

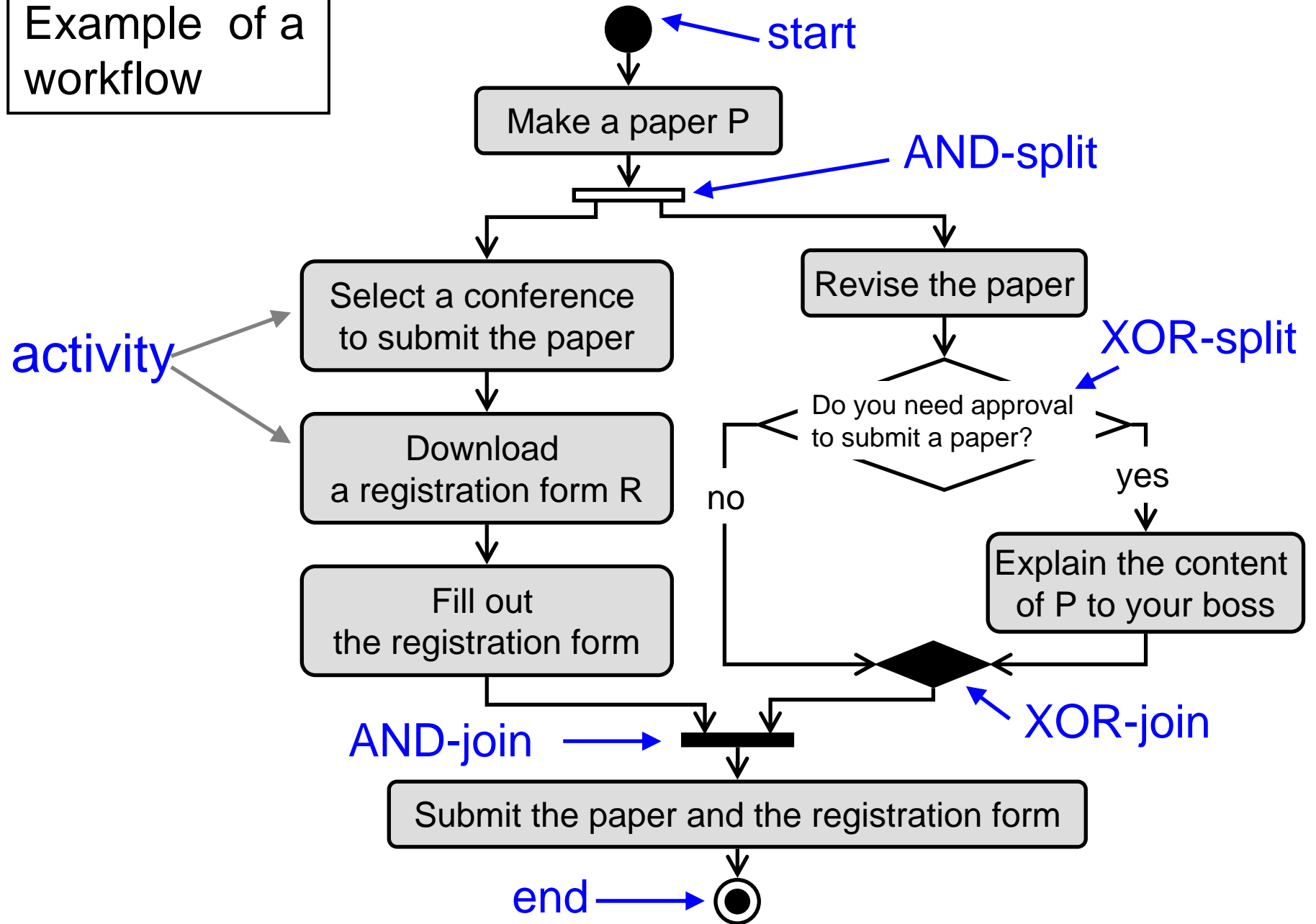
Title	Correctness Properties for Workflows with Multiple Starts and/or Ends
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Citation	
Issue Date	2008-03-03
Type	Presentation
Text version	publisher
URL	http://hdl.handle.net/10119/8288
Rights	
Description	5th VERITE : JAIST/TRUST-AIST/CVS joint workshop on VERification TEchnologyでの発表資料, 開催 : 2008年3月3日, 開催場所 : 北陸先端科学技術大学院大学・情報科学研究科棟 5F コラボレーションルーム 7, JAIST 21世紀COEシンポジウム 2008「検証進化可能電子社会」と共催

Correctness Properties for Workflows with Multiple Starts and/or Ends

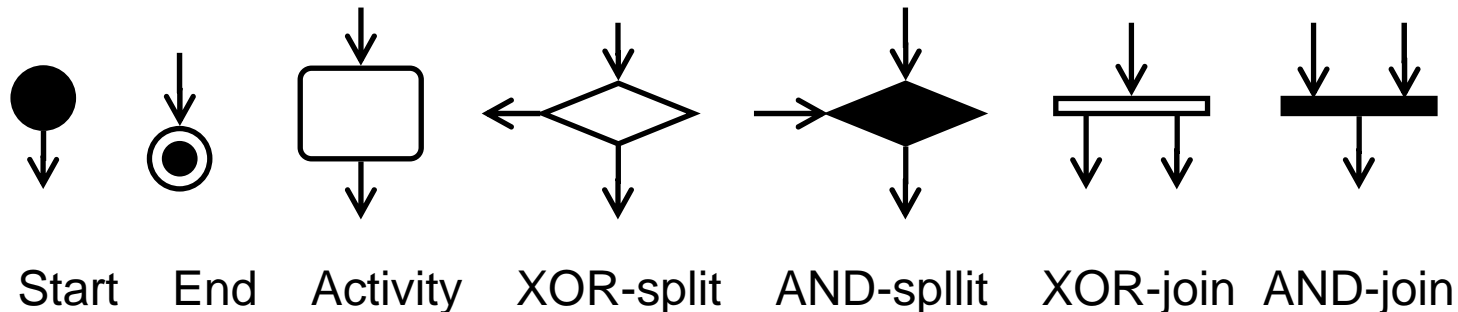
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- Contents of this talk
 - Motivation
 - Correctness of Workflows
 - Why do we consider multiple starts and/or ends?
 - General Correctness of Workflows with Multiple starts and/or Ends

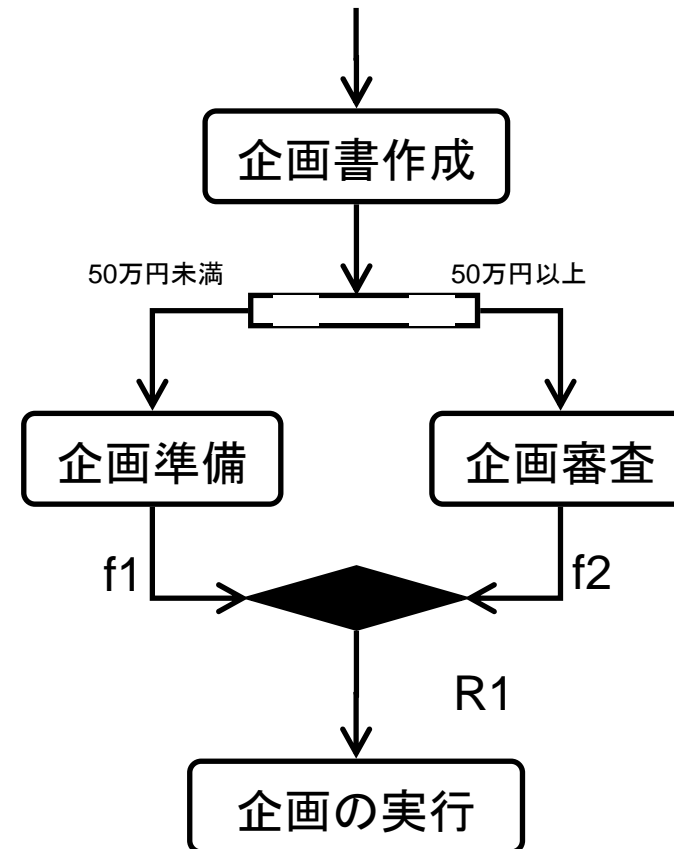
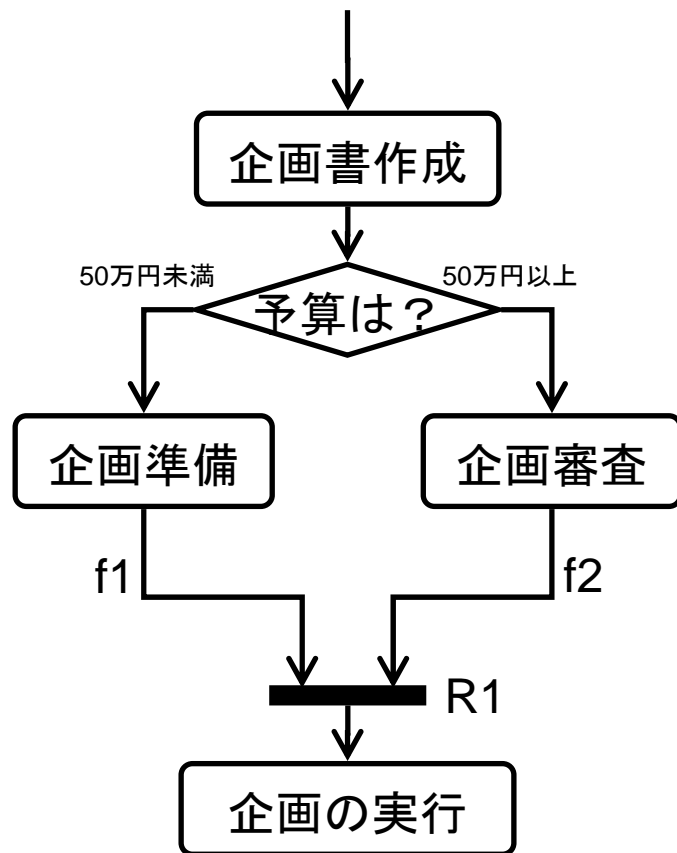
Example of a workflow



- Workflow (Nodes, Arcs):
 - a simple connected directed graph
 - Nodes = Starts \cup Ends \cup Activities
 - \cup XOR-splies \cup XOR-joins
 - \cup AND-splies \cup AND-joins
 - For each $n \in \text{Nodes}$ there exists a path from a start to n .
 - For each $n \in \text{Nodes}$ there exists a path from n to an end.
 - In this talk, we consider only acyclic workflows.

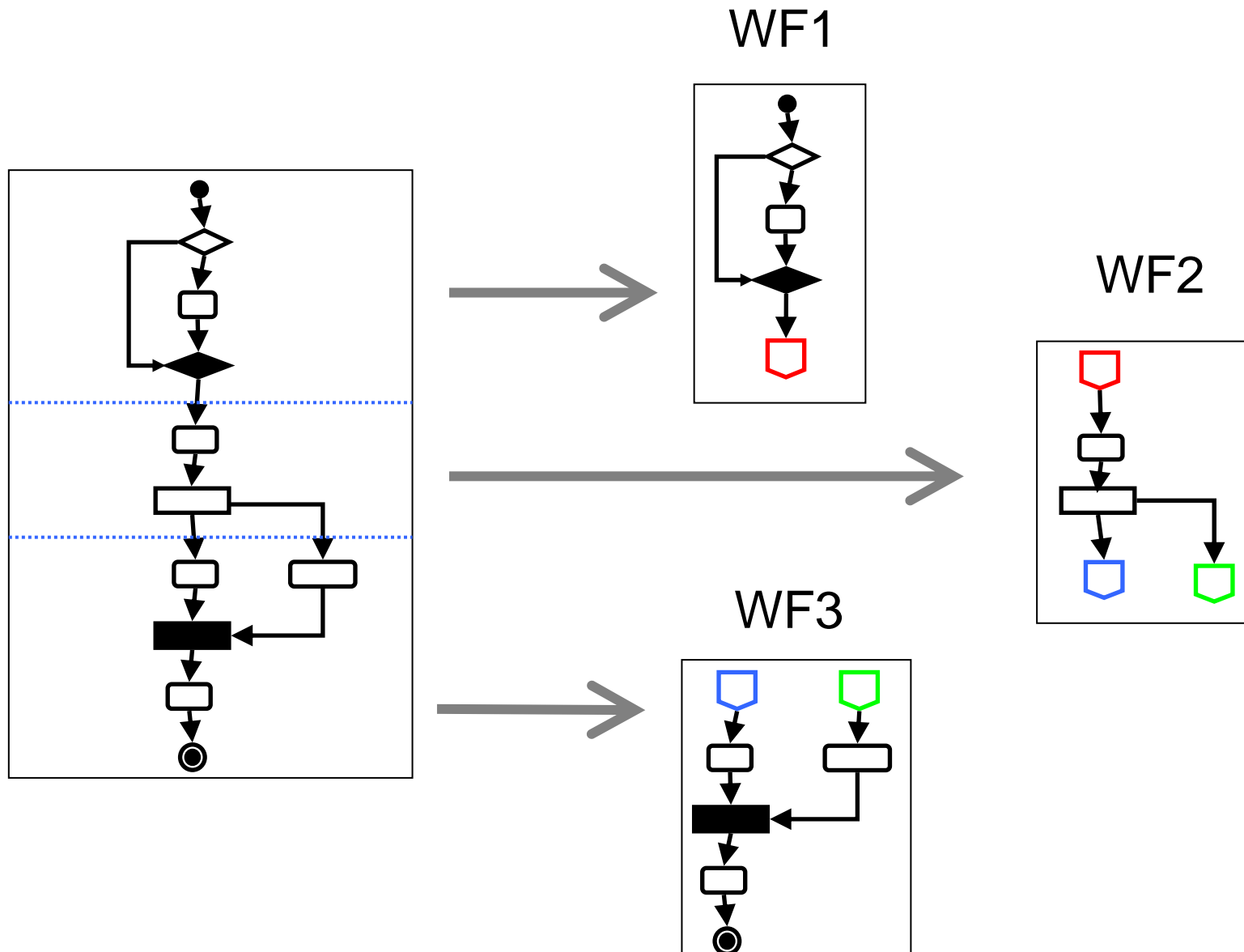


- Correctness of Workflow with one start and one end [Sadiq & Owlrowska 00]
 - Deadlock free
 - Lack of synchronization free

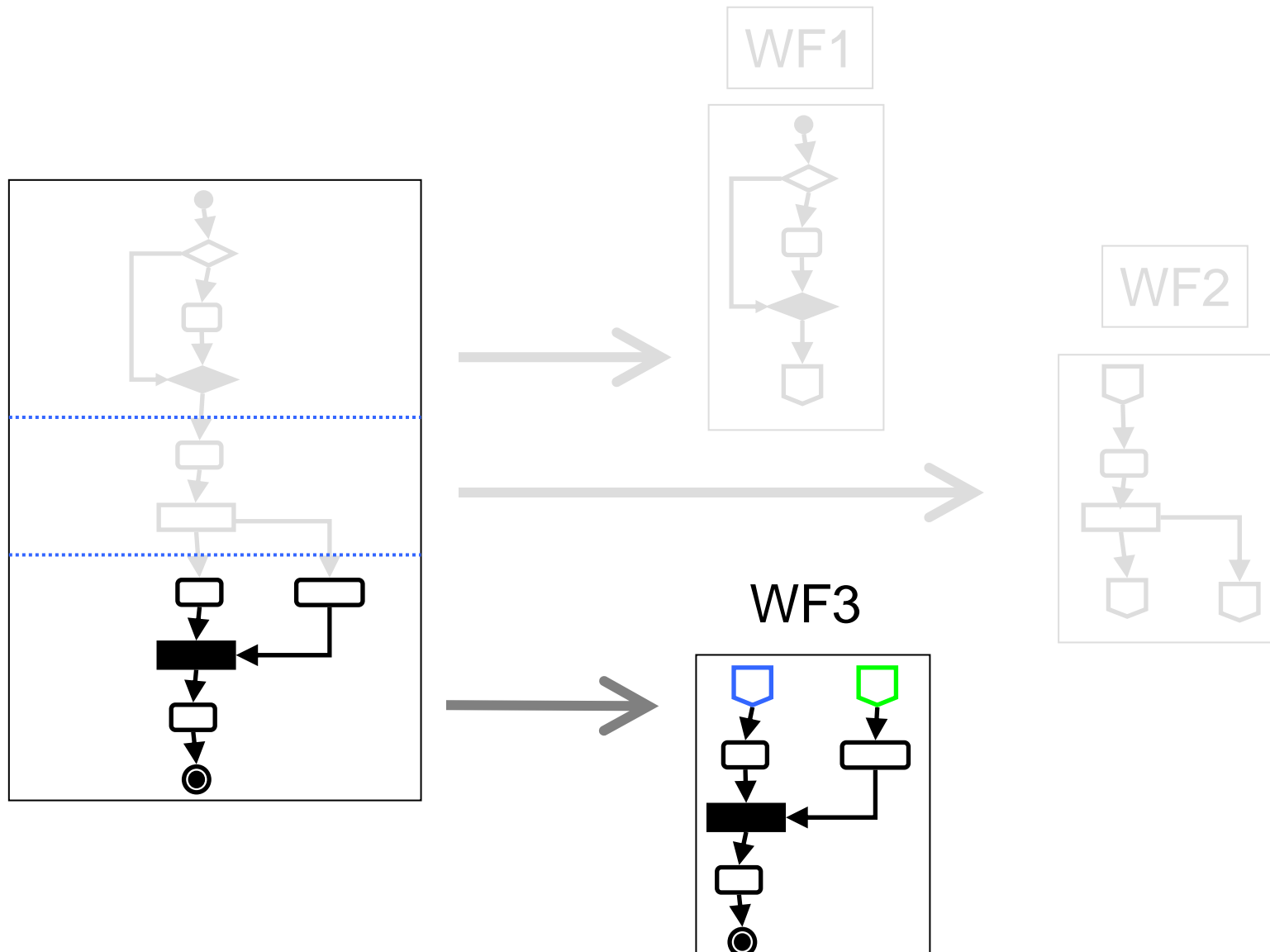


- Verification of Correctness
 - Graph Reductions
(Sadiq and Orłowska 1998, 2000)
 - WF-nets (van der Aalst 1997, 1998)
 - Global-Local Correctness
(Kindler, Martens and Reisig 2000)
 - Woflan (Verbeek, Basten and van der Aalst 2001)
 - Improvement of Sadiq-Orłowska's works
(Lin, Zhao, Li and Chen 2002)
 - Standard Workflow Models (Kiepuszewski, ter Hofstede and van der Aalst 2003)
 - EPCs (Dehnert and van der Aalst 2004~2006)

Why do we consider Multiple starts/ends?



Why do we consider Multiple starts/ends?



- Our purpose

- Extend the concept of correctness over workflows with one start and end to that over workflows with multiple starts and/or end.
- Verification Algorithms of the extended correctness of given workflows.
- Improvement of the algorithms to develop design assistant system of workflow.

- Our purpose
 - **Extend the concept of correctness over workflows with one start and end to that over workflows with multiple starts and/or end.**
 - Verification Algorithms of the extended correctness of given workflows.
 - Improvement of the algorithms to develop design assistant system of workflow.

- **General Correctness**

- We define “general correctness”, that is a generalized version of correctness of workflow to satisfies the following properties.

1. General correctness is a natural extension of correctness, that is, for a workflow with one start, general correctness is the same as original one.
2. General correctness is preserved by the operation of connection and/or division of workflows.
3. General correctness assures the possibility for a workflow to be completed to a correct workflow.

- **start**(W): the set of starts in W
- **end**(W): the set of ends in W

Definition 0.1 For a workflow W , an intermediary graph of W denotes the minimal subgraph V of W that satisfies the following properties.

1. V contains just one start of W .
 2. If V contains an XOR-split c , then V contains just one out-degree of c .
 3. If V contains a node c other than XOR-split, then V contains all out-degrees.
- **IG**(W) denotes the set of intermediary graph of W .
 - **IG**(W, s) the set of intermediary graph of W with start s .

Definition 0.2 (Sadiq and Orłowska 2000) Let W be a workflow with one start.

- An intermediary graph V of W is said to be deadlock free if, for every AND-join r in V , V contains all in-degrees of m .
- An intermediary graph V of W is said to be lack of synchronization free if, for every XOR-join m in V , V contains just one in-degree of m .

Definition 0.3 (Sadiq and Orłowska 2000) A workflow W with one start is said to be correct if every intermediary graph V of W is deadlock free and lack of synchronization free.

Definition 0.4 For a workflow W , a trace graph of W denotes a non-empty subgraph V of W that satisfies the following properties. Let n be a node in V .

1. If n is an XOR-split, then V contains just one out-degree of n as well as the in-degree of n .
2. If n is an XOR-join, then V contains just one in-degree of n as well as the out-degree of n .
3. Otherwise, V contains all in-degrees and all out-degrees of n .

- $\mathbf{TG}(W)$: the set of trace graphs of W
- $\mathbf{TG}(W, S)$: the set of trace graphs V of W with $\mathbf{start}(V) = S$

Definition 0.5 For a workflow W and $U_1, U_2 \in \mathbf{IG}(W)$, U_1 and U_2 are said to conflict on an XOR-split c if U_1 and U_2 share c but they do not share any out-degree of c .

Definition 0.6 Let W be a workflow, \mathbf{U} a set of intermediary graphs of W and n an XOR-split. Then, \mathbf{U} is said to conflict on n there exists a pair (U_i, U_j) on \mathbf{U} that conflicts on n .

Definition 0.7 Let W be a workflow and $W_1, \dots, W_n \in \mathbf{TG}(W)$ with $W_i \cap W_j = \emptyset$ for each $i \neq j$. Then, the non-connected graph $W_1 \cup \dots \cup W_n$ is called a summation of trace graphs.

$\mathbf{TG}_s(W, S)$: the set of summations $W_1 \cup \dots \cup W_n$ with

$$S = \mathbf{end}(W_1) \cup \dots \cup \mathbf{end}(W_n).$$

Definition 0.8 Let W be a workflow and S a set of starts s_1, \dots, s_n of W . Then, S is called an import of W if, for every set consisting of $U_i \in \mathbf{IG}(W, s_i)$ ($i = 1, \dots, n$) that is not conflict on any XOR-split in W , there exists a summation \mathbb{V} with $\mathbb{V} = U_1 \cup \dots \cup U_n$.

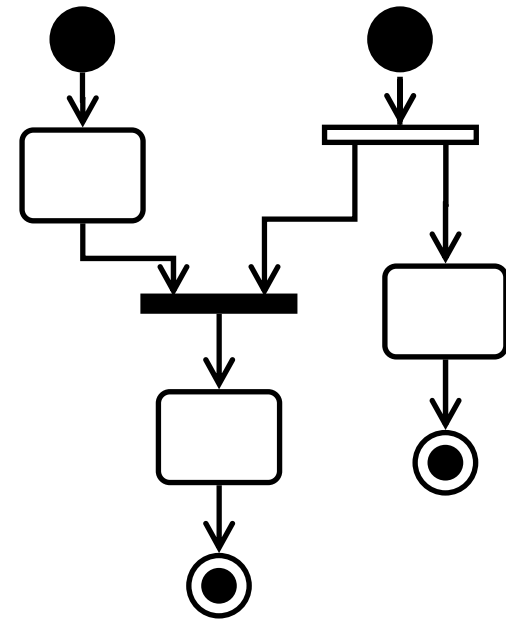
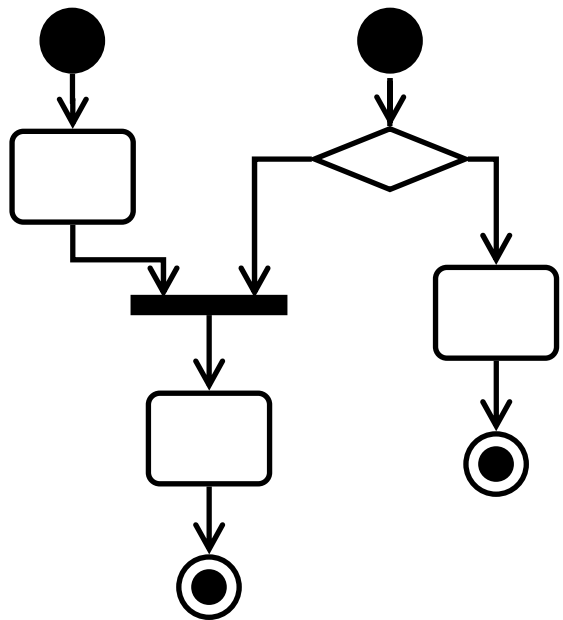
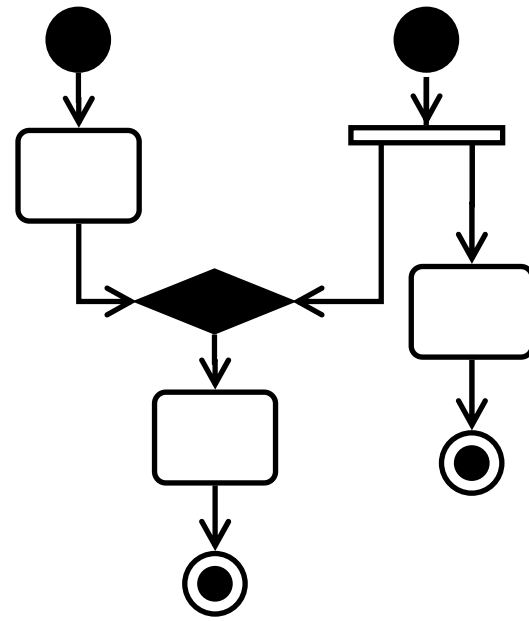
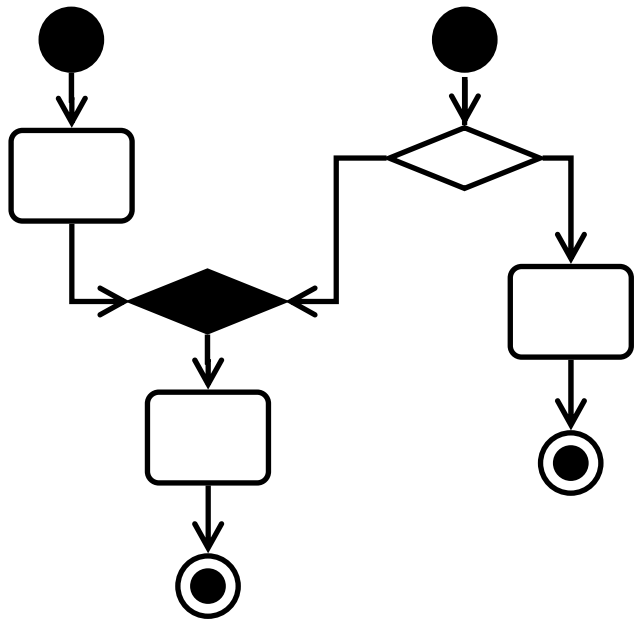
Definition 0.9 A workflow W is said to be generally correct if W has a set \mathbb{I} of imports of W with $\bigcup_{I \in \mathbb{I}} I = \mathbf{start}(W)$.

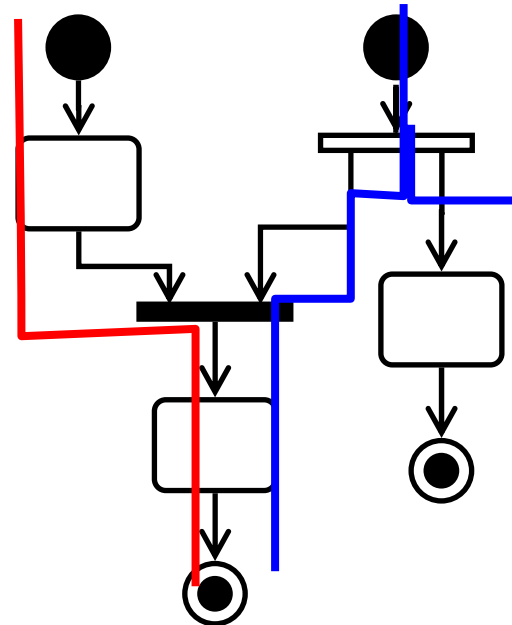
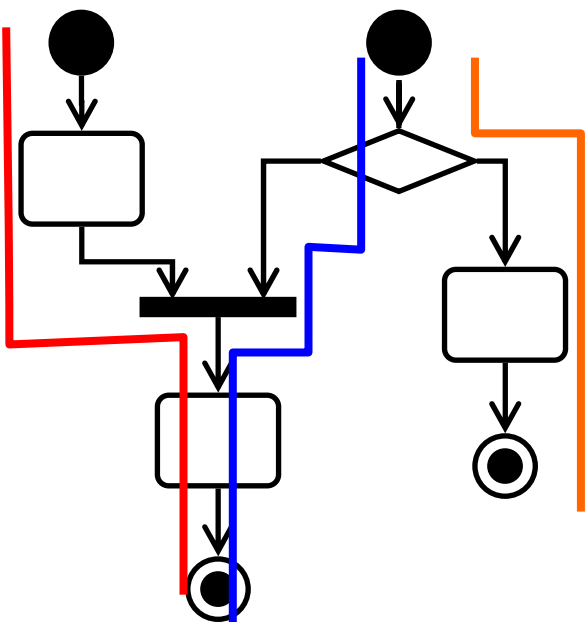
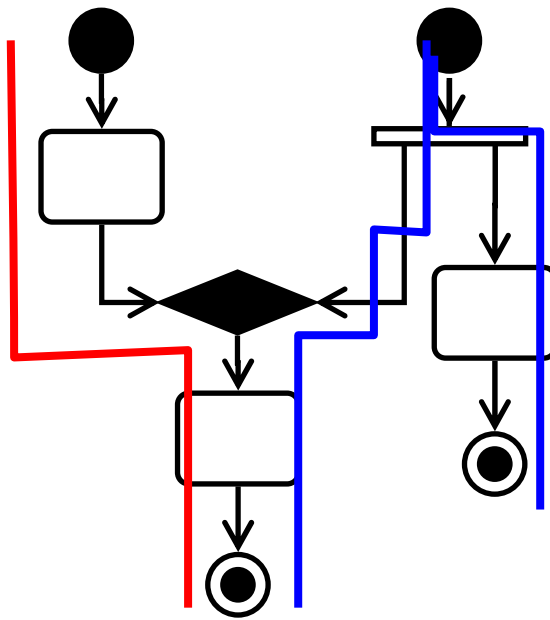
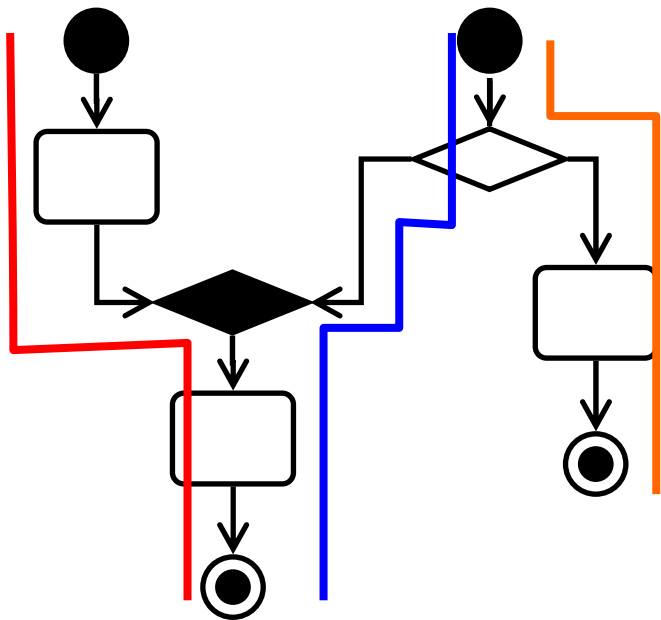
Definition 0.10 Let W be a workflow.

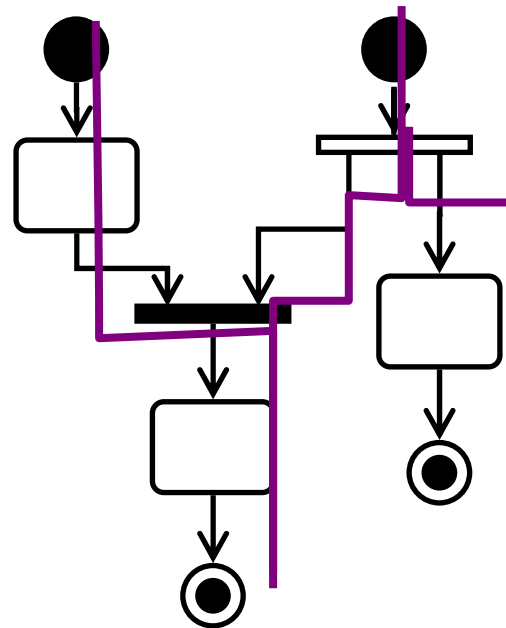
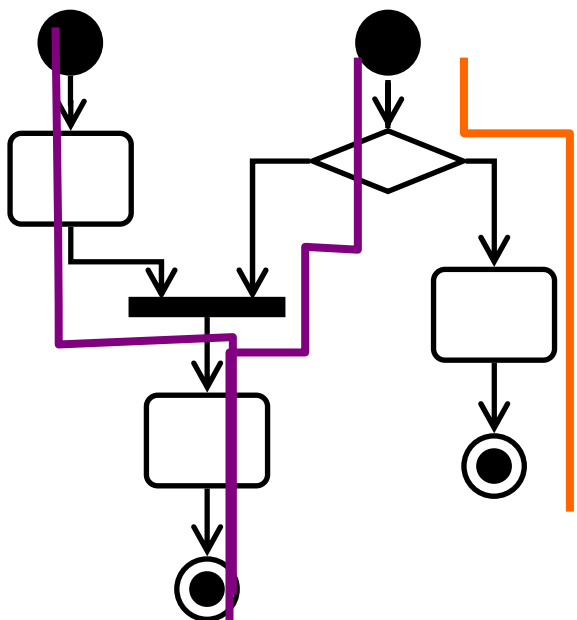
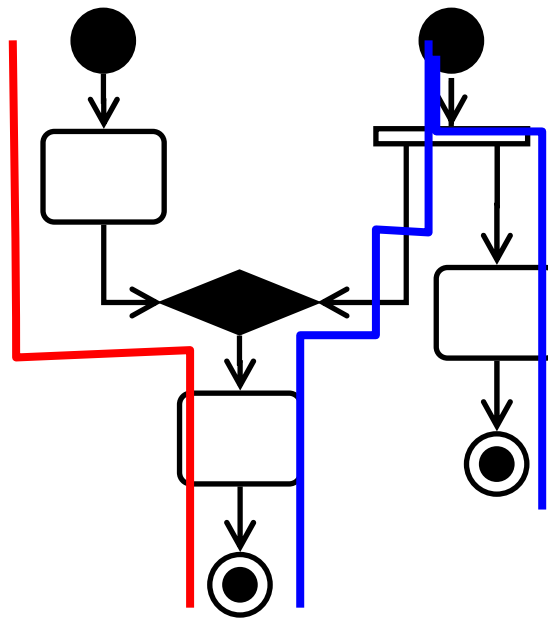
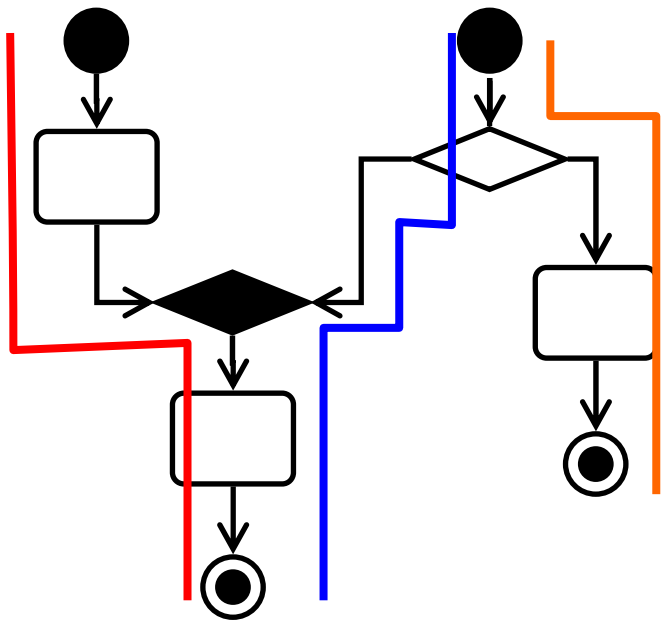
(1) For a summation \mathbb{V} of trace graphs V_1, \dots, V_n of W , the export $\mathbf{ex}(\mathbb{V})$ of \mathbb{V} denotes $\mathbf{end}(V_1) \cup \dots \cup \mathbf{end}(V_n)$.

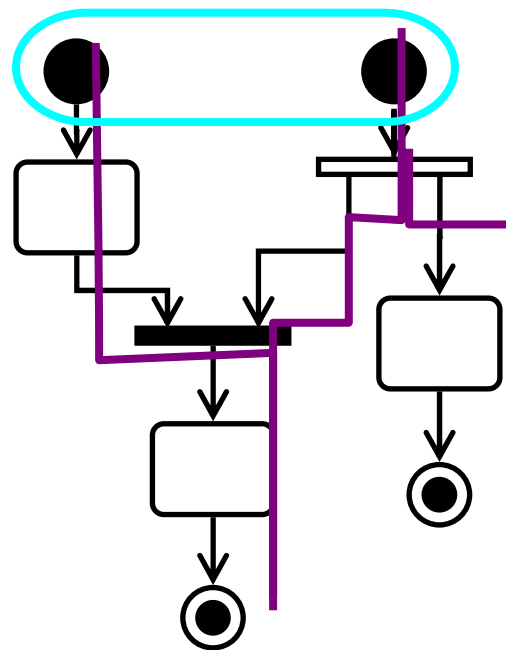
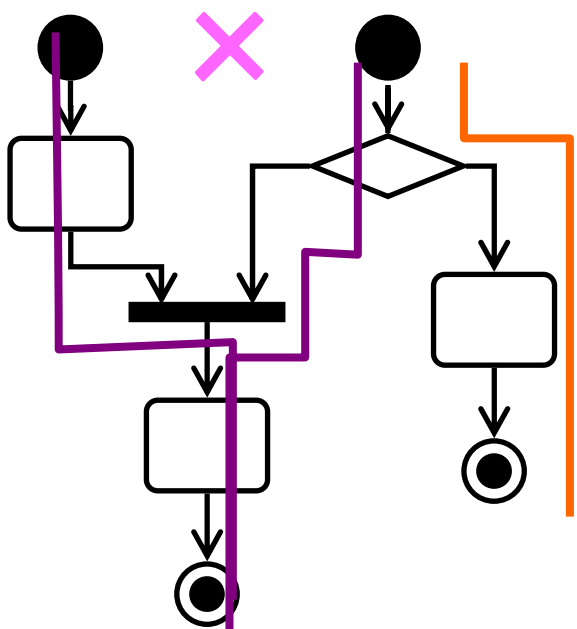
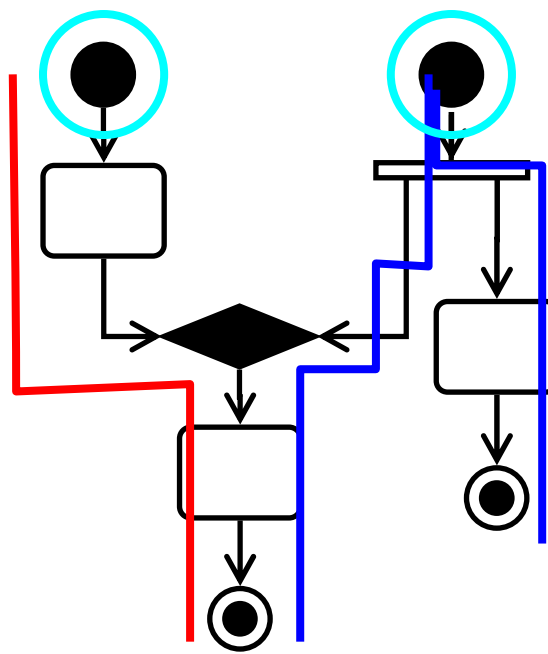
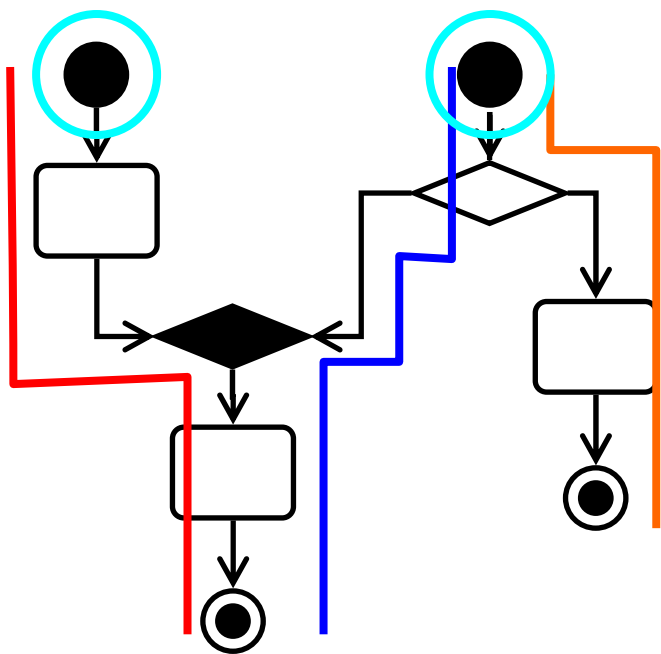
(2) For an import I of W , the set $\{\mathbf{ex}(\mathbb{V}) \mid \mathbb{V} \in \mathbf{TG}_s(W, I)\}$ is called by the export family of I and denoted by $\mathbb{E}_s(W, I)$.

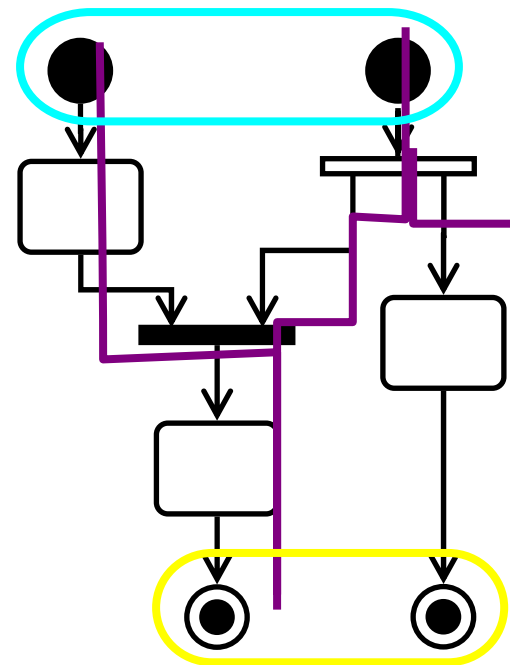
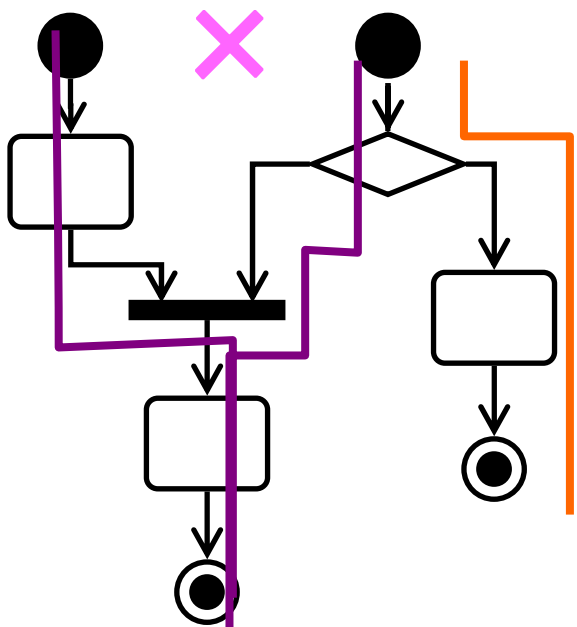
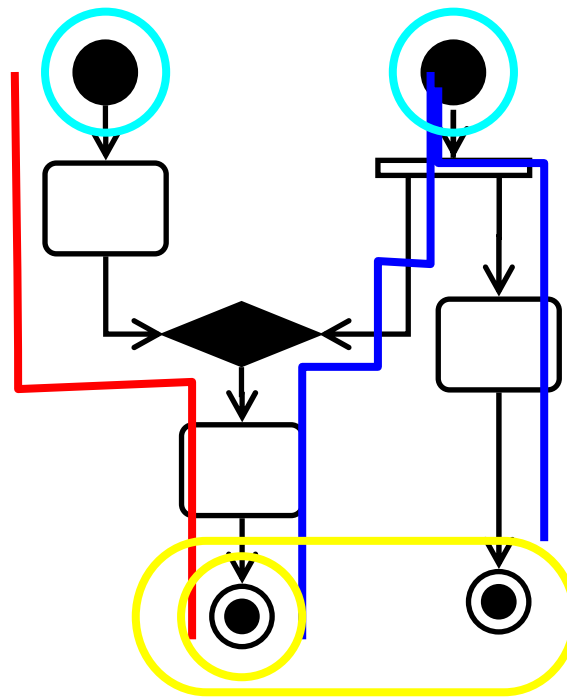
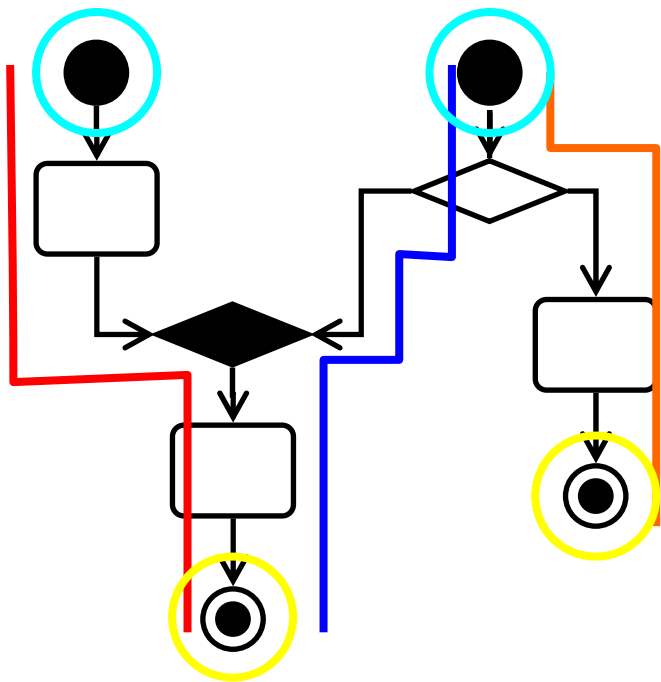
(3) For an import family \mathbb{I} of W , the set $\bigcup_{I \in \mathbb{I}} \mathbb{E}_s(W, I)$ is called by the export family of \mathbb{I} and denoted by $\mathbb{E}_s^*(W, \mathbb{I})$.











- $\mathbf{WF}(n, m)$: the set of workflows with n starts and m ends
- $\mathbf{WF} := \bigcup_{n,m} \mathbf{WF}(n, m)$.

Definition 0.14 Let $W_1 \in \mathbf{WF}(n, m)$, $W_2 \in \mathbf{WF}(m, l)$ and f a bijection from $\mathbf{end}(W_1)$ to $\mathbf{start}(W_2)$. Then, $W_1 *_f W_2$ denotes the workflow obtained from W_1 and W_2 by executing the following procedures.

1. Remove all ends of W_1 and their in-degrees.
2. Remove all starts in W_1 and their out-degrees.
3. For the source n of the in-degree of each end e in W_1 and the target n' of the out-degree of each start $f(e)$ in W_2 , add the arc from n to n' .

In the remainder of this paper, we omit “ f ” in $W_1 *_f W_2$ and identify each $e \in \mathbf{end}(W_1)$ with $f(e) \in \mathbf{start}(W_2)$.

Theorem 0.15 Let $W_1 \in \mathbf{WF}(n, m)$ and $W_2 \in \mathbf{WF}(m, l)$. Then, $W_1 * W_2$ is generally correct for an import family \mathbb{I} if and only if

- $W_1 \in \mathbf{WF}(n, m)$ is generally correct for \mathbb{I} , and
- $W_2 \in \mathbf{WF}(m, l)$ is generally correct for $\mathbb{E}_s^*(W, \mathbb{I})$.

Definition 0.16 For a workflow $W \in \mathbf{WF}(n, m)$, W is said to be extendible if there exists a workflow $W_0 \in \mathbf{WF}(1, n)$ such that $W_0 * W$ is correct.

Lemma 0.17 For every finite set S with $\#S = n > 0$ and every subset \mathbb{S} of the power set of S , there exists a (generally) correct workflow $W \in WF(1, n)$ with $\mathbb{S} = \mathbb{E}_s^*(W, \{\mathbf{start}(W)\})$.

Corollary 0.18 Let $W \in \mathbf{WF}(n, m)$ be a generally correct workflow. Then, W is extendible if and only if W is generally correct.

- **Conclusion**

- We extend the concept of correctness over workflows with one start and end to that over workflows with multiple starts and/or end.
- 1. General correctness is a natural extension of correctness, that is, for a workflow with one start, general correctness is the same as original one. (Theorem 0.13)
- 2. General correctness is preserved by the operation of connection and/or division of workflows. (Theorem 0.15)
- 3. General correctness assures the possibility for a workflow to be completed to a correct workflow. (Corollary 0.18)