

Title	Model Checking Web Specifications : Verification of design specifications for Web applications
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Model Checking Web Specifications

- Verification of design specifications
for Web applications -

JAI ST/AIST Workshop Sep. 21 2005

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Outline of My Talk

- Background
- Proposed methods
- Experiments
- Conclusion

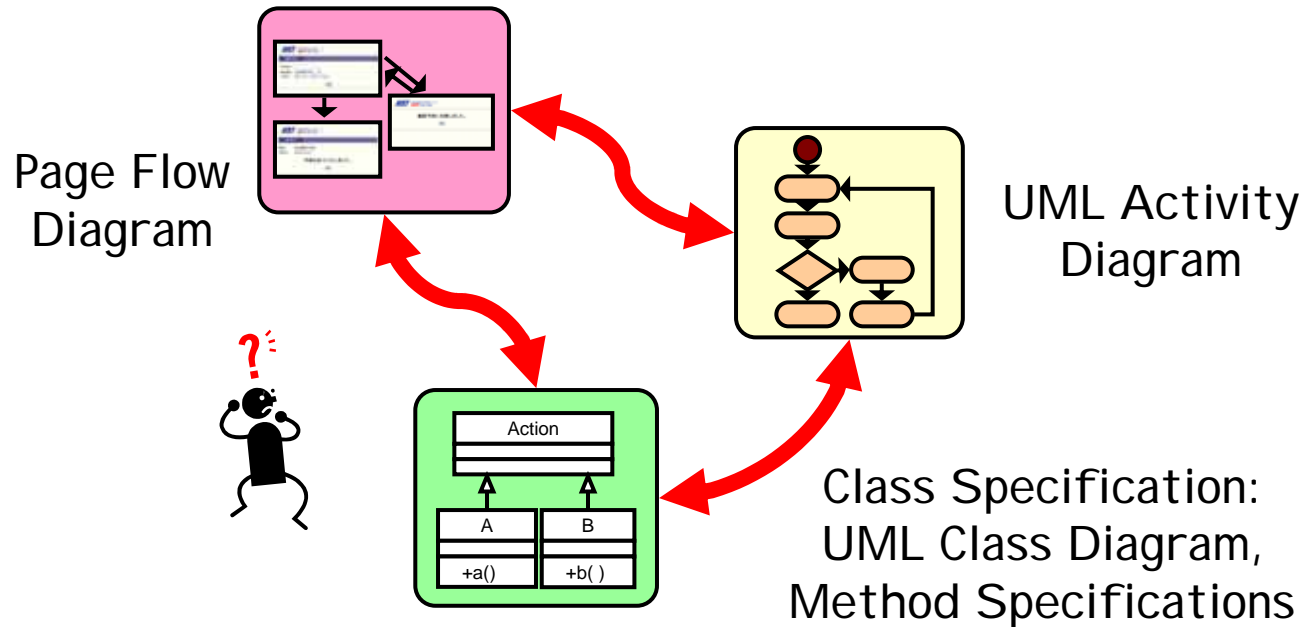
Background

- A certain company, A, provided us a set of design specifications which was used for an actual Web-based business processing application.
- In our fieldwork, we tackled proposing a verification technique for the given design specifications that is easily applicable to existing design process for Web applications.



Design Specifications for a Web Application

*Preliminary
Design
Phase*

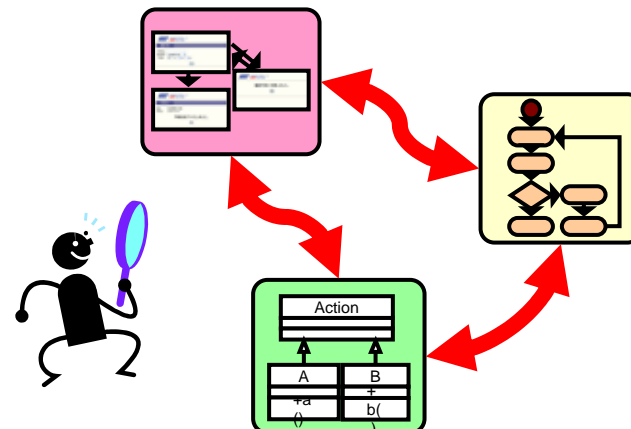


*Detailed
Design
Phase*

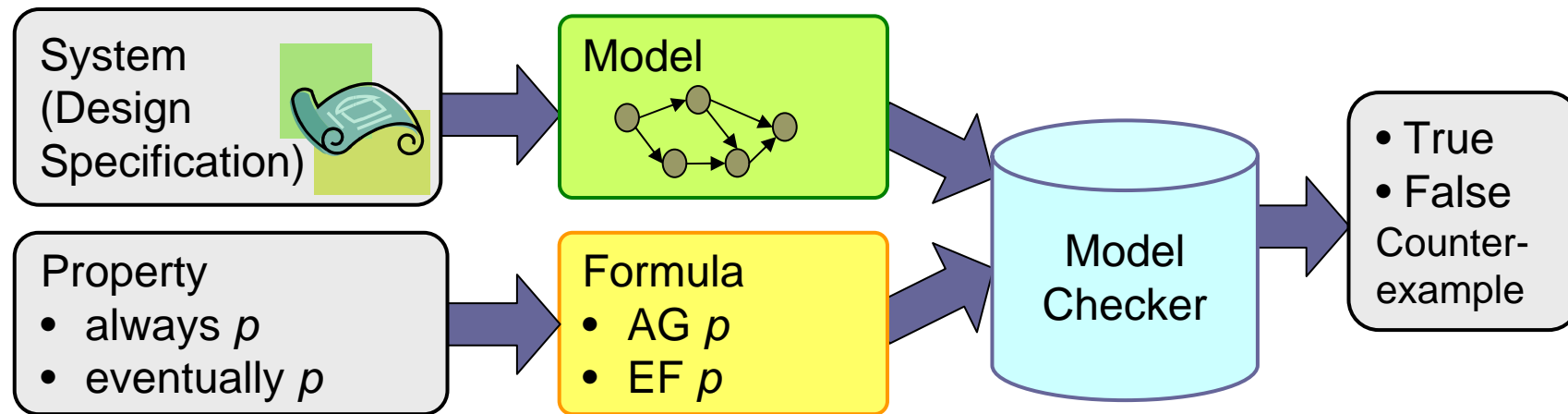
- Consistency checking for design specifications is important in terms of not only reliability but also maintenance and reuse of a Web application.

Our Work

- We proposed verification methods to check
 - I. Consistency between a page flow diagram & an activity diagram
 - II. Consistency between a page flow diagram & a class specification
 - III. Consistency between a class specification & an activity diagram
- The proposed methods are based on a model checking technique.



Model Checking

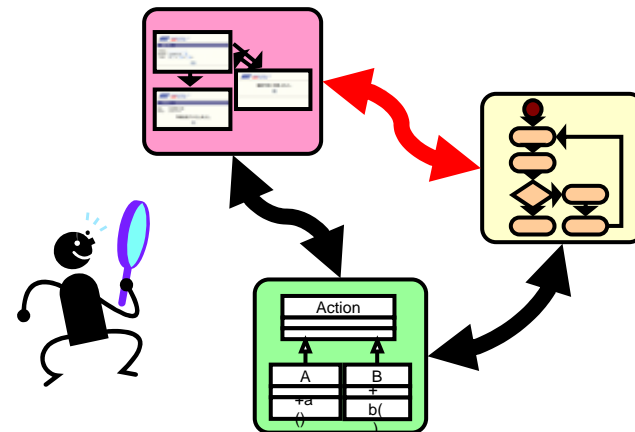


- Model checking is a verification technique that can exhaustively check whether a finite transition system satisfies a temporal logic formula or not.
- Model checking is also helpful to allocate errors because, when the system does not satisfy the property, counterexample is output with the result.

Outline of My Talk

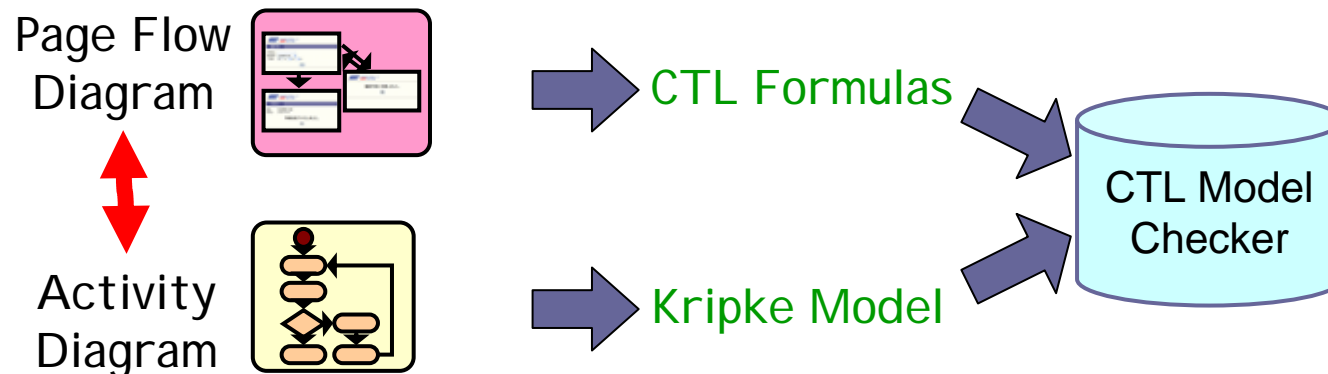
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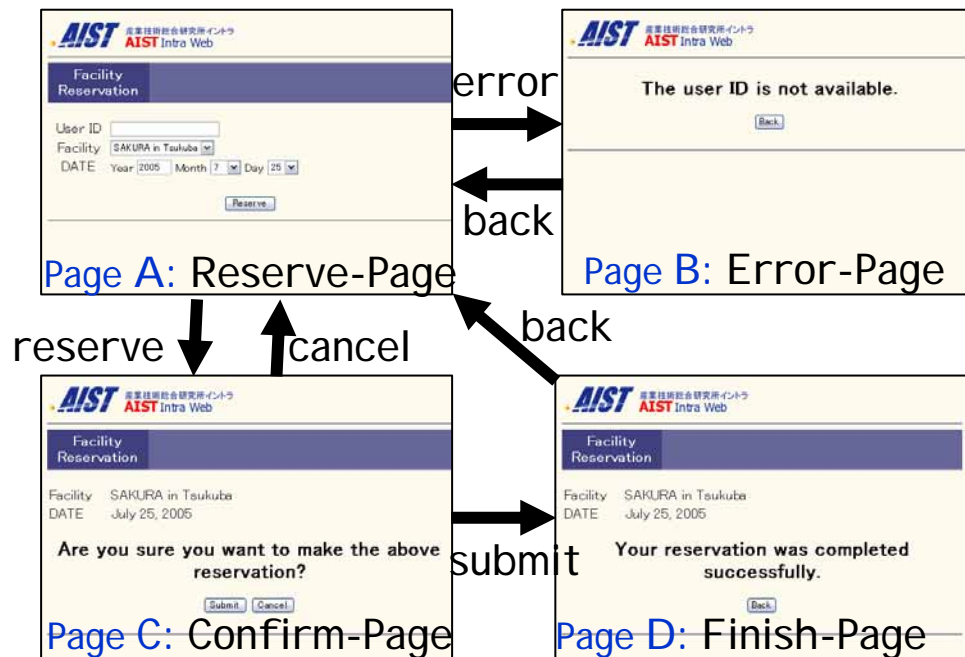


Outline of the proposed method I

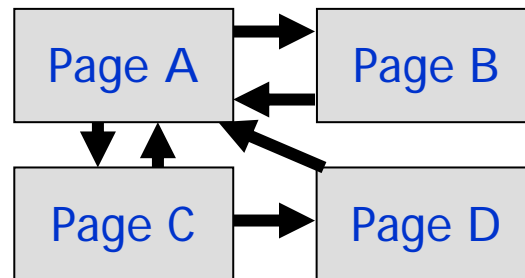
- Proposed method to verify the consistency between a page flow diagram and an activity diagram
 1. Define the consistency between a page flow diagram and an activity diagram.
 2. Represent the consistency using CTL formulas, which are generated from the page flow diagram.
 3. Model check if the CTL formulas hold for a Kripke model constructed from the activity diagram.



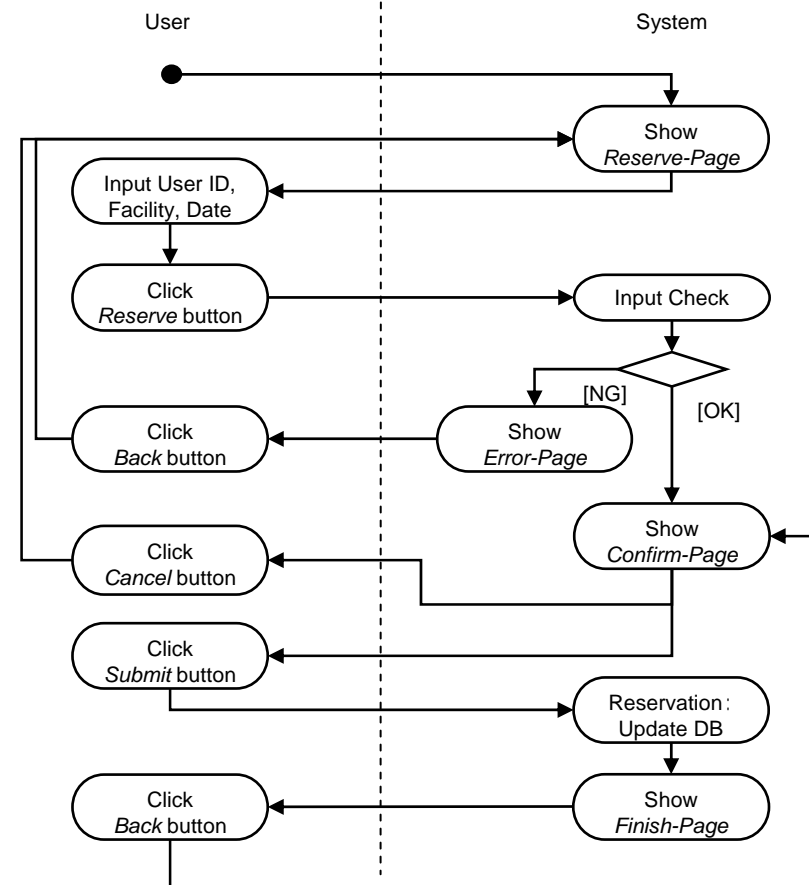
Example Application and Specifications



Page Flow Diagram

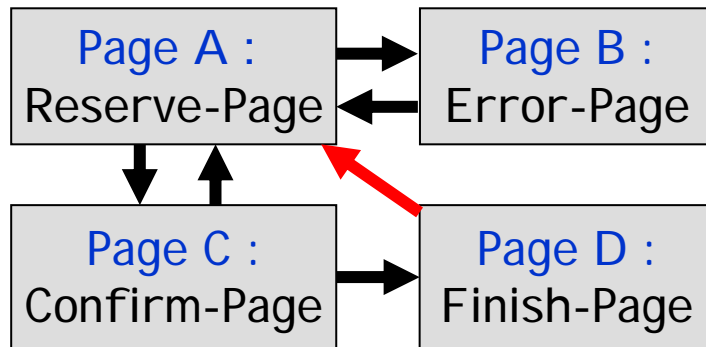


Activity Diagram

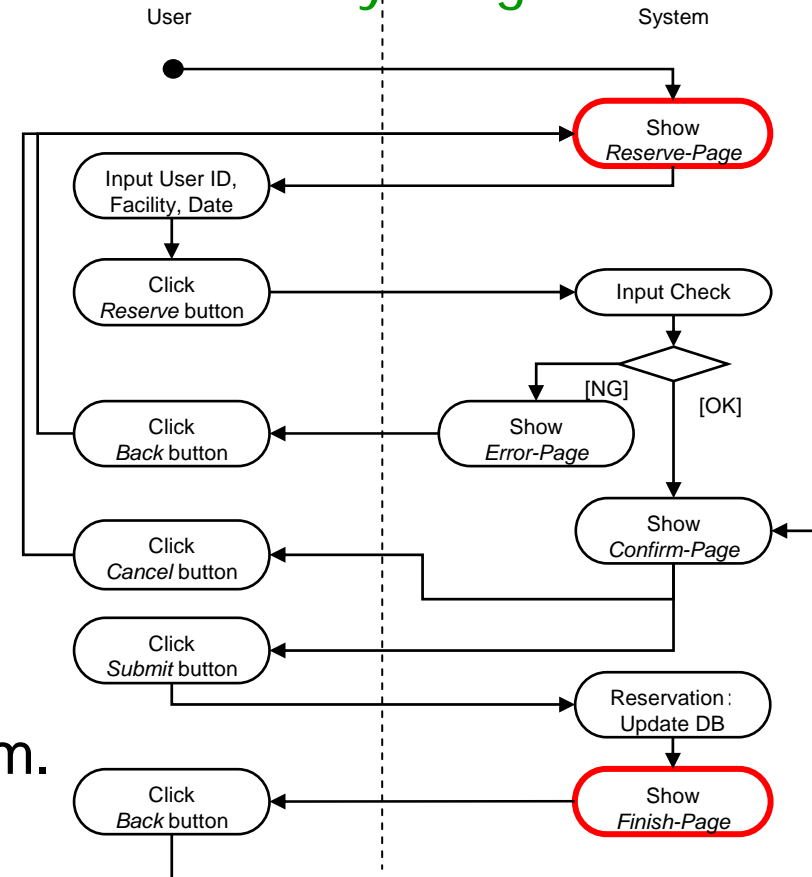


Inconsistency 1

Page Flow Diagram



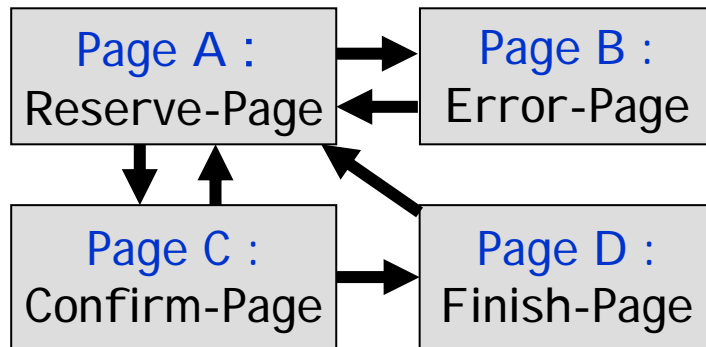
Activity Diagram



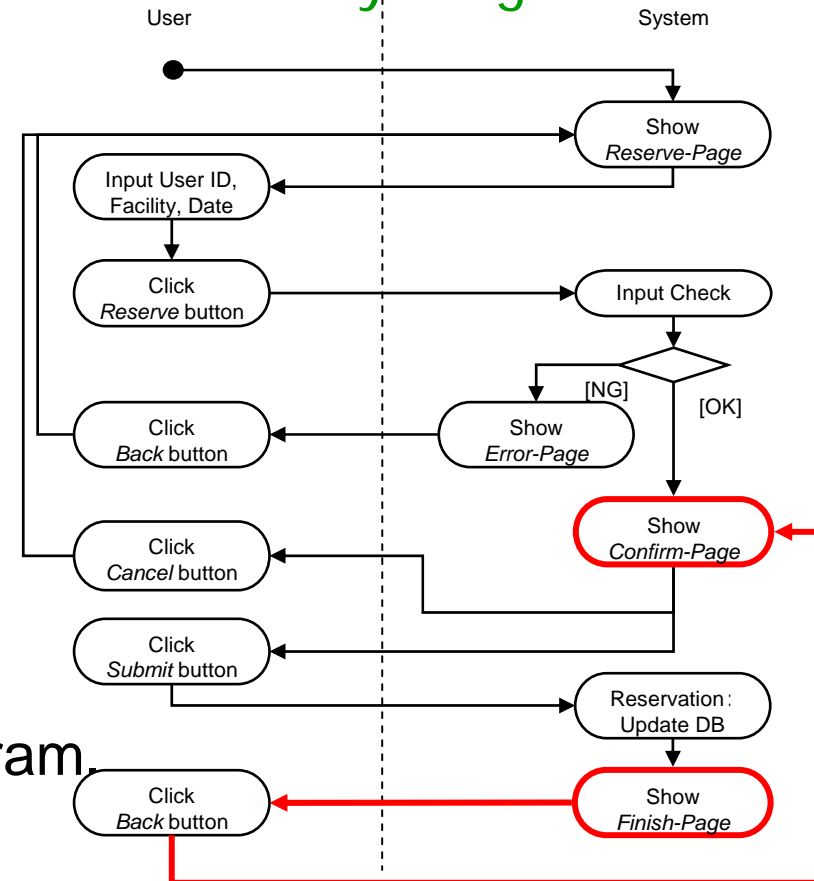
A page transition (page D, page A) in the page flow diagram does not occur in the activity diagram.

Inconsistency 2

Page Flow Diagram



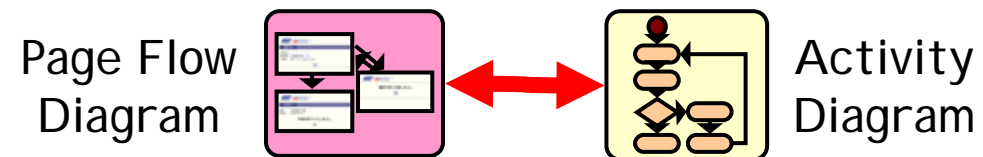
Activity Diagram



A page transition (page D, page C) in the activity diagram does not exist in the page flow diagram.

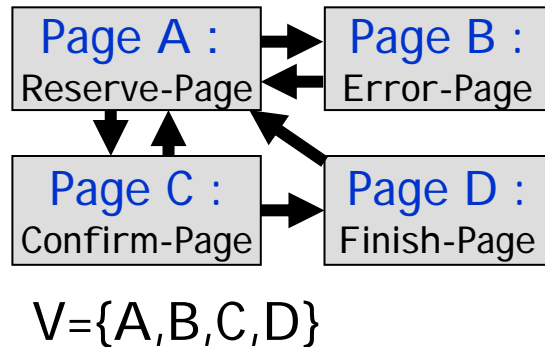
Consistency of Page Flow Diagram and Activity Diagram

- We employ the following two conditions for a definition of the consistency between a page flow diagram and an activity diagram:
 - C1** : For each page transition in the page flow diagram, a transition corresponding to the page transition exists in the activity diagram.
 - C2** : Every transition in the activity diagram corresponds to a stuttering or a page transition in the page flow diagram.

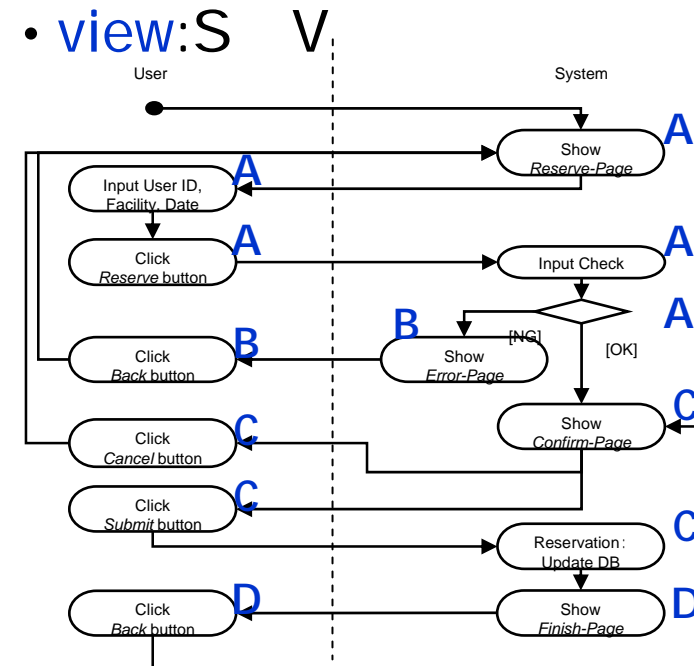


Definition of Consistency

Page Flow Diagram: $(V, E (\subseteq V \times V))$



Activity Diagram: $(S, T (\subseteq S \times S))$



Def. Consistency between (V, E) and $(S, T, view)$ holds iff

C1 : For each $(x, x') \in E$, there exists $(s, s') \in T$ such that $view(s) = x$ and $view(s') = x'$.

C2 : For each $(s, s') \in T$, $view(s) = view(s')$ or $(view(s), view(s')) \in E$.

Consistency Checking Problem

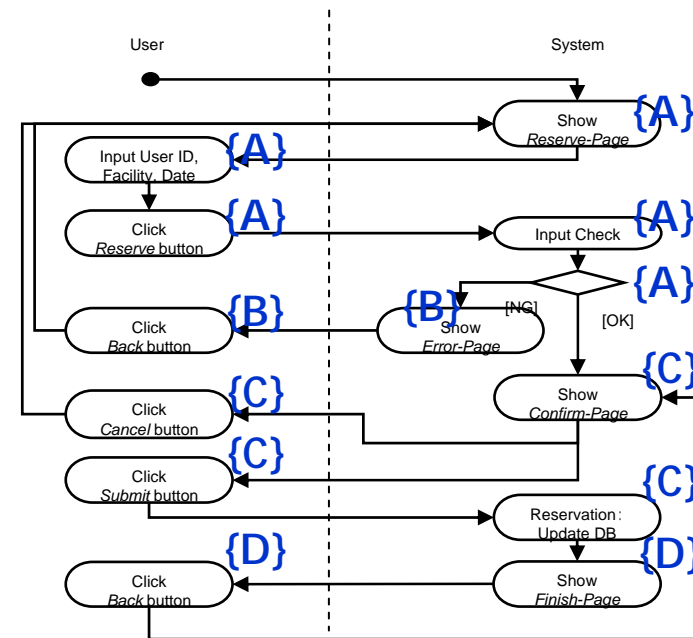
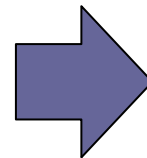
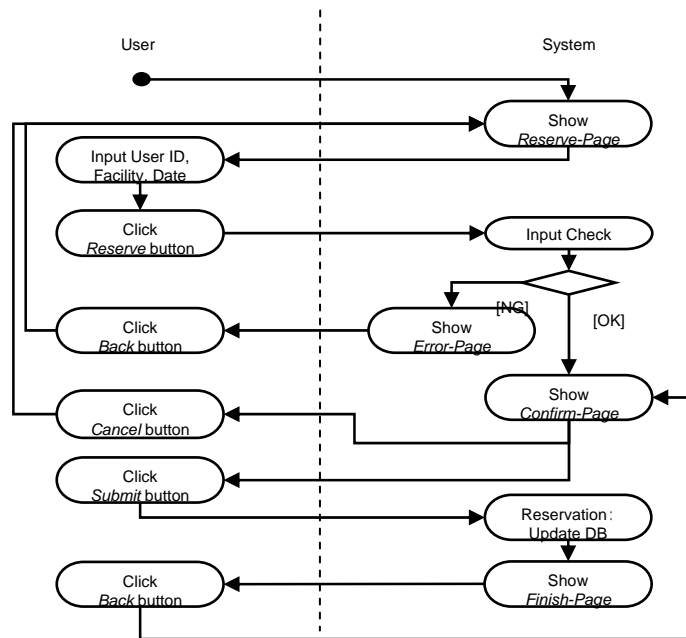
- Consider a page flow diagram (V, E) and a Kripke structure $K = (S, T, s. \{view(s)\}: S \rightarrow 2^V)$ where (S, T) denotes an activity diagram. Let s_0 denote the initial state of K .
- The two conditions for the consistency are represented using CTL (Computation Tree Logic) as follows:
 - C1** : $\forall (x, x') \in E, (K, s_0) \models EF (x \neq EX x')$
 - C2** : $\forall x \in V, (K, s_0) \models AG (x \neq AX (\bigvee_{(x, x') \in E} EX x'))$
- The consistency between a page flow diagram and an activity diagram is verified by model checking the above CTL formulas for Kripke structure K .

Model input to Model Checker

- Input model is constructed from an activity diagram.

Activity Diagram: (S,T)

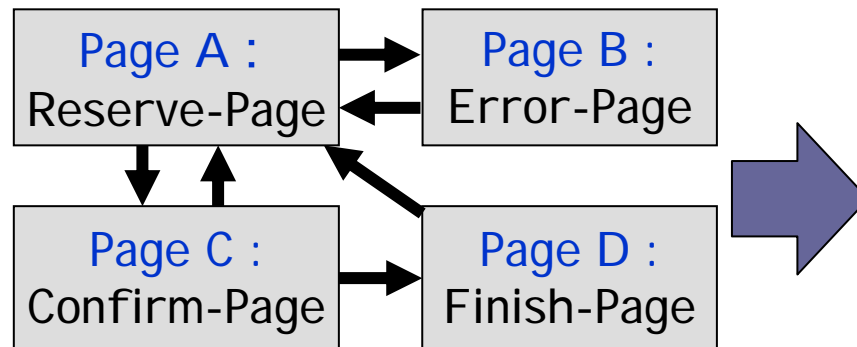
Kripke Model K: (S,T, s.{view(s)})



Formulas input to Model Checker

- Input formulas are generated from a page flow diagram.

Page Flow Diagram

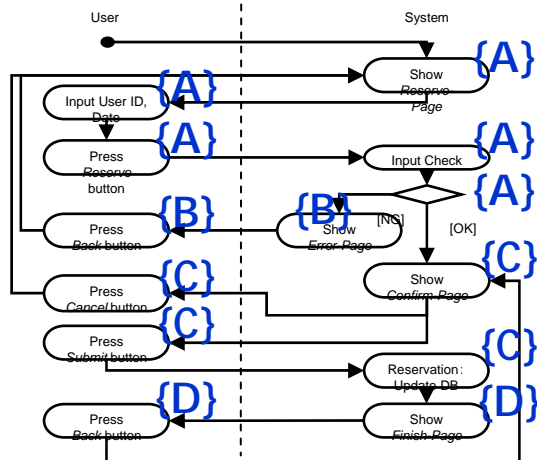


CTL Formulas

- 1. EF (A EX B)
 - 2. EF (A EX C)
 - 3. EF (B EX A)
 - 4. EF (C EX A)
 - 5. EF (C EX D)
 - 6. EF (D EX A)
 - 7. AG (A AX (A B C))
 - 8. AG (B AX (B A))
 - 9. AG (C AX (C A D))
 - 10. AG (D AX (D A))

Model Checking

Kripke Model K:
(S, T, s.{view(s)})



CTL Formulas

- *Cond1*

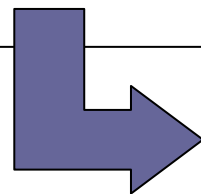
1. EF (A EX B)
2. EF (A EX C)
3. EF (B EX A)
4. EF (C EX A)
5. EF (C EX D)
6. EF (D EX A)

- *Cond2*

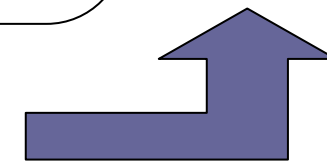
7. AG (A AX (A B C))
8. AG (B AX (B A))
9. AG (C AX (C A D))
10. AG (D AX (D A))

Result

-----	True
-----	True
-----	True
-----	True
-----	True
-----	False
-----	True
-----	True
-----	True
-----	False



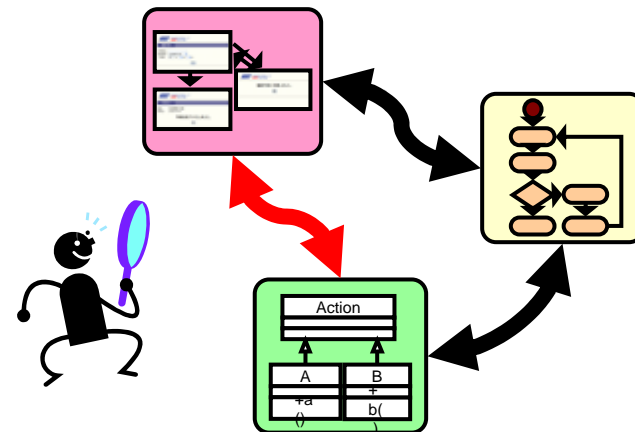
CTL Model Checker



Outline of My Talk

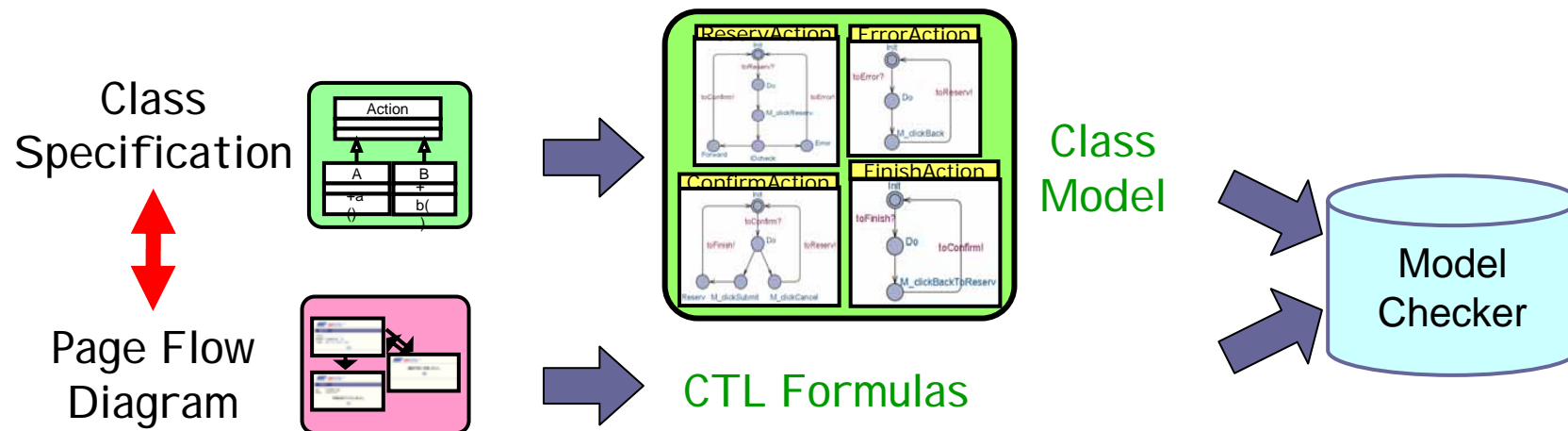
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Outline of the proposed method I I

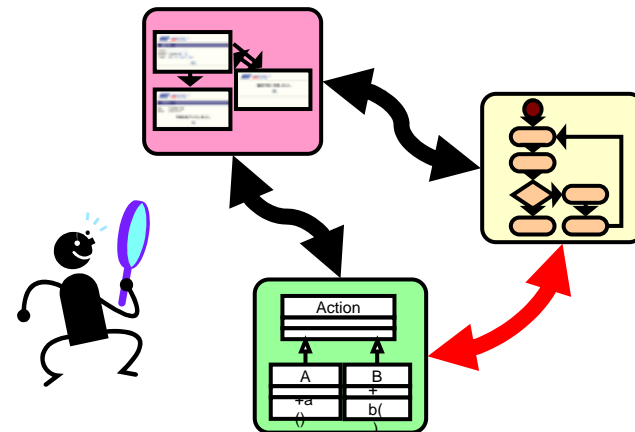
- Proposed method to verify the consistency between a page flow diagram and a class specification
 - From the given class specification consisting of a class diagram and method specifications, we model its behavior by a parallel composition of labeled transition systems.
 - Apply the proposed method I to the behavior model of the class specification.



Outline of My Talk

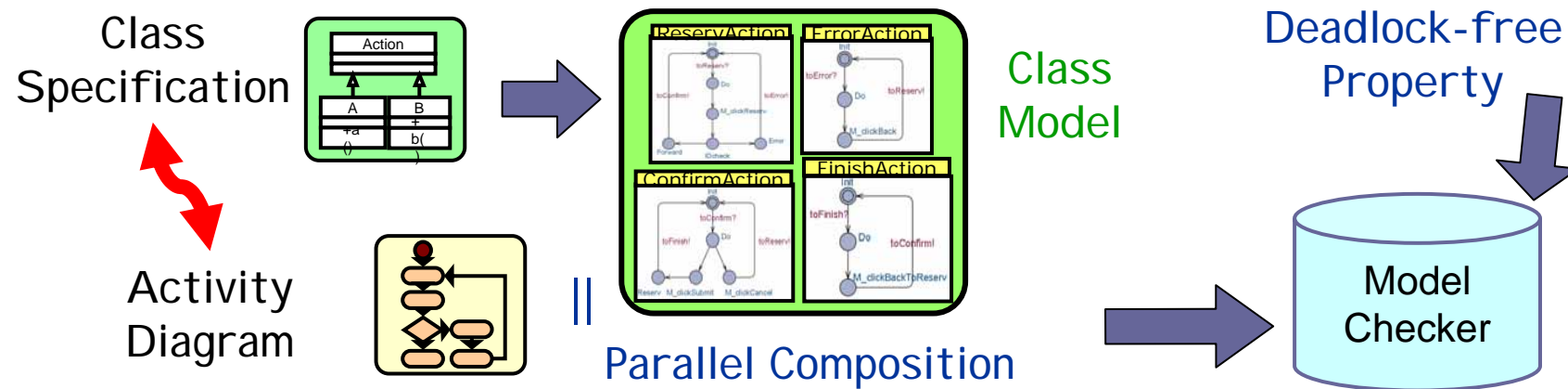
- Background
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 - III. **Consistency between a class specification & an activity diagram**

- Experiments
- Conclusion



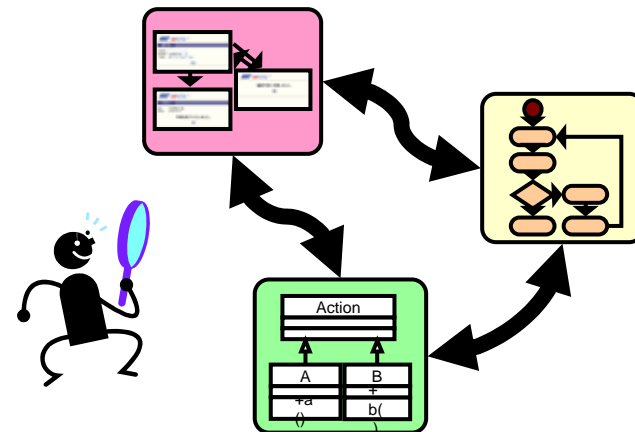
Outline of the proposed method I I I

- Proposed method to check the consistency between a class specification and an activity diagram
 1. Compose the class model and the activity diagram.
 2. Model check the deadlock-free property for the composed model.
- If a deadlock occurs in the composed model, there exists an inconsistency between the two specifications.



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- **Experiments**
- Conclusion

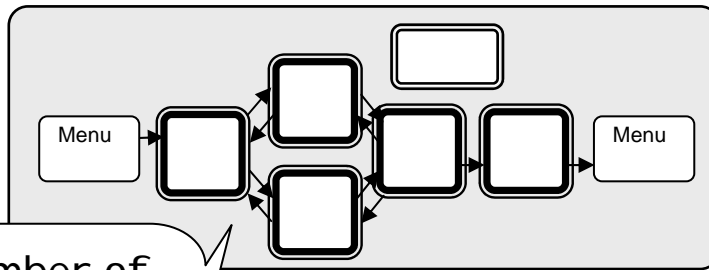


Case Study

- We applied the proposed methods to the real specifications of the given Web application.
 - Developed by Java using Jakarta Struts framework.
 - Classified into several tens of modules.
- We chose one module M and checked the consistencies for a page flow diagram, an activity diagram, and a class specification for module M .

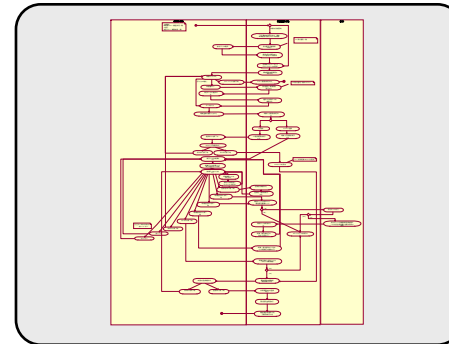
Experiment I

Page Flow Diagram



- Number of pages: 6
- Number of transitions: 9

Activity Diagram



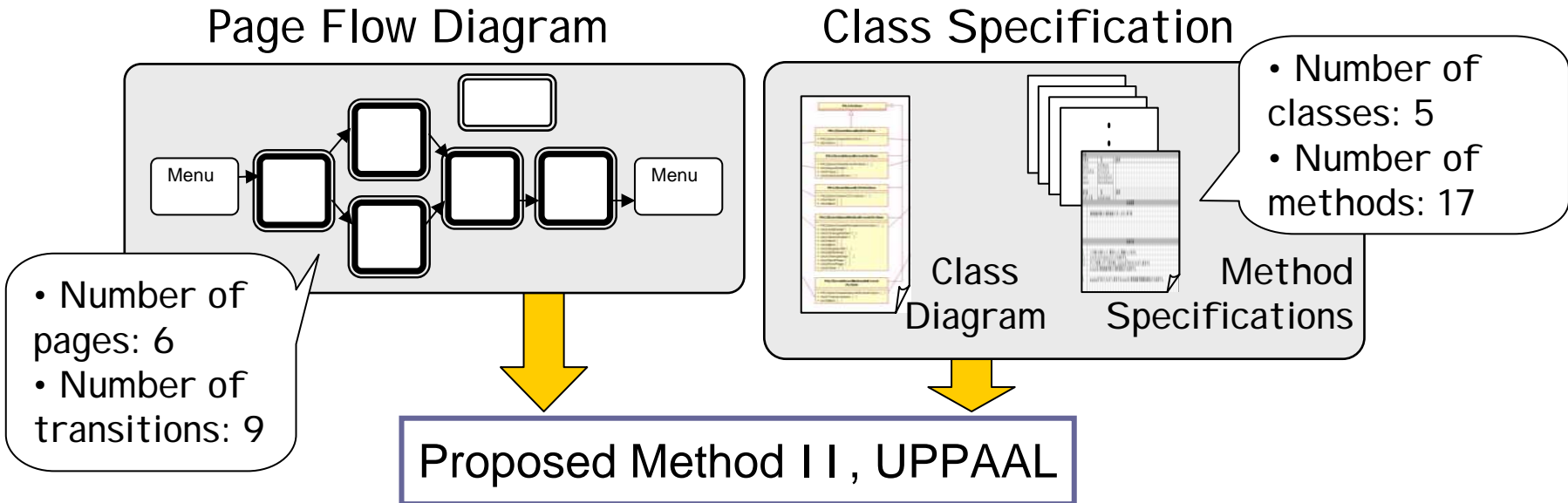
- Number of states: 66
- Number of transitions: 83

Proposed Method I, NuSMV

Proprietary Secret

X : page transition that does not exist in the activity diagram
 : page transition that exists only in the activity diagram

Experiment II

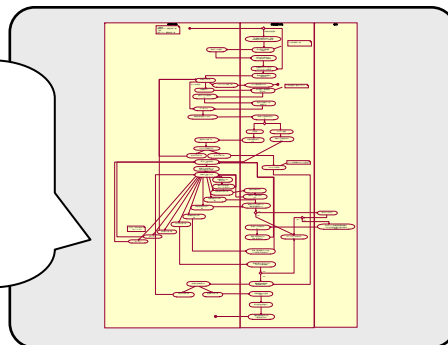


Proprietary Secret

X : page transition that does not exist in the class specification
 : page transition that exists only in the class specification

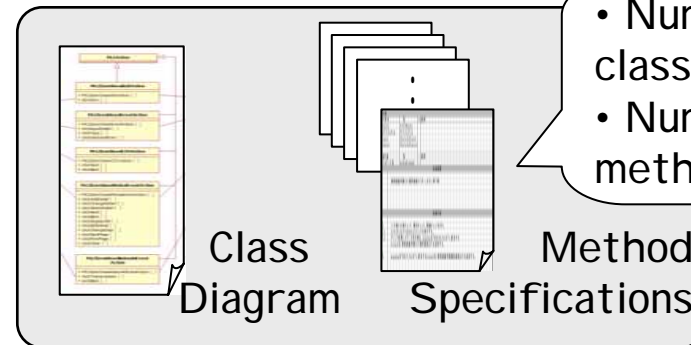
Experiment III

Activity Diagram



- Number of states: 66
- Number of transitions: 83

Class Specification



- Number of classes: 5
- Number of methods: 17

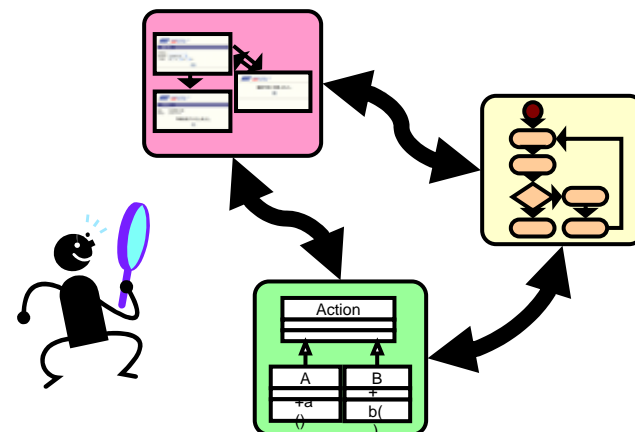
Proposed Method III, UPPAAL

Result: Deadlock-free --- False

Proprietary Secret

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Conclusion

- As a fieldwork, we proposed the methods to verify the consistency between design specifications for Web applications.
- By applying the proposed methods, we found several faults in the real specifications that had not been detected in actual reviews.
- Future work includes a full automation and an evaluation of scalability and efficiency of the proposed methods.