition-subtraction method to the computation in the verification, we can calculate \( yP + eP_A \) in \( m \) curve doublings and \( \frac{m}{2}(m - t) + \frac{m}{2}t \) curve additions on the average with precomputations of \( \pm(P \pm P_A) \), which require about the same computational amount as one curve addition.

Here we set \( n, m = 128 \). Then the known attacks on such an elliptic curve cryptosystems requires at least \( 2^{14} \) elliptic curve operations. This is roughly equal to that of the original Schnorr on \( F_p \) (\( p \) is 512 bits). If lower security is required, then \( n, m \) can be replaced by a smaller number like 97. For the security parameter, here we set \( t = 128 \).

We will present two versions of Schnorr signature on an elliptic curve. One is the basic Schnorr signature on an elliptic curve described above, called Basic EC version. Another is called Reducing data EC version. In this version, only \( x(P_A) \) and the least significant bit of \( y(P_A) \) are published as a public key to reduce the data size. The same is done for the basepoint \( P \). On the other hand, the original Schnorr signature scheme on \( F_p \), called Finite field version (\( p \) is 512 bits, the security parameter \( t = 128 \)) roughly has the same security as that on the above elliptic curves. So the size of the definition field of Finite field version is four times as large as that of Basic and Reducing data EC versions.

We compare Basic EC version, Reducing data EC version and Finite field version, with respect to data size. Table 1 shows the comparison.

• Basic EC version

The system key is \( a, p, P, \) and \( l \) (640 bits). The secret key is \( s \) (128 bits). They are stored in a smart card. So the data size stored in a smart card is 768 bits. The public key is \( (P_A) \) (256 bits) and the signature is \( e \) and \( y \) (256 bits).

• Reducing data EC version

In this version, we have to publish one more parameter \( a, b, P, x(P) \) as a system key to recover a point by the \( x \)-coordinate of the point and the least significant bit of the \( y \)-coordinate of the point. It requires no residue to recover the \( y \)-coordinate of \( P \) and increases computation for signature. The system key is \( a, b, p, x(P), \) the least significant bit of \( y(P) \) and \( l \) (641 bits). The secret key is \( s \) (128 bits). So the data size stored in a smart card is 769 bits. It is almost equal to that of Basic EC version. The public key is \( (x(P_A) \) and the least significant bit of \( y(P_A) \) (129 bits) and the signature is \( e \) and \( y \) (256 bits).

• Finite field version

The system key of Finite field version is a set of the definition field, the basepoint and the order of basepoint (1164 bits), where the size of the definition field is 512 bits and the order of basepoint is 140 bits. The secret key is 140 bits. So the data size stored in a smart card is 1304 bits.

The size of the definition fields of both EC versions is reduced to 25% of Finite field version. But the stored data size is not so reduced (39%). This is because an elliptic curve point has 2 coordinates and we need a parameter to
decide $E$.

Let us compare the three cases with respect to the computation amount. We assume the computation method that we described in Section 2. Table 2 shows the comparison of the computation amount of signature generation and verification. Here we assume $m(n) = (n/t)^2 m(t)$, where $m(n)$ denotes the amount of work to perform one modular multiplication whose modulus size is $n$ bits. We assume the ratio of the computation amount of division in $K$ to that of multiplication in $K$ to 10. We see the computation amount of signature generation of Reducing data EC version is reduced to 67% of Finite field version. It is not so reduced as the size of the definition field. This is because the computation amount of one elliptic curve addition is much more than that of one multiplication in the same definition field and we need to recover a basepoint.

We see that both EC versions seem to be better than Finite field version for both points of the data size and the computation amount. But actually they are not so efficient considering the less size of the definition field of $E$. For the stored data size, the ratio of the stored data size to the definition field for both EC versions is 6. On the other hand, for Finite field version, the ratio is 2.5. For the computation amount, one elliptic curve addition requires about 12 multiplications. If we require higher security, for example $t = 160$, then we will have to construct an elliptic curve over at least a 160-bit finite field. Then the advantage for EC versions shown in Table 1 and 2 decreases.

<table>
<thead>
<tr>
<th></th>
<th>System Key</th>
<th>Secret Key</th>
<th>Public Key</th>
<th>Signature size</th>
</tr>
</thead>
<tbody>
<tr>
<td>Basic EC version</td>
<td>640</td>
<td>128</td>
<td>256</td>
<td>256</td>
</tr>
<tr>
<td>Reducing data EC version</td>
<td>641</td>
<td>128</td>
<td>129</td>
<td>256</td>
</tr>
<tr>
<td>Finite field version</td>
<td>1164</td>
<td>140</td>
<td>512</td>
<td>268</td>
</tr>
</tbody>
</table>

Table 1. Comparison of data size(in bits)

<table>
<thead>
<tr>
<th></th>
<th>Signature Generation</th>
<th>Signature Verification</th>
</tr>
</thead>
<tbody>
<tr>
<td>Basic EC version</td>
<td>129</td>
<td>151</td>
</tr>
<tr>
<td>Reducing data EC version</td>
<td>141</td>
<td>175</td>
</tr>
<tr>
<td>Finite field version</td>
<td>210</td>
<td>242</td>
</tr>
</tbody>
</table>

Table 2. Comparison of the computation amount(number of 512-bit modular multiplications)

In the next section, we construct an elliptic curve cryptosystem, which has
4 Elliptic curves suitable for Cryptosystems

If $E(F_p)$ and the basepoint $P \in E(F_p)$ are appropriately chosen, then the only known attacks on the cryptosystems are the square root attacks. We first discuss a method to construct such elliptic curves and then investigate what elliptic curve among them is suitable for implementation with respect to less data size (key length) and less computation amount.

4.1 Decision of the class of elliptic curves

One method to avoid the recent attack is to construct EDLP on $E/F_p$ with $p$ elements ([7]). We describe a modified method to decide the class of such elliptic curves. There are two phases for the decision of $E/F_p$ with $p$ elements.

The first phase is to find an appropriate prime $p$. Such $p$ is a form of $p = db^2 + db + \frac{4d+1}{4}$ ($b$ is an integer) for $d \in \{3, 11, 19, 43, 67, 163\}$. Such integers $d$ enable us to construct easily the $j$-invariant $j_d$ of $E/F_p$ with $p$ elements for the prime $p$, which is uniquely determined by $d$. Table 3 lists integers $d$ and the $j$-invariant $j_d$.

<table>
<thead>
<tr>
<th>$d$</th>
<th>$j_d$</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>$0$</td>
</tr>
<tr>
<td>11</td>
<td>$(-2^5)^3$</td>
</tr>
<tr>
<td>19</td>
<td>$(-2^3 \cdot 3)^3$</td>
</tr>
<tr>
<td>43</td>
<td>$(-2^6 \cdot 5)^3$</td>
</tr>
<tr>
<td>67</td>
<td>$(-2^3 \cdot 3 \cdot 5 \cdot 11)^3$</td>
</tr>
<tr>
<td>163</td>
<td>$(-2^3 \cdot 3 \cdot 5 \cdot 23 \cdot 29)^3$</td>
</tr>
</tbody>
</table>

Table 3. Integers $d$ and $j$-invariant $j_d$

Once the prime $p = db^2 + db + \frac{4d+1}{4}$ and $j_d$ are given, then the next phase is to decide the class of $E/F_p$ with $p$ elements. There is a little difference between the case of $d = 3$ and others. First we investigate the case of $d \in \{11, 19, 43, 67, 163\}$. Then the elliptic curves over $F_p$ with the $j$-invariant $j_d$ are given as follows.

$$E_{c,d}: y^2 = x^3 + 3c^2a_dx + 2c^3a_d, a_d = \frac{j_d}{1728 - j_d} \quad (\forall c \in F_p^*).$$
For each $d$, we can classify $\{E_{c,d}|c \in F_p^*\}$ into two equivalence classes of twists, namely

$$\mathcal{E}_d = \{E_{c,d}|c \in F_p^*, \left(\frac{c}{p}\right) = 1\} \quad \text{and} \quad \mathcal{E}'_d = \{E_{c,d}|c \in F_p^*, \left(\frac{c}{p}\right) = -1\},$$

where $\left(\frac{c}{p}\right)$ denotes the Legendre symbol. Then only one of the two classes gives the elliptic curves with $p$ elements. A general condition to decide the class was investigated ([1]). In our case, the condition can be simplified as follows.

**Theorem 1.** Let $p$ be a prime represented by $p = db^2 + db + \frac{d+1}{4}$ ($b$ is an integer) for $d \in \{11, 19, 43, 67, 163\}$. Then the class which gives elliptic curves with $p$ elements is determined as:

$$\mathcal{E}_d \quad \text{if} \quad \left(\frac{\alpha_d}{p}\right) = -1,$$

$$\mathcal{E}'_d \quad \text{if} \quad \left(\frac{\alpha_d}{p}\right) = 1,$$

where $\alpha_d$ is an integer determined by $d$. Table 4 shows the values of $\alpha_d$.

<table>
<thead>
<tr>
<th>$d$</th>
<th>$\alpha_d$</th>
</tr>
</thead>
<tbody>
<tr>
<td>11</td>
<td>3 * 7</td>
</tr>
<tr>
<td>19</td>
<td>3</td>
</tr>
<tr>
<td>43</td>
<td>2 * 5 * 7</td>
</tr>
<tr>
<td>67</td>
<td>3 * 7 * 11 * 31</td>
</tr>
<tr>
<td>163</td>
<td>2 * 3 * 5 * 7 * 11 * 19 * 23 * 29 * 127</td>
</tr>
</tbody>
</table>

**Table 4.** Integers $d$ and $\alpha_d$

Now we get the following procedure to decide the class of elliptic curves with $p$ elements.

**Procedure 1**

1. Search a large prime $p$ such that $p = db^2 + db + \frac{d+1}{4}$ ($b$ is an integer) for $d \in \{11, 19, 43, 67, 163\}$.
2. Calculate $\left(\frac{\alpha_d}{p}\right)$. If $\left(\frac{\alpha_d}{p}\right) = -1$, then $\mathcal{E}_d$ is the class. Else if $\left(\frac{\alpha_d}{p}\right) = 1$, then $\mathcal{E}'_d$ is the class.

Next we will investigate the case of $d = 3$. Then the elliptic curves over $F_p$ ($p = 3b^2 + 3b + 1$) with the $j$-invariant $j_d$ are given as follows.

$$E_{\xi} : y^2 = x^3 + \xi \quad (\forall c \in F_p^*).$$

(6)
In this case, we can classify $\{E_\xi|c \in F^*_p\}$ into six equivalence classes of twists, namely

$$E_{3,i} = \{E_\xi|xi \in F^*_p, \left(\frac{\xi}{p}\right)_6 = (-\omega)^i\} \quad (0 \leq i \leq 5, \omega = \frac{-1 + \sqrt{-3}}{2}),$$

where $\left(\frac{\xi}{p}\right)_6$ denotes the sixth power residue symbol. Then exactly one of the six classes gives the elliptic curves with $p$ elements. We have a next formula on the number of rational points of the elliptic curves (6).

**Theorem 2 ([4]).** If $p \equiv 1 \pmod{3}$, let $p = \pi \overline{\pi}$ with $\pi \in \mathbb{Z}[\omega]$ and $\pi \equiv 2 \pmod{3}$. Then

$$\#E_\xi(F_p) = p + 1 + \left(\frac{4\xi}{\pi}\right)_6 + \left(\frac{4\xi}{\overline{\pi}}\right)_6 \pi.$$  \hfill (7)

Using the formula (7), the condition to decide the class can be given as follows.

**Theorem 3.** Let $p$ be a prime represented by $p = 3b^2 + 3b + 1$ ($b$ is an integer). Then the class which gives elliptic curves with $p$ elements is determined as:

- $E_{3,1}$ if $b \equiv 0, 2, 4 \pmod{6}$,
- $E_{3,5}$ if $b \equiv 1, 3, 5 \pmod{6}$.

**Proof.** We prove only the case of $b \equiv 1 \pmod{6}$. As for the other cases, we can do the same way. Let $\pi = (2b + 1)\omega + (b + 1)$. Then $p = \pi \overline{\pi}$ and $\pi \equiv 2 \pmod{3}$. Since $\left(\frac{\xi}{\pi}\right)_6 = \omega$, we get that $\#E_\xi(F_p) = p$ if and only if

$$\left(\frac{\xi}{\pi}\right)_6 \omega^{3\pi} + \left(\frac{\xi}{\overline{\pi}}\right)_6 \omega^{\overline{\pi}} = -1,$$

that is, $tr(\omega \left(\frac{\xi}{\pi}\right)_6 \pi) = -1$. So we get $\left(\frac{\xi}{\pi}\right)_6 = -\omega^2$. This means that the class which gives elliptic curves with $p$ elements is $E_{3,5}$.

Now we get the following procedure to decide the class of elliptic curves with $p$ elements.

**Procedure 2**

1. Search a large prime $p$ such that $p = 3b^2 + 3b + 1$ ($b$ is an integer).
2. If $b \equiv 0, 2, 4 \pmod{6}$, then $E_{3,1}$ is the class. Else if $b \equiv 1, 3, 5 \pmod{6}$, then $E_{3,5}$ is the class.

We have seen that the time to decide the class of $E/F_p$ with $p$ elements depends on the time finding $p = db^2 + db + \frac{d+1}{4}$ for $d \in \{3, 11, 19, 43, 67, 163\}$. We can easily find such a prime. In fact we were convinced experimentally that finding a prime $p = db^2 + db + \frac{d+1}{4}$ in the range of $30 \sim 90$ digits is as easy as finding a prime in that range. So we can easily decide the class of $E/F_p$ with $p$ elements which gives secure cryptosystems.
4.2 Selection of an elliptic curve and a basepoint

Elliptic curve cryptosystems require the computation of $kP$, where $P = (X_1, Y_1, 1)$ is a fixed point called basepoint. It is accomplished by repeated doubling, adding and subtracting of $P$. If we can select a basepoint $P$ with a small $x$-coordinate $X_1$ or a small $y$-coordinate $Y_1$, the amount of computation of $kP$ will be reduced. Especially in the case of signature and identification by a smart card, reducing of total data size stored in a smart card and the computation amount by a smart card is important. If fewer parameters represent an elliptic curve and a basepoint, the data stored in a smart card is reduced. Furthermore we wish to recover $P$ easily from the parameters.

In the last section, we have decided the class of elliptic curves which gives the secure cryptosystems. Note that any elliptic curve $E/F_p$ of the class and any basepoint $P \in E(F_p)$ give cryptosystems with the same security. We will discuss how to select $E$ of the class and $P$ in $E$ suitable for cryptosystems, in the sense that it reduces computation amount of $kP$ and necessary data size to be stored. We will classify $d$ into two cases, $d = 3$ and others.

- **Proposed scheme A**

  First we deal with the case of $d \in \{11, 19, 43, 67, 163\}$. For a given $p = db^2 + db + \frac{d+1}{4}$, we know which class, $E_d$ or $E'_d$, gives an elliptic curve with $p$ elements in Section 4.1. Without loss of generality, we will discuss the case of $E_d$. Let $y_0 = x_0^3 + 3a_dx_0 + 2a_y$ for $x_0 \in F_p$. Then we get one elliptic curve in $E_d$ and the basepoint following (8).

\[
E_d \ni E_{y_0,d}, \quad E_{y_0,d} \ni P = (y_0x_0, y_0^2) \text{ if } \left( \frac{y_0}{p} \right) = 1 \tag{8}
\]

If $y_0$ satisfies the condition of (8) for $x_0 = 0$, then we get $E_d \ni E_{y_0,d}$ and $E_{y_0,d} \ni P = (0,4a_y^2)$. In fact such $y_0$ satisfies the condition of (8) if and only if

\[
\left( \frac{y_0}{p} \right) = \left( \frac{2a_y}{p} \right) = 1.
\]

Except for $d = 19$, there exists $p = db^2 + db + \frac{d+1}{4}$ which satisfies $\left( \frac{2a_y}{p} \right) = 1$.

Combining the condition on $p$ to decide a class (i.e. $\left( \frac{2a_y}{p} \right) = -1$ or 1), we obtain that such an elliptic curve over $F_p$ exists if and only if $\left( \frac{2a_y}{p} \right) = -1$ in both cases, $E_d$ and $E'_d$. Table 5 shows the value of $\beta_d$.

We were also convinced experimentally that, for all $p = db^2 + db + \frac{d+1}{4}$ ($d \in \{11, 43, 67, 163\}$), such an elliptic curve exists with a probability of about one half. Here is one example for a 128-digit prime in the case of $d = 11$.

\[
E : y^2 = x^3 + 12a^3x + 16a^4; \quad E(F_p) \ni P = (0,4a^2),
\]

\[
p = 17014118346046923956673067859662240771716369,
\]

\[
a = 52713357398698261607887303078701255349.
\]
Let us use this elliptic curve $E_{y_0, d}$ and basepoint $P = (0, 4a_d^2)$ for Schnorr signature, where $E_{y_0, d} = E$ and $a = a_d$. We further assume that the public key $P_A$ is represented by $x(P_A)$ and the least significant bit of $y(P_A)$. The computation of $kP$ requires the addition to the basepoint $P$, which is calculated in 9 modular multiplications. So the computation of $kP$ requires $1932m(128)$. We can recover the basepoint in one modular multiplication, only if we store $a_d$. Since $ord(P)$ equals $p$, the system key is $a_d$ and $p$ (256 bits). Table 6 shows the data size and Table 7 shows the computation amount. The data size stored in a smart card is reduced to one half of that of Reducing data EC version and Basic EC version. The public key size is the same as that of Reducing data EC version.

The computation amount of the signature generation is reduced by 6% (resp. 14%) of that of Basic EC version (resp. Reducing data EC version). The computation amount of the signature verification is reduced by 10% of that of Reducing data EC version. It is increased by 5% of that of Basic EC version. This is because we need one power residue to recover one’s public key in the signature verification. If we publish $P_A$ instead of $x(P_A)$ and the least significant bit of $y(P_A)$ as a public key, then the computation amount of the signature verification is reduced by 3% of that of Basic EC version. Even in this case, the public key size is only 50% of Finite field version.

We can choose a prime $p$ and an elliptic curve $E/F_p$ as follows.

$$E : y^2 = x^3 + 12x^3 + 16a^4; \quad E(F_p) \ni P = (0, 4a^2),$$

$$p = 2^{128} - 89 235 631 848 000 472 139 4037$$

$$a = 1887 65172 00252 43003 83780 59753 00282 08521$$

The form of $p$ simplifies the arithmetic modulo $p$ and we can store $p$ with only 73 bits. Of course, the particular form of $p$ provides no disadvantage on the security for now.

- **Proposed scheme B**

Next we deal with the case of $d = 3$. For a given $p = 3b^2 + 3b + 1$, we know which class, $\mathcal{E}_{3,1}$ or $\mathcal{E}_{3,5}$, gives the elliptic curve with $p$ elements in Section 4.1. We only discuss the case of $\mathcal{E}_{3,1}$. As for the other case, we can do in the same way.
An elliptic curve \( E/F_p \) with \( p \) elements and a basepoint \( P \) is given as follows,
\[
E_\xi : y^2 = x^3 + \xi y_0^3; \ E_\xi(F_p) \ni P = (x_0y_0, y_0^2),
\]
\[
(\forall \xi \text{ such that } \left( \frac{\xi}{p} \right)_6 = -\omega, \ \forall y_0 = x_0^3 + \xi \in F_p^{*2}).
\]
In this case, there doesn't exist an elliptic curve with the point whose \( x \)-coordinate equals 0 because of \( \xi \not\in F_p^{*2} \). But we can select a small \( \xi \) such that \( \left( \frac{\xi}{p} \right)_6 = -\omega \) and a small \( x_0 \) such that \( y_0 = x_0^3 + \xi \in F_p^{*2} \). Here is one example for a 128-digit prime.
\[
E : y^2 = x^3 + 3 \cdot 4^3; \ E(F_p) \ni P = (4,16),
p = 1701\,411\,834\,604\,692\,480\,631\,577\,203\,304\,937\,639\,647
\]
\[
(x_0 = 1, \ \xi = 3)
\]
Let us use the elliptic curve \( E_\xi \) and the basepoint \( P = (x_0y_0, y_0^2) \) for Schnorr signature. We further assume that one's public data \( P_A \) is represented by \( x(P_A) \) and the least significant bit of \( y(P_A) \). Then the addition to \( P = (x_0y_0, y_0^2) = (X_1, Y_1) \) is accomplished in 9 modular multiplications because we can neglect the multiplications by a small constants \( X_1 \) and \( Y_1 \). Furthermore the simple equation of \( E \) reduces the computation amount of doubling. It is accomplished in 10 modular multiplications. As for the computation amount of \( kP \), it requires 1676\,m(128). As for the recovering the basepoint, we can recover it in a negligible computation amount only if we store \( x_0 \) and \( \xi \) whose data size is enough small. As for the data size, the data size of \( x_0 \) and \( \xi \) is neglected and \( \text{ord}(P) \) equals \( p \). So the size of system parameters \( x_0, \xi \) and \( p \) of Schnorr signature scheme on such \( E_\xi \) is about the same as that of the definition field. Table 6 shows the data size and Table 7 shows the computation amount.

We see that the elliptic curves and the basepoints in the case of \( d = 3 \) give good properties for the cryptosystems, especially in the application of digital signature and identification by a smart card. The data size stored in a smart card is reduced to one third of that of Reducing data EC version and Basic EC version. The public key size is the same as that of Reducing data EC version. The computation amount of the signature generation is reduced by 19\% (resp. 26\%) of that of Basic EC version (resp. Reducing data EC version). The computation amount of the signature verification is reduced by 6\% (resp. 19\%) of Basic EC version (resp. Reducing data EC version). If we publish \( P_A \) as a public key, then the computation amount of the signature verification is reduced by 14\% of that of Basic EC version.

In the same way as Proposed scheme A, we can choose a prime \( p \) and an elliptic curve \( E/F_p \) as follows.
\[
E : y^2 = x^3 + 3 \cdot 4^3; \ E(F_p) \ni P = (4,16),
p = 2^{128} - 86\,617\,554\,926\,587\,606\,009\,85
\]
\[
(x_0 = 1, \ \xi = 3)
\]
The form of \( p \) simplifies the arithmetic modulo \( p \) and we can store \( p \) with only 73 bits.
<table>
<thead>
<tr>
<th></th>
<th>System Key</th>
<th>Secret Key</th>
<th>Public Key</th>
<th>Signature size</th>
</tr>
</thead>
<tbody>
<tr>
<td>Proposed scheme A</td>
<td>256(201)</td>
<td>128</td>
<td>129</td>
<td>256</td>
</tr>
<tr>
<td>Proposed scheme B</td>
<td>131 (76)</td>
<td>128</td>
<td>129</td>
<td>256</td>
</tr>
</tbody>
</table>

**Table 6.** Data size of the Proposed schemes (in bits)

<table>
<thead>
<tr>
<th></th>
<th>Signature Generation</th>
<th>Signature Verification</th>
</tr>
</thead>
<tbody>
<tr>
<td>Proposed scheme A</td>
<td>121</td>
<td>158</td>
</tr>
<tr>
<td>Proposed scheme B</td>
<td>105</td>
<td>142</td>
</tr>
</tbody>
</table>

**Table 7.** Computation amount of the Proposed schemes (number of 512-bit modular multiplications)

5 Conclusion

Elliptic curve cryptosystems often require the computation of $kP$, where $P$ is a fixed basepoint. We have proposed the elliptic curves and basepoints suitable for cryptosystems, in the sense that they require less data size and less computation amount for $kP$. Especially if we use the Proposed version B in Schnorr signature scheme by a smart card, we have seen that

1. the data size stored in a smart card is reduced to one third of that of Basic EC version and Reducing data EC version;
2. the data size of public key is reduced to one half of that of Basic EC version and is the same as Reducing data EC version;
3. the computation amount of the signature generation is reduced by 19% (resp. 26%) of that of Basic EC version (resp. Reducing data EC version);
4. the computation amount of the signature verification is reduced by 6% (resp. 19%) of Basic EC version (resp. Reducing data EC version);
5. in the case where we publish the point $P_A$ as a public key, the computation amount of the signature verification is reduced by 14% of that of Basic EC version.

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References


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