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Description	



# On the Outage Probability of MMSE Turbo Equalization in Frequency-Selective Rayleigh Fading Channels

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**Abstract**—We analyze the outage probability of frequency domain minimum mean squared error turbo equalization over frequency-selective Rayleigh fading channels with exponential delay power profile. A correlation chart analysis is used to evaluate the convergence property of the iterative system. Based on the convergence characteristic of the equalizer and decoder, we derive a closed form approximation to the outage probability of the turbo equalizer. Specifically, an upper bounding technique and a central limit theorem are used to show that the outage probability is well approximated by a sum of complementary Gaussian error functions. Numerical results of outage probability simulations of the turbo equalizer in channels with different delay power profiles are presented to demonstrate the accuracy of the proposed method.

## I. INTRODUCTION

Turbo equalization [1]-[4] is one of the most promising techniques, without requiring excessive computational complexity, for coded transmissions over frequency-selective fading channels. The complexity advantage of turbo equalization is due to the separation of channel equalization and decoding into two basic soft-input soft-output (SfiSfo) processing components, while such high performance can be achieved by exchanging soft information between the SfiSfo components in an iterative manner. Turbo equalization was originally proposed in [1], utilizing a maximum a posteriori probability (MAP) algorithm for iterative processing in frequency-selective fading channels. However, because of its exponentially increasing complexity, the MAP-based equalizer is only practical for simple modulation formats, like binary phase shift keying (BPSK), and for channels with few multipath components. In [2], the optimal MAP algorithm has been replaced by a low-cost alternative based on the soft canceling (SC) and minimum mean-squared error (MMSE) principle. The SC-MMSE filtering approach in [2], originally proposed for the detection of random coded code-division multiple-access (CDMA) signals, has been applied to channel equalization in [3]. In [4], a turbo equalizer for single carrier transmission over broadband channels is proposed that performs the MMSE filtering in frequency domain (FD). Due to using the computationally efficient fast Fourier transform, the SC FD-MMSE equalizer in [4] has much lower complexity than its time-domain counterpart presented in [3].

In this paper, we analyze the performance of SC FD-MMSE turbo equalization in frequency-selective Rayleigh fading channels with exponential delay power profile. Specifically, we focus on the outage probability, which is defined by the probability of unsuccessful convergence of the equalizer, given random channel realizations.

We first provide a correlation chart analysis [5], similar to the well known extrinsic information transfer (EXIT) chart analysis [6], to evaluate the convergence property of the turbo system. In fact, as shown in [7], the correlation chart method provides us with an analytical expression describing the convergence property of the equalizer. We then use the correlation functions of the equalizer and decoder to derive a closed-form approximation on the outage probability. In particular, using a union bounding technique and a central limit theorem, we show that the outage probability is well approximated by a sum of complementary Gaussian error functions.

## II. SYSTEM MODEL AND TURBO EQUALIZATION

Consider a single carrier cyclic prefix assisted block-transmission wireless communication system. The transmission scheme is based on bit interleaved coded modulation (BICM), where the information bit sequence is encoded by a rate- $r_c$  binary encoder, randomly bit-interleaved, BPSK modulated, and grouped into  $N$  ( $n = 1, \dots, N$ ) blocks

$$\mathbf{b}(n) \equiv [b_0(n), \dots, b_q(n), \dots, b_{Q-1}(n)]^T \quad (1)$$

that are transmitted over the frequency-selective fading channel.

The channel  $\mathbf{h} \equiv [h(0), \dots, h(L-1)]^T$  is composed of  $L$  taps and assumed to be constant during the transmission of one frame (comprised of  $N$  blocks), but varying randomly and independently frame-by-frame. Thus, we consider a slowly time-varying fading channel. We also assume that the  $L$  channel gains are perfectly known at the receiver.

Employing a cyclic prefix of length  $P = L - 1$  to each transmit block, the received signals can be expressed as

$$\mathbf{r}(n) = \mathbf{H}\mathbf{b}(n) + \mathbf{v}(n), n = 1, \dots, N, \quad (2)$$

where  $\mathbf{H} = \text{circ}_Q\{\mathbf{h}\}$  is the circulant channel matrix of size  $Q \times Q$ , and  $\mathbf{v}(n) \sim \mathcal{CN}(\mathbf{0}, \sigma^2 \mathbf{I})$  is the additive white Gaussian noise (AWGN). Note that  $\mathbf{H}$  may be decomposed into a diagonal matrix  $\Xi$  by the Fourier matrix,

$$\mathbf{H} = \mathbf{F}^H \Xi \mathbf{F}, \quad (3)$$

where  $\mathbf{F}$  denotes the Fourier matrix of size  $Q$ , whose  $(l, j)$ -th element is given by  $Q^{-\frac{1}{2}} e^{-i\frac{2\pi}{Q}lj}$ ,  $0 \leq l, j \leq Q-1$ , and  $\Xi = \text{diag}\{\boldsymbol{\tau}\}$  is the FD channel matrix with  $\boldsymbol{\tau}$  being a vector comprising the FD channel coefficients, i.e.,  $\boldsymbol{\tau} = \mathbf{F}^H \tilde{\mathbf{h}}$  with  $\tilde{\mathbf{h}} = [\mathbf{h}^T, \mathbf{0}_{1 \times (Q-L-1)}]^T$ . At the receiver side, iterative processing for joint equalization and decoding is performed. The receiver consists of a SC FD-MMSE equalizer and a single-user a posteriori probability decoder. Within the iterative processing, *extrinsic* log likelihood ratios (LLRs) of the coded bits are exchanged between the equalizer and decoder, following the turbo principle [2]. Inputs to the equalizer are the received signals  $\mathbf{r}(n)$  and the *a priori* LLR sequences

$$\boldsymbol{\zeta}(n) \equiv [\zeta_0(n), \dots, \zeta_q(n), \dots, \zeta_{Q-1}(n)]^T \text{ for all } n, \quad (4)$$

where  $\zeta_q(n) \equiv \log(\text{Prob}(b_q(n) = +1)/\text{Prob}(b_q(n) = -1))$ . The equalizer first generates an estimate of the received signal  $\mathbf{r}(n)$  and then subtracts it from  $\mathbf{r}(n)$ , yielding the residual:

$$\tilde{\mathbf{r}}(n) = \mathbf{r}(n) - \mathbf{H}\bar{\mathbf{b}}(n) \quad (5)$$

where  $\bar{\mathbf{b}}(n) = [\bar{b}_0(n), \dots, \bar{b}_{Q-1}(n)]^T$  is a vector comprising the expected values of the elements in  $\mathbf{b}(n)$ , i.e.,

$$\bar{\mathbf{b}}(n) = \tanh\left(\frac{\boldsymbol{\zeta}(n)}{2}\right). \quad (6)$$

After the residual  $\tilde{\mathbf{r}}(n)$  is calculated, linear adaptive FD filtering is performed to extract the desired signal components  $\mathbf{b}(n)$  as [3] [4]

$$\mathbf{z}(n) = (1 + \gamma\varphi)^{-1} (\gamma\bar{\mathbf{b}}(n) + \mathbf{F}^H \boldsymbol{\Psi} \mathbf{F} \tilde{\mathbf{r}}(n)), \quad (7)$$

where  $\mathbf{z}(n) = [z_0(n), \dots, z_q(n), \dots, z_{Q-1}(n)]^T$ ,  $\varphi = \frac{1}{QN} \sum_{n=1}^N \bar{\mathbf{b}}^T(n) \bar{\mathbf{b}}(n)$  is the mean energy of the symbol estimates in  $\bar{\mathbf{b}}(n)$ ,  $\boldsymbol{\Psi} = \Xi^H [(1 - \varphi)\Xi\Xi^H + \sigma^2 \mathbf{I}]^{-1}$  is the frequency domain filter, and  $\gamma = (1/Q)\text{Trace}\{\boldsymbol{\Psi}\Xi\}$ . The equalizer then computes the extrinsic LLR for each transmitted bit  $b_q(n)$  as

$$\lambda_q(n) \equiv \log\left(\frac{\text{Prob}(z_q(n)|b_q(n) = +1)}{\text{Prob}(z_q(n)|b_q(n) = -1)}\right). \quad (8)$$

Following [3], we may approximate  $z_q(n)$  by an equivalent AWGN channel having  $b_q(n)$  as its input, so that

$$\text{Prob}(z_q(n)|b_q(n) = b) \sim \mathcal{N}(\mu_b b, \sigma_b^2), b \in \{\pm 1\}, \quad (9)$$

where  $\mu_b = \gamma/(1 + \gamma\varphi)$ , and  $\sigma_b^2 = \mu_b(1 - \mu_b)$ .

Note that during the first iteration of turbo equalization,  $\zeta_q(n)$  is zero for all  $n, q$ , and later on  $\zeta_q(n)$  is provided via the interleaver in the form of extrinsic LLRs of the decoder.

### A. Channel Model and Correlation Properties

We consider Rayleigh block-fading channels, where the channel coefficients  $h(l)$ ,  $l = 0, \dots, L-1$  are assumed to be independent and identically distributed (i.i.d) circularly-symmetric complex Gaussian random variables. Furthermore, we presume that the delay autocorrelation function of the channel  $\mathbf{h}$  is described by an exponential delay power profile with normalized root mean square (rms) delay  $\tau_d$ ,

$$p(l) = v \exp\left(-\frac{l}{\tau_d}\right), \text{ for } l = 0, \dots, L-1, \quad (10)$$

where  $v = L/\sum_l p(l)$  is a normalization constant.

Let us write the  $q_1$ th and  $q_2$ th FD channel coefficients as

$$\begin{aligned} \tau_{q_1} &= \tau_{c,q_1} + i\tau_{s,q_1}, \\ \tau_{q_2} &= \tau_{c,q_2} + i\tau_{s,q_2}. \end{aligned} \quad (11)$$

In (11), the random variables  $\tau_{c,q_1}$ ,  $\tau_{s,q_1}$ ,  $\tau_{c,q_2}$  and  $\tau_{s,q_2}$  are identically zero-mean Gaussian distributed with variance  $\eta/2$ , where  $\eta = \frac{L}{Q}$ . Assume that, as  $L$  and  $Q$  increase, the normalized frequency separation  $\Delta f \equiv \tau_d/Q$  remains fixed. Following [8], we then may write the cross-correlations of the random variables in (11) for  $L, Q \rightarrow \infty$  as

$$\begin{aligned} \mathbb{E}[\tau_{c,q_1} \tau_{s,q_1}] &= \mathbb{E}[\tau_{c,q_2} \tau_{s,q_2}] = 0 \\ \mathbb{E}[\tau_{c,q_1} \tau_{c,q_2}] &= \mathbb{E}[\tau_{s,q_1} \tau_{s,q_2}] = \frac{\eta/2}{1 + (2\pi\Delta f \Delta q)^2} \\ \mathbb{E}[\tau_{c,q_1} \tau_{s,q_2}] &= -\mathbb{E}[\tau_{c,q_2} \tau_{s,q_1}] = -\frac{\eta\pi\Delta f \Delta q}{1 + (2\pi\Delta f \Delta q)^2}, \end{aligned} \quad (12)$$

where  $\Delta q = |q_1 - q_2|$ . From (12), we observe that  $(\tau_{c,q_1}, \tau_{s,q_1})$  and  $(\tau_{c,q_2}, \tau_{s,q_2})$  form a *circular pair* [9]. Therefore, the exponential distributed FD channel gains  $\kappa_q = \tau_{c,q}^2 + \tau_{s,q}^2$  with mean  $\eta$  and variance  $\eta^2$  for all  $q = 0, \dots, Q-1$ , have the correlation coefficient [9]

$$\delta_{\Delta q} = \frac{1}{1 + (2\pi\Delta f \Delta q)^2}. \quad (13)$$

### III. CONVERGENCE CHARACTERISTIC OF THE TURBO EQUALIZER

In this section, the correlation chart analysis [5] is used to study the convergence characteristic of the SC FD-MMSE turbo equalizer.

Let  $\varphi_e \equiv \mathbb{E}[b_q(n)a_q(n)]$  be the correlation between the true binary transmit signal  $b_q(n)$  and the MMSE estimate  $a_q(n) \equiv \mathbb{E}[b_q(n)|\lambda_q(n)] = \tanh((1/2)\lambda_q(n))$  of  $b_q(n)$  given  $\lambda_q(n)$ . Following [6], we model  $(1/2)b_q(n)\lambda_q(n)$  as i.i.d. Gaussian random variables. We assume that the *symmetry condition* [6] for all LLR messages is satisfied, so that  $p(x) = p(-x)\exp(x)$ , where  $p(x)$  is the probability density function (PDF) of an LLR message. By enforcing this condition on the random variables  $(1/2)b_q(n)\lambda_q(n)$ , we obtain

$$\mathbb{E}\left[\frac{1}{2}b_q(n)\lambda_q(n)\right] = \text{Var}\left[\frac{1}{2}b_q(n)\lambda_q(n)\right] = \frac{2\mu_b}{1 - \mu_b}. \quad (14)$$

Similarly, let  $\varphi_d \equiv \mathbb{E}[b_q(n)c_q(n)]$  be the correlation between  $b_q(n)$  and  $c_q(n) = \tanh((1/2)\zeta_q(n))$ . Then using (7), we can

express the effective signal-to-noise ratio (SNR)  $\Psi$  at the equalizer output as

$$\begin{aligned}\Psi(\varphi_d) &= \frac{2\mu_b}{1-\mu_b} \\ &= \frac{2\gamma}{1-\gamma(1-\varphi_d)} \\ &= \frac{2}{1-\varphi_d} \left[ Q \left( \sum_{q=0}^{Q-1} \frac{1}{1+\rho\kappa_q} \right)^{-1} - 1 \right],\end{aligned}\quad (15)$$

where  $\rho \equiv (1-\varphi_d)/\sigma^2$ . Under the Gaussian assumption, the correlation  $\varphi_e$  is then given by

$$\varphi_e = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \tanh(z\sqrt{\Psi(\varphi_d)} + \Psi(\varphi_d)) e^{-\frac{z^2}{2}} dz, \quad (16)$$

$$\equiv \phi(\Psi(\varphi_d)), \quad (17)$$

$$\equiv f_e(\varphi_d), \quad (18)$$

where we have defined  $\phi(x) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \tanh(z\sqrt{x} + x) e^{-\frac{z^2}{2}} dz$ . The function  $f_e(\cdot)$  in (18) is referred to as the correlation characteristic of the equalizer in what follows. Similarly, the correlation  $\varphi_d$  can be expressed as

$$\varphi_d = f_d(\varphi_e), \quad (19)$$

where  $f_d(\cdot)$  is denoted as the correlation characteristic of the decoder. We obtain  $f_d$  by a Monte Carlo method, as described in [5].

With the definitions in (18) and (19), the convergence behavior of the turbo equalizer can now be analyzed by evaluating the correlation sequence  $\{\varphi_e^{(l)}, \varphi_d^{(l)}\}$ ,  $l = 0, \dots, T$ , over  $T$  iterations between the equalizer and decoder, generated by

$$\varphi_e^{(l)} = f_e(\varphi_d^{(l)}) \text{ with } \varphi_d^{(0)} = 0, \quad (20)$$

$$\varphi_d^{(l+1)} = f_d(\varphi_e^{(l)}). \quad (21)$$

The functions  $f_e(\cdot)$  and  $f_d(\cdot)$  are bounded and monotonically increasing in  $\varphi_d$  and  $\varphi_e$ , respectively. Thus, the sequence  $\{\varphi_e^{(l)}, \varphi_d^{(l)}\}$  converges asymptotically to fixed values  $\{\tilde{\varphi}_e, \tilde{\varphi}_d\}$  with  $\tilde{\varphi}_e = \lim_{l \rightarrow \infty} \varphi_e^{(l)}$  and  $\tilde{\varphi}_d = \lim_{l \rightarrow \infty} \varphi_d^{(l)}$ , for  $T \rightarrow \infty$ . The function  $f_e(\cdot)$  is shown for an example snapshot in Fig. 1. Also shown is the inverse function of  $f_d(\cdot)$ , denoted as  $f_{I,d}(\cdot)$ , and the decoding trajectory that visualizes the correlation exchange between the equalizer and decoder. The function  $f_{I,d}(\cdot)$  has been calculated for a rate-1/2 memory-three convolutional code.

Convergence of turbo equalization is achieved when  $\varphi_d^{(l)}$  attains the maximum value  $\tilde{\varphi}_d = 1$ . Obviously, this is possible for  $T$  being sufficiently large, if the following constraint holds:

$$f_e(\varphi_d) > f_{I,d}(\varphi_d), \quad \forall \varphi_d \in [0, 1]. \quad (22)$$

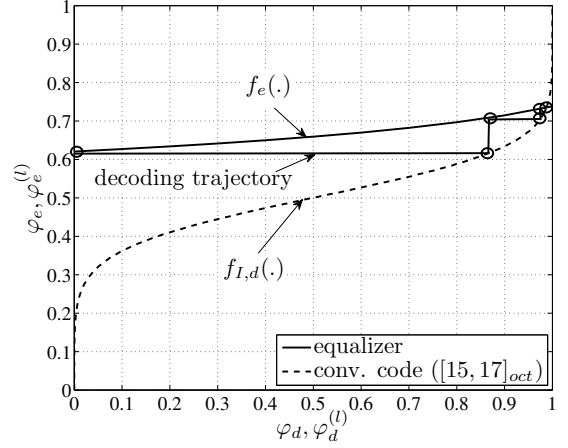


Fig. 1. Equalizer and decoder correlation characteristics for a random channel realization ( $L = 32$ ,  $\tau_d = 5$ ) at  $E_b/N_0 = 3$  dB.

#### IV. OUTAGE PROBABILITY ANALYSIS

The turbo equalizer is in outage if for the specific channel realization  $\tilde{\varphi}_d < 1$ . Hence, an outage event  $\mathcal{O}$  occurs when  $f_e(\varphi_d) \leq f_{I,d}(\varphi_d)$  for at least one value of  $\varphi_d \in [0, 1]$ . Thus, we can define the outage probability of the turbo equalizer as

$$\text{Prob}(\mathcal{O}) \equiv \text{Prob}(f_e(\varphi_d) \leq f_{I,d}(\varphi_d), \exists \varphi_d \in [0, 1]). \quad (23)$$

A direct evaluation of the constraint  $f_e(\varphi_d) \leq f_{I,d}(\varphi_d)$  in (23) on the continuous interval  $\varphi_d \in [0, 1]$  is computationally intractable. Therefore, we proceed by imposing the constraint on a discrete set of  $D$  values  $\{\varphi_{d,k}\}$ ,  $k = 1, \dots, D$ . Then, using the union bound [13], we find an approximate upper bound on (23) as

$$\text{Prob}(\mathcal{O}) \leq \sum_{k=1}^D \text{Prob}(f_e(\varphi_{d,k}) \leq f_{I,d}(\varphi_{d,k})). \quad (24)$$

Using (17) and (15), we can write (24) as

$$\begin{aligned}\text{Prob}(\mathcal{O}) &\leq \sum_{k=1}^D \text{Prob}(\phi(\Psi(\varphi_{d,k})) \leq f_{I,d}(\varphi_{d,k})) \\ &= \sum_{k=1}^D \text{Prob}\left(\frac{1}{Q} \sum_{q=0}^{Q-1} \frac{1}{1+\rho_k \kappa_q} > A_k\right),\end{aligned}\quad (25)$$

where

$$\rho_k = (1 - \varphi_{d,k})/\sigma^2, \quad (26)$$

$$A_k = \frac{1}{(1 - \varphi_{d,k})\phi_I(f_{I,d}(\varphi_{d,k})) + 2} \quad (27)$$

with  $\phi_I(\cdot)$  being the inverse function of  $\phi(\cdot)$ . For convenience, we define

$$\begin{aligned}S_k &\equiv \frac{1}{Q} \sum_{q=0}^{Q-1} \frac{1}{1 + \rho_k \kappa_q} \\ &= \frac{1}{Q} \sum_{q=0}^{Q-1} c_k(\kappa_q) = \frac{1}{Q} \sum_{q=0}^{Q-1} s_k(\tau_{c,q}, \tau_{s,q}),\end{aligned}\quad (28)$$

where  $c_k(\kappa_q) = s_k(\tau_{c,q}, \tau_{s,q}) \equiv \frac{1}{1 + \rho_k \kappa_q}$ . Exact calculation of the distribution of  $S_k$  in (28) is not easy, since as shown in Section II-A, the FD channel coefficients  $\kappa_q$  are correlated with order  $1/(\Delta q)^2$ . We use a theorem from *Arcones* [10] to show that  $S_k$  is asymptotically Gaussian distributed. Note that this theorem has also been adapted in [11] to calculate the capacity of OFDM systems.

We state *Arcones'* theorem below. A proof of it can be found in [10].

*Theorem 1:* Let  $\{\mathbf{X}_j\}$ ,  $\mathbf{X}_j \equiv [X_j^{(1)}, \dots, X_j^{(d)}]^T$ ,  $0 \leq j < \infty$  be a stationary zero-mean sequence of Gaussian random vectors in  $\mathbb{R}^d$  with covariance function

$$r^{(i,l)}(k) = \mathbb{E}[X_m^{(i)} X_{m+k}^{(l)}] \quad (29)$$

for  $k \in \mathbb{Z}$ ,  $0 \leq m < \infty$  and  $m+k \geq 0$ . Let  $f: \mathbb{R}^d \rightarrow \mathbb{R}$  be a real-valued function with Hermite rank  $\nu(f)$  such that  $1 \leq \nu(f) < \infty$ . Suppose that

$$\sum_{k=-\infty}^{\infty} |r^{(i,l)}(k)|^{\nu(f)} < \infty \quad (30)$$

for  $1 \leq i, l \leq d$ . Then, as  $Q$  tends to infinite

$$\frac{1}{\sqrt{Q}} \sum_{j=0}^{Q-1} (f(\mathbf{X}_j) - \mathbb{E}[f(\mathbf{X}_j)]) \xrightarrow{d} \mathcal{N}(0, \sigma_{\{\mathbf{X}_j\}}^2), \quad (31)$$

where  $\xrightarrow{d}$  denotes convergence in distribution, and

$$\begin{aligned} \sigma_{\{\mathbf{X}_j\}}^2 &= \mathbb{E}[(f(\mathbf{X}_0) - \mathbb{E}[f(\mathbf{X}_0)])^2] \\ &+ 2 \sum_{j=0}^{\infty} \mathbb{E}[(f(\mathbf{X}_0) - \mathbb{E}[f(\mathbf{X}_0)])(f(\mathbf{X}_k) - \mathbb{E}[f(\mathbf{X}_k)])]. \end{aligned} \quad (32)$$

We apply Theorem 1 to the sequence of Gaussian FD channel coefficients  $\{\mathbf{X}_q\}$ ,  $0 \leq q \leq Q-1$  with  $\mathbf{X}_q = [\tau_{c,q}, \tau_{s,q}]^T$ ,  $d = 2$ . In [11], it was shown that for a zero-mean  $d = 2$  stationary Gaussian sequence with the correlation properties described in (12), the condition in (30) is satisfied if the Hermite rank of the function  $f(\cdot)$  is at least two. Let  $f(\mathbf{X}_q) = s_k(\tau_{c,q}, \tau_{s,q})$ . In Appendix A, we show that  $s_k(\cdot)$  has Hermite rank  $\nu(s_k) \geq 2$ . Requirement (30) is then satisfied, so that

$$\begin{aligned} \frac{1}{\sqrt{Q}} \sum_{q=0}^{Q-1} (s_k(\tau_{c,q}, \tau_{s,q}) - \mathbb{E}[s_k(\tau_{c,q}, \tau_{s,q})]) \\ \xrightarrow{d} \mathcal{N}(0, \sigma_{\{s_k\}}^2). \end{aligned} \quad (33)$$

Therefore, for large finite  $Q$ , the distribution of  $S_k$  in (28) may be approximated by a Gaussian random variable having mean  $\mu_{S_k} = \mathbb{E}[c_k(\kappa_q)]$  and variance

$$\begin{aligned} \sigma_{S_k}^2 &= \text{Var}[S_k] \\ &= \frac{1}{Q} \text{Var}[c_k(\kappa_q)] + \frac{2}{Q^2} \sum_{j=1}^{Q-1} (Q-j) \text{Cov}[c_k(\kappa_0), c_k(\kappa_j)]. \end{aligned} \quad (34)$$

The quantities  $\mathbb{E}[c_k(\kappa_q)]$  and  $\text{Var}[c_k(\kappa_q)]$  can be calculated as

$$\begin{aligned} \mathbb{E}[c_k(\kappa_q)] &= \int_0^{\infty} \frac{1}{1 + \rho_k \kappa_q} p(\kappa_q) d\kappa_q, \\ \text{Var}[c_k(\kappa_q)] &= \int_0^{\infty} \left( \frac{1}{1 + \rho_k \kappa_q} - \mathbb{E}[c_k(\kappa_q)] \right)^2 p(\kappa_q) d\kappa_q, \end{aligned}$$

where  $p(\kappa_q)$  is the PDF of  $\kappa_q$ . Since  $\kappa_q$  is exponential distributed,  $p(\kappa_q) = 1/\eta \exp(-\kappa_q/\eta)$ . Therefore,  $\mathbb{E}[c_k(\kappa_q)]$  and  $\text{Var}[c_k(\kappa_q)]$  may be written as

$$\begin{aligned} \mathbb{E}[c_k(\kappa_q)] &= \frac{\eta}{\rho_k} e^{\frac{\eta}{\rho_k}} E_1\left(\frac{\eta}{\rho_k}\right), \quad (35) \\ \text{Var}[c_k(\kappa_q)] &= \frac{\eta}{\rho_k} \left[ 1 - \frac{\eta}{\rho_k} e^{\frac{\eta}{\rho_k}} E_1\left(\frac{\eta}{\rho_k}\right) \left( 1 + e^{\frac{\eta}{\rho_k}} E_1\left(\frac{\eta}{\rho_k}\right) \right) \right], \quad (36) \end{aligned}$$

where  $E_1(x) = \int_x^{\infty} e^{-t}/t dt$  denotes the exponential integral function. A derivation of  $\text{Cov}[c_k(\kappa_0), c_k(\kappa_j)]$  in (34) is given in Appendix B. Using (34) and (35), we may express (24) as

$$\begin{aligned} \text{Prob}(\mathcal{O}) &\leq \sum_{k=1}^D \frac{1}{\sqrt{2\pi}\sigma_{S_k}} \int_{A_k}^{\infty} \exp\left[-\frac{(x - \mu_{S_k})^2}{2\sigma_{S_k}^2}\right] dx \\ &= \frac{1}{2} \sum_{k=1}^D \text{erfc}\left(\frac{A_k - \mu_{S_k}}{\sqrt{2}\sigma_{S_k}}\right), \end{aligned} \quad (37)$$

where  $\text{erfc}(x) = \frac{2}{\sqrt{\pi}} \int_x^{\infty} e^{-t^2} dt$  is the complementary Gaussian error function.

## V. NUMERICAL RESULTS

In this section, we provide results of simulations for the outage performance of the SC FD-MMSE turbo equalizer and the union bounding technique from Section IV. We consider a single carrier CP-assisted block-transmission system with each block having  $Q = 128$  BPSK symbols. The binary encoder at the transmitter is a serially concatenated convolutional code (SCCC), consisting either of a rate-1/2 or rate-3/4 outer encoder and a recursive rate-1 inner encoder. The systematic rate-1/2, memory-4 code is defined by the generator  $(g_r, g_0) = (23, 35)$ , where  $g_r$  denotes the feedback polynomial. The higher rate-3/4 code is obtained by puncturing, as specified in [14]. The recursive rate-1 inner encoder has generator polynomials  $(g_r, g_0) = (3, 2)$ . We evaluate the outage performance for Rayleigh block-fading channels with  $L = 32$  path components and exponential delay power profiles with  $\tau_d = 5$  and  $\tau_d = 8$ . The length of a frame is fixed to  $NQ = 8192$  BPSK symbols, and thus the channel is assumed to be constant over  $N = 64$  transmitted blocks. The CP length is set to  $P = 31$ . The turbo equalizer performs 10 iterations between the equalizer and the SCCC decoder, and 20 iterations between the inner and outer channel decoder. For the calculation of the union bound in (37), the constraint in (24) is computed on a grid of  $D = 5$  points, such that  $\varphi_{d,1} = 0.01$ ,  $\varphi_{d,2} = 0.3$ ,  $\varphi_{d,3} = 0.6$ ,  $\varphi_{d,4} = 0.9$ , and  $\varphi_{d,5} = 0.99$ .

The outage performances of the SC FD-MMSE turbo equalizer and the union bound in (37) for transmissions over

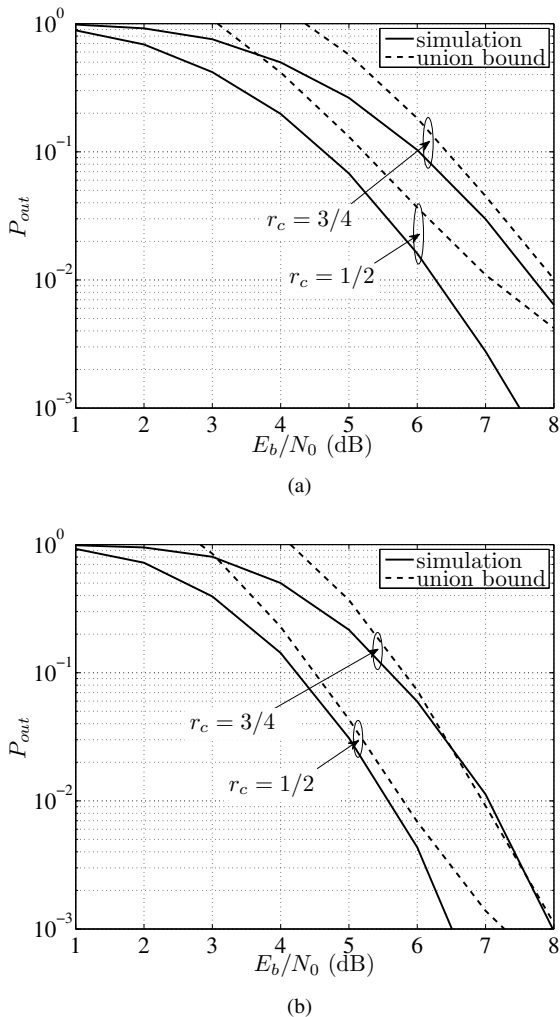


Fig. 2. Outage probability and union bound of the SC FD-MMSE turbo equalizer utilizing SCCCs with rates  $r = 1/2$  and  $3/4$  for  $L = 32$ -tap Rayleigh fading channel having an exponential delay power profile with normalized rms delay (a)  $\tau_d = 5$  and (b)  $\tau_d = 8$ .

$\tau_d = 5$  and  $\tau_d = 8$  channels are shown in Fig. 2 (a) and 2 (b), respectively. The outage probabilities ( $P_{out}$ ) have been computed by averaging over 50000 random channel realizations. We observe that the performance improves with increasing values of rms delay  $\tau_d$  due to increasing channel diversity. Also, observe that we obtain a reasonable analytical approximation of the outage probability using the proposed union bounding technique for all configurations, where the performance gap between simulation and bound is smaller than  $E_b/N_0 = 0.8$  dB for  $P_{out} = 10^{-2}$ .

## VI. CONCLUSION

We have considered the performance of FD SC-MMSE turbo equalization over frequency-selective Rayleigh fading channels with exponential delay power profile. A correlation analysis has been provided to evaluate the convergence property of the turbo equalization system. Using the correlation functions of the equalizer and decoder, we have derived

a closed form expression to the outage probability of the turbo equalizer. Numerical results show that the technique presented in this paper yields a reasonable approximation of the outage performance of the FD SC-MMSE turbo equalizer in frequency-selective Rayleigh fading channels.

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## APPENDIX A

Let  $\mathbf{X}$  be a Gaussian vector with zero-mean and unit-variance. Further, let  $f : \mathbb{R}^d \rightarrow \mathbb{R}$  be a measurable function with  $E[f(\mathbf{X})^2] < \infty$ . The Hermite rank  $\nu(f)$  of  $f$  with respect to  $\mathbf{X}$  is defined as [10]

$$\nu(f) \equiv \inf \left\{ \tau : \exists \text{ polynomial } P \text{ of degree } \tau \text{ with } E[(f(\mathbf{X}) - E[f(\mathbf{X})])P(\mathbf{X})] \neq 0 \right\}. \quad (38)$$

In the following, we show that the Hermite rank of  $s_k(\tau_{c,q}, \tau_{s,q})$  is at least 2. We basically follow the same derivation as in [11], where the Hermite rank of the sub-channel capacity as a function of the channel gains for an OFDM system has been calculated.

First, we show that  $\nu(s_k) \neq 0$ . Consider a zero-order polynomial  $P(\mathbf{X}) = A_0$  with  $\mathbf{X} = (X_1, X_2)$ . Then, the condition in (38) becomes

$$\begin{aligned} E[(s_k(\mathbf{X}) - E[s_k(\mathbf{X})])P(\mathbf{X})] \\ = A_0 E[s_k(\mathbf{X})] - A_0 E[s_k(\mathbf{X})] \\ = 0, \end{aligned} \quad (39)$$

for all  $A_0 \in \mathbb{R}$ , and thus  $\nu(s_k) \neq 0$ . Next, consider a first-order polynomial  $P(\mathbf{X}) = A_0X_1 + A_1X_2 + A_3$ . Then, we have

$$\begin{aligned} E[(s_k(\mathbf{X}) - E[s_k(\mathbf{X})])P(\mathbf{X})] \\ = A_0E[X_1s_k(\mathbf{X})] + A_1E[X_2s_k(\mathbf{X})], \end{aligned} \quad (40)$$

where we have used the property  $E[X_1] = E[X_2] = 0$ . The random variables  $X_1$  and  $X_2$  are identically Gaussian distributed. Thus, we get

$$\begin{aligned} E[X_1s_k(\mathbf{X})] &= E[X_2s_k(\mathbf{X})] \\ &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} x_1s_k(x_1, x_2)p(x_1, x_2)dx_1dx_2, \end{aligned} \quad (41)$$

where  $p(x_1, x_2)$  is the joint probability function of the two correlated Gaussian random variables  $X_1$  and  $X_2$ . Observe that  $p(x_1, x_2)$  and  $s_k(x_1, x_2)$  are even in  $(X_1, X_2)$ . Thus, the integral in (41) is zero and it follows that  $E[X_1s_k(\mathbf{X})] = E[X_2s_k(\mathbf{X})] = 0$ , and therefore,  $\nu(s_k) \neq 1$ . We conclude that  $\nu(s_k)$  is at least 2.

## APPENDIX B

The covariance  $\text{Cov}[c_k(\kappa_0), c_k(\kappa_j)]$  between  $c_k(\kappa_0)$  and  $c_k(\kappa_j)$  for  $1 \leq j \leq Q - 1$  is given by

$$\text{Cov}[c_k(\kappa_0), c_k(\kappa_j)] = E[c_k(\kappa_0), c_k(\kappa_j)] - \mu_{S_k}^2, \quad (42)$$

where

$$\begin{aligned} E[c_k(\kappa_0), c_k(\kappa_j)] &= \int_0^{\infty} \int_0^{\infty} \frac{1}{1 + \rho_k x} \\ &\quad \times \frac{1}{1 + \rho_k y} f_{\kappa_0, \kappa_j}(x, y) dx dy. \end{aligned} \quad (43)$$

In (43),  $f_{\kappa_0, \kappa_j}(x, y)$  is the joint PDF of  $\kappa_0$  and  $\kappa_j$ , which follows a bivariate exponential distribution [9] with the form:

$$f_{\kappa_0, \kappa_j}(x, y) = \alpha \exp[-\beta(x + y)] I_0(\theta\sqrt{xy}), \quad (44)$$

where

$$\alpha = \frac{1}{\eta^2(1 - \delta_j)}, \quad \beta = \alpha\eta, \quad \theta = 2\alpha\eta\sqrt{\delta_j}$$

and  $I_0(\cdot)$  denotes the modified zero-order Bessel function of the first kind. We may substitute (44) into (43) to obtain

$$\begin{aligned} E[c_k(\kappa_0), c_k(\kappa_j)] &= \alpha \int_0^{\infty} \int_0^{\infty} \frac{1}{1 + \rho_k x} \frac{1}{1 + \rho_k y} \exp[-\beta(x + y)] \\ &\quad \times I_0(\theta\sqrt{xy}) dx dy \\ &= \alpha \sum_{n=0}^{\infty} \left( \frac{\theta^{2n}}{4^n (n!)^2} \left[ \int_0^{\infty} \frac{x^n}{1 + \rho_k x} \exp(-\beta x) dx \right]^2 \right). \end{aligned} \quad (45)$$

Note that (46) is obtained by using the series expansion of  $I_0(\cdot)$  [12]. Using a substitution from [12], we may express the

integral expression in (46) as

$$\begin{aligned} &\int_0^{\infty} \frac{x^n}{1 + \rho_k x} \exp(-\beta x) dx \\ &= \frac{(-1)^n}{\rho_k^{n+1}} \left[ \exp\left(\frac{\beta}{\rho_k}\right) E_1\left(\frac{\beta}{\rho_k}\right) + \sum_{s=1}^n (s-1)! \left(-\frac{\rho_k}{\beta}\right)^s \right]. \end{aligned} \quad (47)$$

Substituting (47) into (46), we get

$$\begin{aligned} E[c_k(\kappa_0), c_k(\kappa_j)] &= \frac{\alpha}{\rho_k^2} \sum_{n=0}^{\infty} \left( \frac{\left(\frac{\theta}{2\rho_k}\right)^{2n}}{(n!)^2} \left[ \exp\left(\frac{\beta}{\rho_k}\right) E_1\left(\frac{\beta}{\rho_k}\right) \right. \right. \\ &\quad \left. \left. + \sum_{s=1}^n (s-1)! \left(-\frac{\rho_k}{\beta}\right)^s \right]^2 \right). \end{aligned} \quad (48)$$

Numerical calculations show that the series in (48) is rapidly convergent. The series representation (48) may therefore be used to efficiently calculate (42).